

High-performance Computing for Flows in Porous Media

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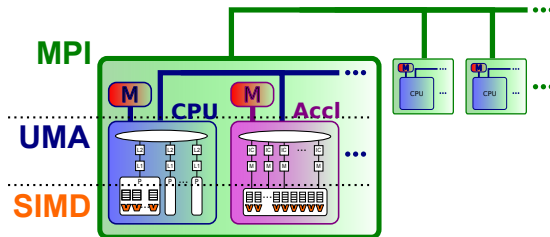
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Trends in Computer Architecture

- Power wall
 - ▶ Power consumption is the limiting factor for exascale computing!
 - ▶ Tianhe 2: 33,8 PF (Linpack), 17.8 MW $\Rightarrow$ 500 MW
 - ▶ Clock rate stagnates but Moore's law is still valid
- Memory wall
 - ▶ Bandwidth not sufficient to sustain anywhere near peak performance
 - ▶ (Still) significant latency for main memory access
 - ▶ Memory access is energetically costly
 - ▶ Memory is hierarchical: register, cache, local, remote, ...
- ILP wall
 - ▶ Automatic extraction of instruction level parallelism stagnates
 - ▶ Revival of vectorization in form of SIMD instructions
- Other
 - ▶ Heterogeneous node architectures: CPU + coprocessors
 - ▶ Reduced reliability due to immense number of components?

EXA-SCALE $\Rightarrow 10^6$ cores $\times 10^3$ threads $\times 10^9$ GFLOPS

Typical Architecture



(Figure by Christian Engwer)

- Three levels of parallelism
 - ▶ Instruction level parallelism: SIMD vector units
 - ▶ Multithreading/shared memory: multiple sockets, multiple cores
 - ▶ Message passing: nodes connected by fast interconnect
- Coprocessor architectures
 - ▶ Triggered by economical and then power considerations
 - ▶ CPUs are catching up: E5-2699v3: 0.9TF@145W, K40: 1.4TF@235W
 - ▶ My opinion: coprocessor architectures will not last ...

What Does This Mean for PDE Numerics?



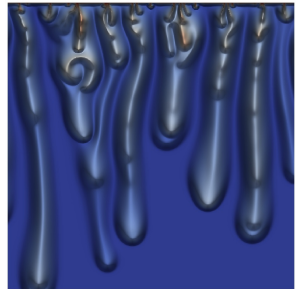
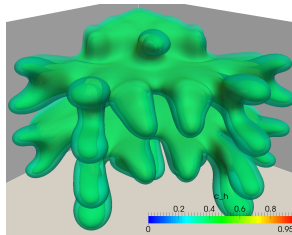
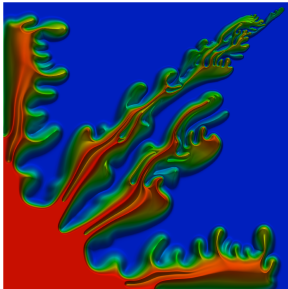
- Exploit hybrid parallelism:
MPI + shared memory + SIMD,
i.e. Intel TBB + auto vectorization
- Transparent accelerator support (concentrate on Intel Phi)
- FLOPS are for free, memory access is expensive
- Matrix-based methods
 - ▶ Robust preconditioners available (AMG)
 - ▶ Poor performance, bound by memory bandwidth
- Matrix-free methods
 - ▶ Restricted to explicit time-stepping or certain solvers, e.g. CG, GMG
 - ▶ Low-order FE schemes:
 - ★ Exploit problem structure: linear transformation, regular grids, constant coefficients, ...
 - ★ Vectorization over several elements
 - ▶ High-order FE schemes:
 - ★ Complexity reduction through tensor-product approaches
 - ★ Medium high order: Vectorization within one element
 - ★ Very high order: Several threads to one element

Unstable Flows in Porous Media

$$\nabla \cdot v = 0, \quad v = -\frac{K}{\mu(c)}(\nabla p - \rho(c)g)$$

$$\partial_t(\Phi c) + \nabla \cdot (cv - D(v)\nabla c) = 0$$

Density driven flow, miscible displacement
Can exploit high resolution schemes



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Discontinuous Galerkin (DG) Finite Element Methods

- Advantages

- ▶ Potentially higher order of convergence
- ▶ Higher-order methods beneficial even without sufficient regularity
- ▶ Element-wise conservative
- ▶ Full permeability tensors, discontinuous coefficients
- ▶ Handles elliptic, parabolic and hyperbolic equations equally well
- ▶ Unstructured grids, nonconforming refinement
- ▶ Contiguous memory access
- ▶ Block matrix structure: less indirect access, BLAS level 3

- Disadvantages

- ▶ No monotonicity
- ▶ Free parameter to tune
- ▶ Higher number of DOFs compared to standard FE or FV
- ▶ Linear systems are more difficult to solve

DG for Flow Equation

$$\begin{aligned}
 a_h(p, w; c) &= \sum_{T \in \mathcal{T}_h} \int_T [\mu(c)^{-1} K(\nabla p - \rho(c)g)] \cdot \nabla w - \sum_{F \in \mathcal{F}_h^{iD}} \int_F \nu_F \cdot \{\mu(c)^{-1} K(\nabla p - \rho(c)g)\}_\omega \llbracket w \rrbracket \\
 &\quad - \sum_{F \in \mathcal{F}_h^{iD}} \int_F \theta \nu_F \cdot \{\mu(c)^{-1} K \nabla w\}_\omega \llbracket p \rrbracket + \sum_{F \in \mathcal{F}_h^{iD}} \int_F \gamma_{F,f} \llbracket w \rrbracket \llbracket p \rrbracket \\
 l(w) &= \sum_{T \in \mathcal{T}_h} \int_T q_f w - \sum_{F \in \mathcal{F}_h^N} \int_F j_f w \quad \theta = 1 \text{ (SIPG)}, \theta = 0 \text{ (OBB)}, \theta = -1 \text{ (NIPG)}
 \end{aligned}$$

SIPG + weighted averages (*Di Pietro, Ern, Guermond '2008*):

$$\delta_{Kn}^\pm = \nu_F^t K^\pm \nu_F / \mu(c)^\pm, \quad \omega^- = \frac{\delta_{Kn}^+}{\delta_{Kn}^- + \delta_{Kn}^+}, \quad \omega^+ = \frac{\delta_{Kn}^-}{\delta_{Kn}^- + \delta_{Kn}^+}.$$

Penalty term ($\langle \cdot, \cdot \rangle$ denotes harmonic average, k polynomial degree):

$$\gamma_{F,f} = \alpha \langle \delta_{Kn}^-, \delta_{Kn}^+ \rangle M, \quad M = k(k+n-1) \frac{|F|}{\min(|T^-(F)|, |T^+(F)|)}$$

Transport Equation

$$b_h(c, z; v) = \partial_t \int_{\Omega} \Phi c z + \sum_{T \in \mathcal{T}_h} \int_T (D \nabla c - v c) \cdot \nabla z + \sum_{F \in \mathcal{F}_h^{iD}} \int_F c^{\uparrow} \nu_F \cdot v \llbracket z \rrbracket$$

$$- \sum_{F \in \mathcal{F}_h^{iD}} \int_F \nu_F \cdot \{D \nabla c\}_{\omega} \llbracket z \rrbracket - \theta \sum_{F \in \mathcal{F}_h^{iD}} \int_F \nu_F \cdot \{D \nabla z\}_{\omega} \llbracket c \rrbracket + \sum_{F \in \mathcal{F}_h^{iD}} \int_F \gamma_{F,t} \llbracket z \rrbracket \llbracket c \rrbracket$$

$$r(z) = \sum_{T \in \mathcal{T}_h} \int_T q_t z - \sum_{F \in \mathcal{F}_h^{Nn}} \int_F j_t z$$

$$m(c, z) = \int_{\Omega} \Phi c z$$

Upwinding, weighting and penalty parameter:

$$c^{\uparrow} = \begin{cases} c^- & \nu_F \cdot v \geq 0 \\ c^+ & \text{else} \end{cases}$$

$$\delta_{Dn}^{\pm} = \nu_F^t D^{\pm} \nu_F, \quad \omega^{\pm} = \frac{\delta_{Dn}^{\mp}}{\delta_{Dn}^- + \delta_{Dn}^+}, \quad \gamma_{F,t} = \langle \delta_{Dn}^-, \delta_{Dn}^+ \rangle M.$$

Decoupled Solution Algorithm

- 1: Set time t , substeps S
- 2: Initialize $c_h(t)$
- 3: **while** $t < t_{end}$ **do**
- 4: Solve flow equation $a_h(p_h, w; c_h(t)) = l(w) \quad \forall w \in W_h^{k_1}$
- 5: Reconstruct Darcy velocity $\hat{v} \in \hat{V} \subset H(\text{div}; \Omega)$ from p_h and data
- 6: Choose appropriate time step Δt , set $\tau = \Delta t/S$
- 7: **for** $i = 1$ to S **do**
- 8: Given $c_h(t + (i-1)\tau)$ compute $c_h(t + i\tau)$ by solving
- 9: ODE system $\partial_t m(c_h, z) + b_h(c_h, z; \hat{v}) = r(z) \quad \forall z \in W_h^{k_2}$
- 10: **end for**
- 11: **end while**

- Line 4: May also use cell-centered finite volume scheme
- Line 5: Uses RT_0 or BDM_1 finite element spaces
- Line 6: Based on CFL criterion
- Line 9: May use explicit or implicit method; here: SSP2/3

AMG For DG

Idea: Subspace correction in continuous Galerkin subspace [*B., Blatt, Scheichl, NLAA, 2012*]:

$$E_{AMG4DG} = (I - B_{SG}A_{DG})^{\nu_2}(I - R^T B_{AMG}RA_{DG})(I - B_{SG}A_{DG})^{\nu_1}$$

- B_{SG} : Preconditioner for DG system, e.g. Block SSOR
- R given by $\varphi_i^{CG} = \sum_{j=1}^{n_{DG}} r_{ij} \varphi_j^{DG}$
- Compute $A_{CG} = RA_{DG}R^T$ via sparse matrix multiplication
- B_{AMG} is one V-cycle of AMG preconditioner applied to A_{CG}
- We use our agglomeration-based AMG
- For heterogeneous problems CG may be replaced by P_0
- Purely algebraic, except the definition of R

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Exploiting Tensor Product Structure I

- Tensor product basis: $\phi_i(x) = \phi_{i_1, \dots, i_d}(x_1, \dots, x_d) = \theta_{i_d}(x_d) \dots \theta_{i_1}(x_1)$
- Tensor product quadrature: $\Xi_j = (\xi_{j_1}, \dots, \xi_{j_d})^T$
- $i = (i_1, \dots, i_d) \in I^d = \{1, \dots, m\}^d$, $j = (j_1, \dots, j_d) \in J^d = \{1, \dots, n\}^d$

Evaluate finite element function on ref elem at quadrature points:

$$\begin{aligned}
 u_h(\Xi_i) &= \sum_{j \in J^d} c_j \phi_j(\Xi_i) && (i \in I^d) \\
 &= \sum_{j_d \in J} \dots \sum_{j_1 \in J} \theta_{j_d}(\xi_{i_d}) \dots \theta_{j_1}(\xi_{i_1}) \quad c_{j_1, \dots, j_d} && \text{(tensor product structure)} \\
 &= \sum_{j_d \in J} \dots \sum_{j_1 \in J} A_{i_d, j_d}^{(d)} \dots A_{i_1, j_1}^{(1)} \quad c_{j_1, \dots, j_d} && (A^{(k)} \in \mathbb{R}^{m \times n})
 \end{aligned}$$

All basis functions at all quadrature points as matrix-vector product:

$$y = Ax = (A^{(d)} \otimes \dots \otimes A^{(1)}) x, \quad A_{(i_1, \dots, i_d), (j_1, \dots, j_d)} = \prod_{k=1}^d A_{i_k, j_k}^{(k)}$$

Sum Factorization I

Extract common factors:

$$\begin{aligned}
 Y_{i_1, \dots, i_d} &= \sum_{j_d \in J} \dots \sum_{j_1 \in J} A_{i_d, j_d}^{(d)} \dots A_{i_1, j_1}^{(1)} X_{j_1, \dots, j_d}^{(0)} \\
 &= \sum_{j_d \in J} \dots \sum_{j_2 \in J} A_{i_d, j_d}^{(d)} \dots A_{i_2, j_2}^{(2)} \underbrace{\left(\sum_{j_1 \in J} A_{i_1, j_1}^{(1)} X_{j_1, \dots, j_d}^{(0)} \right)}_{X_{j_2, \dots, j_d}^{(1)}} \\
 &= \sum_{j_d \in J} \dots \sum_{j_3 \in J} A_{i_d, j_d}^{(d)} \dots A_{i_3, j_3}^{(3)} \underbrace{\left(\sum_{j_2 \in J} A_{i_2, j_2}^{(2)} X_{j_2, \dots, j_d}^{(1)} \right)}_{X_{j_3, \dots, j_d}^{(2)}} \\
 &= \dots = \sum_{j_d \in J} A_{i_d, j_d}^{(d)} X_{j_d}^{(d-1)}
 \end{aligned}$$

Do this for all $(i_1, \dots, i_d)$ simultaneously

Sum Factorization II

- Algorithm by *Buis, Dyksen (1996)*
- For $k = 1, \dots, d$ do („dim by dim approach”)
 - ▶ compute **matrix-matrix product** $X^{(k)} = A^{(k)}X^{(k-1)}$
 - ▶ “rotate” $X^{(k)}$ to make j_{k+1} the row index (transposition in 2d)
- Overall complexity: $O(n^{d+1})$ instead of $O(n^{2d})$ ($n = m = k + 1$) for the steps above!
- Problem: matrices are rather small, e.g. $(4 \times 4) \cdot (4 \times 16)$ for $\mathbb{Q}_3$ in 3d
- Finite element application by e.g. *Melenk, Gerdes, Schwab (2001)*, *Kronbichler, Korman (2012)*:
 - ▶ compute $\hat{p}_h, \nabla \hat{p}_h$ at quadrature points, $O(n^{d+1})$
 - ▶ compute coefficients, geometry factors at quad. points, $O(n^d)$
 - ▶ compute residuals $O(n^{d+1})$
- Can be applied to nonlinear problems on curved elements
- Only tensor product structure of basis functions and quadrature is needed

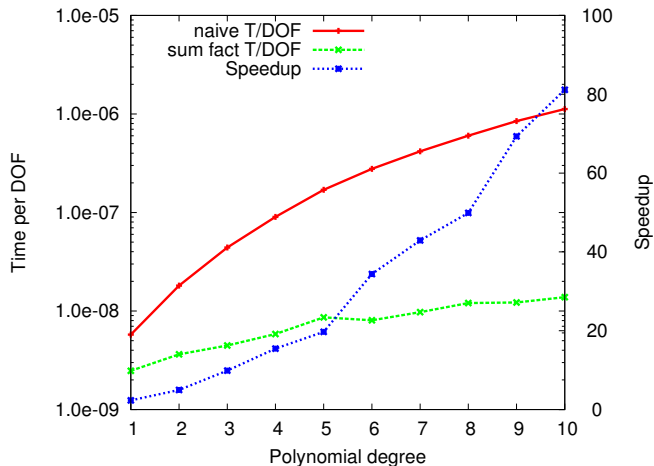
Sum Factorization III

- *Discontinuous Galerkin*: Effort for face terms dominates computation time for say $k < 10$
- Summary of complexities *per element* for $n = m = k + 1$ and $d = 3$

Operation	CG	DG (moderate k)
Matrix-free operator evaluation	n^4	n^3
Assembled operator evaluation	n^6	n^6
Sum factorized matrix assembly	n^7	n^6
Naive matrix assembly	n^9	n^9
Matrix-free point Jacobi application	n^4	n^3
Block SSOR preconditioner	n^9	n^9

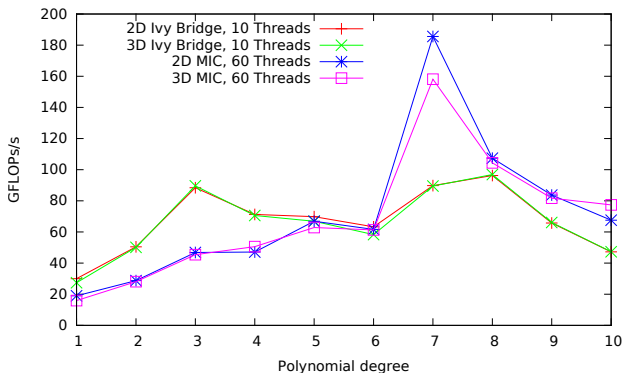
- Storage needed per element is reduced since only $1d$ basis functions need to be evaluated

Sum Factorization for Basis Evaluation in 3D



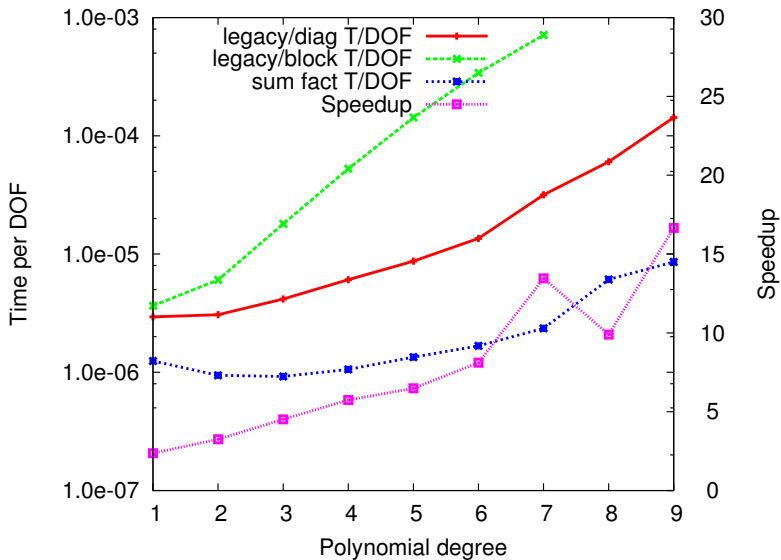
Except Q_1 in 2d, sum factorization is always faster than naive version

Floating-Point Performance: Auto vectorizer, threaded

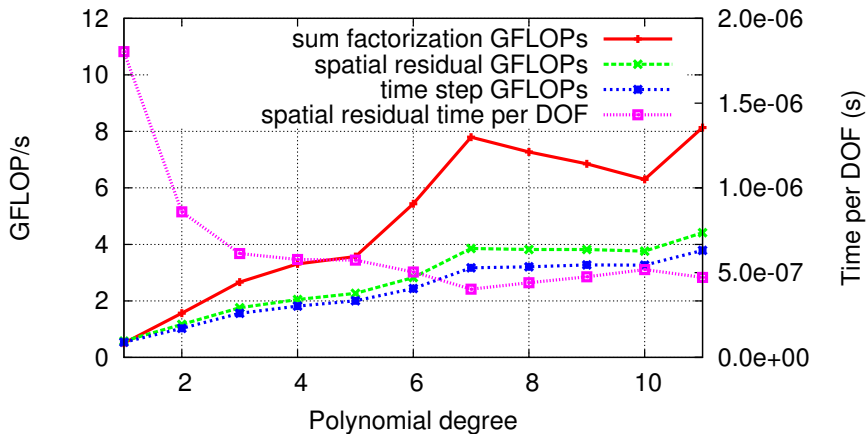


- Each thread runs sum factorization kernel in 3D
- icc with auto vectorization / SIMD hints
- Intel MIC (61 cores) vs. Xeon E5-2680v2 (Ivy Bridge, 10 cores)

Complete Explicit Time Step 3d



GFLOP Rates Per Core



Xeon E5-2680v2 10 core processor (22 GFLOPS peak / core)

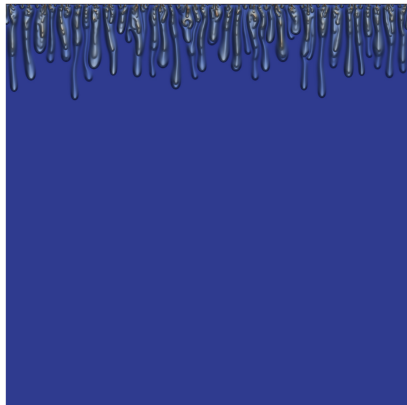
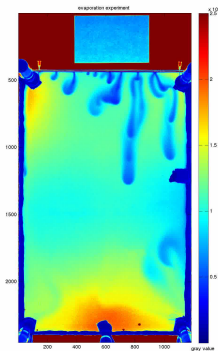
Residual: 4 GFLOPS

Sum factorization alone: 8 GFLOPS

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2d Density Driven Flow



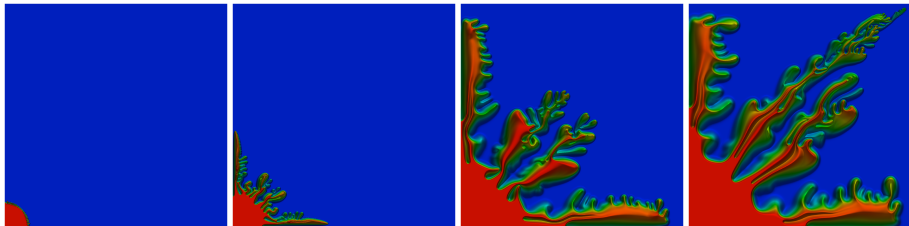
Experiment by Philip Kreyenberg
(IUP, Heidelberg)

$$Ra = 8000$$

25 cores of Xeon E5-2680v2
1200², Q_2 , $1.3 \cdot 10^7$ DOF

2d Miscible Displacement

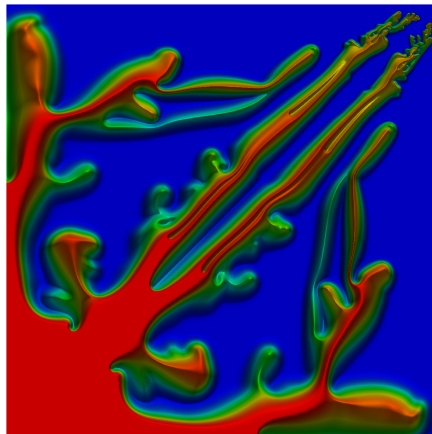
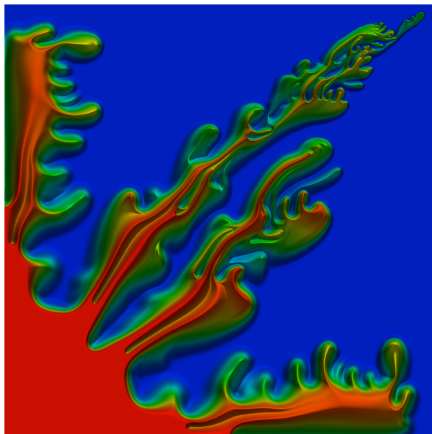
DG(Q_2)+BDM₁+DG(Q_3)+SSP3 — $9 \cdot 10^6$ DOF, 180000 time steps — 40 cores, 19 h



CCFV+RT₀+DG(Q_3)+SSP3 — $6.21 \cdot 10^6$ DOF, 180000 time steps — 40 cores, 19 h

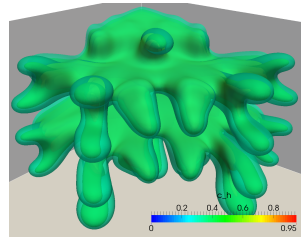
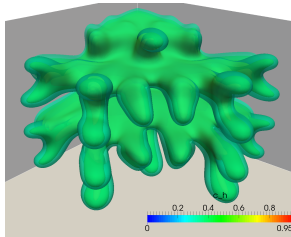
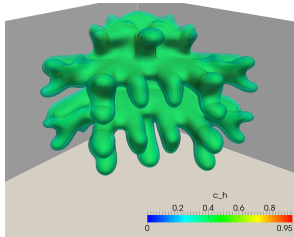
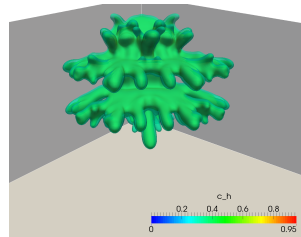
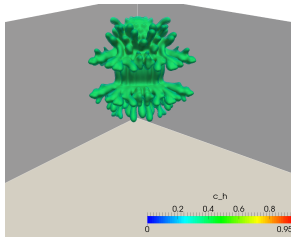
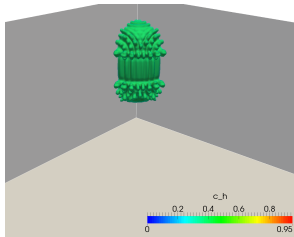


2d Miscible Displacement – Movies



3d Miscible Displacement

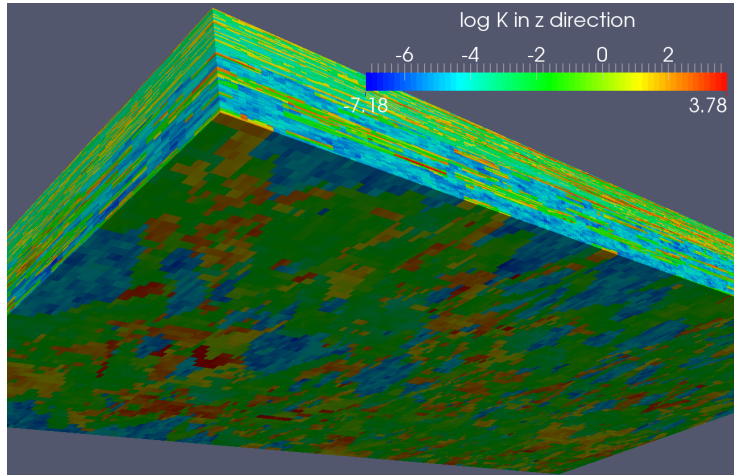
DG(Q_1)+BDM₁+DG(Q_2)+SSP2 — $259 \cdot 10^6$ DOF, 46000 time steps — 100 cores, 3 days



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SPE10 Test Problem I



Domain $1200 \times 2200 \times 170 \text{ ft}^3$, $60 \times 220 \times 85$ mesh (1.1 Mio cells)
Anisotropic, diagonal permeability tensor, 11 orders of magnitude variation

SPE10 Test Problem II

Solve elliptic problem $-\nabla \cdot (K_{SPE10} \nabla u) = 0$; $u = g$ on $\partial\Omega_{SPE10}$

DG(Q_1), hybrid DG/AMG- P_0 and Block-SSOR preconditioner

Discretization	DG- Q_1	DG- Q_1
Preconditioner	AMG-DG(2,2,V)	BSSOR(1)
DOF $\cdot 10^6$	9.0	9.0
#IT(10^{-6})	55	2637
TIT(s)	8.90	1.90
T _{solve} (s)	490	5010

→ Combine matrix-based and matrix-free

Summary

- All methods presented in this talk have been implemented in the DUNE software framework (www.dune-project.org)
- SF based DG implementation gives huge savings for matrix-free and matrix based approaches
- SF can handle variable coeffs, non-affine geometry, nonlinear problems
- Good SIMD performance is achieved for larger polynomial degrees
- Future work
 - ▶ Code generation or intrinsics might be necessary to achieve still better performance
 - ▶ Hybrid (matrix-free/matrix-based) preconditioners for implicit DG formulations
 - ▶ **Disclaimer:** Use moderate polynomial degree in practice!

Thank you for your attention!