

Algebras with short pp-definitions

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TCA - Pocinho

Short pp-definitions

Structures with short pp-definitions

$\mathbb{A} = (A; R_1, \dots, R_k)$... finite relational structure

$Q \subseteq A^n$ is **pp-definable** over \mathbb{A} if

$$Q(x_1, \dots, x_n) \Leftrightarrow \underbrace{\exists y_1, \dots, y_k R_{i_1}(\dots) \wedge \dots \wedge R_{i_j}(\dots)}_{\psi(x_1, \dots, x_n) \text{ pp-formula over } \mathbb{A}}$$

$\langle \mathbb{A} \rangle :=$ all pp-definable relations

Definition

- \mathbb{A} has **pp-definitions of length** $\leq f(n)$ if $\forall Q \in \langle \mathbb{A} \rangle \cap A^n$:
 Q is definable by a pp-formula ψ with $|\psi| \leq f(n)$
- \mathbb{A} has **short pp-definitions** if \mathbb{A} has pp-definitions of length $\leq p(n)$,
for a *polynomial* $p(n)$.

Question: Which \mathbb{A} have short pp-definitions?

Examples

Affine spaces

$$\mathbb{A} = (\{0, 1\}; \{(x, y, z) \mid x + y = z\}, \{0\}, \{1\}),$$

$$Q \in \langle \mathbb{A} \rangle \Leftrightarrow Q \text{ affine subspace of } \mathbb{Z}_2^n$$

$$\Leftrightarrow \text{given by } \leq n \text{ equations:}$$

$$x_{i_1} + x_{i_2} + \dots + x_{i_k} = a \Leftrightarrow$$

$$\exists y_2, \dots, y_k : (x_{i_1} + x_{i_2} = y_2) \wedge (y_2 + x_{i_3} = y_3) \wedge \dots \wedge (y_{k-1} + x_{i_k} = y_k) \wedge (y_k = a).$$

$$\Rightarrow \text{pp-definitions of length } O(n^2).$$

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\Rightarrow pp-definitions of length $O(n^2)$.

2-SAT

$$\mathbb{A} = (\{0, 1\}; (R_{a,b})_{a,b \in \{0,1\}}), \text{ with } R_{a,b} = \{0, 1\}^2 \setminus \{(a, b)\}.$$

$$Q \in \langle \mathbb{A} \rangle \Leftrightarrow Q(x_1, \dots, x_n) = \bigwedge_{1 \leq i, j \leq n} \text{pr}_{\{i,j\}} Q(x_i, x_j).$$

\Rightarrow pp-definitions of length $O(n^2)$.

Observation 1

\mathbb{A} has pp-defs. of length $\leq p(n)$

$\langle \mathbb{A} \rangle = \langle \mathbb{B} \rangle \Rightarrow \mathbb{B}$ has pp-defs. of length $\leq c \cdot p(n)$

Algebras/Clones with short pp-definitions

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$\text{Pol}(\mathbb{A}) = \{f: \mathbb{A}^n \rightarrow \mathbb{A} \mid n \in \mathbb{N}\} \dots$ *polymorphism clone of \mathbb{A}*

$\mathbf{A} \dots$ algebraic structure

$\text{Inv}(\mathbf{A}) = \{R \leq \mathbf{A}^n \mid n \in \mathbb{N}\}$ *invariant relations of \mathbf{A}*

$\text{Inv}(\text{Pol}(\mathbb{A})) = \langle \mathbb{A} \rangle \Rightarrow$ short pp-definitions is a **property of $\text{Pol}(\mathbb{A})$**
(even up to clone isomorphism).

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(even up to clone isomorphism).

Definition

\mathbf{A} has **short pp-definitions**, if $\text{Inv}(\mathbf{A}) = \langle \mathbb{A} \rangle$ has short pp-definitions.

Examples

- Affine subspaces of $\mathbb{Z}_2^n \leftrightarrow \mathbf{A} = (\{0, 1\}, x - y + z)$
- 2-SAT $\leftrightarrow \mathbf{A} = (\{0, 1\}, \text{maj}(x, y, z))$

Observation 2

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So **short pp-definitions** \Rightarrow **few subpowers**.

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So **short pp-definitions** \Rightarrow **few subpowers**.

If \mathbf{A} has few subpowers:

- \mathbf{A} has an edge term t (IMMVW'10):

$$t(y, y, x, x, x, \dots, x) \approx x$$

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\vdots

$$t(x, x, x, x, x, \dots, y) \approx x$$

- $\text{Inv}(\mathbf{A}) = \langle \mathbb{A} \rangle$ for some finite $\mathbb{A} = (A; R_1, \dots, R_n)$ (AMM'14)

A conjecture about few subpowers

Conjecture (Bulín)

- (weak) \mathbf{A} has short pp-defs. $\Leftrightarrow \mathbf{A}$ has few subpowers.
- (strong) \mathbf{A} has pp-defs. of length $O(n^k) \Leftrightarrow \mathbf{A}$ has a k -edge term.

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True for

- \mathbf{A} is affine
- \mathbf{A} has NU-term
 $y \approx t(y, x, \dots, x) \approx t(x, y, x, \dots, x) \approx \dots \approx t(x, \dots, x, y)$
- $|A| = 2$ (Lagerkvist, Wahlström '14)

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$|A| = 3$ not covered by above

Theorem (Bulín, MK '23)

If $\text{HSP}(\mathbf{A})$ is residually finite, then

\mathbf{A} has pp-definition of length $O(n^k) \Leftrightarrow \mathbf{A}$ has a k -edge term.

Main result

Theorem (Bulín, MK '23)

If $\text{HSP}(\mathbf{A})$ is residually finite, then

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\mathbf{B} is *subdirectly irreducible*, if $\text{Con}(\mathbf{B}) = \begin{array}{c} \circlearrowleft \\ \mu \\ \bullet \\ 0_{\mathbf{B}} \end{array}$

$\text{HSP}(\mathbf{A})$ *residually finite*, if

$\mathbf{B} \in \text{HSP}(\mathbf{A})$ is SI $\Leftrightarrow \mathbf{B} \in \{\mathbf{B}_1, \dots, \mathbf{B}_k\}$, $|B_i| < \infty$.

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(folklore) $|A| = 3$, \mathbf{A} few subpowers $\Rightarrow \text{HSP}(\mathbf{A})$ is residually finite.

Corollary (Bulín, MK '23)

If $|A| = 3$, then

\mathbf{A} has pp-definition of length $O(n^k) \Leftrightarrow \mathbf{A}$ has a k -edge term.

Proof idea

Proof step 1: Reduction to critical relations

A relation $R \leq \mathbf{A}^n$ is called **critical** if

- R is \wedge -irreducible ($R_1, R_2 > R \Rightarrow R_1 \cap R_2 > R$)
- R has no dummy variables

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Lemma

\mathbf{A} ... k -edge-term, $R \leq \mathbf{A}^n$. Then

$$R = \bigwedge_{|J| \leq k} (\text{pr}_J R) \wedge R_1 \wedge \dots \wedge R_l \text{ for } l \leq n \cdot |A|^2, R_i \text{ critical, } \textit{parallelogram property}.$$

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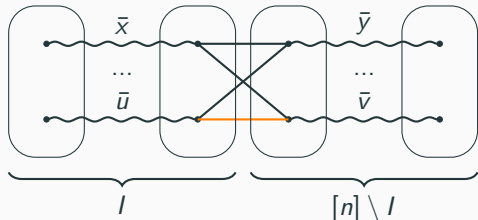
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$R \subseteq A^n$ has the **parallelogram property** if $\forall I \subset [n]$



$$(\bar{x}, \bar{y}), (\bar{x}, \bar{v}), (\bar{u}, \bar{y}) \in R \\ \Rightarrow (\bar{u}, \bar{v}) \in R$$

Proof step 2: Similarity

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Strategy: as for $x_1 + x_2 + \dots + x_n = a$

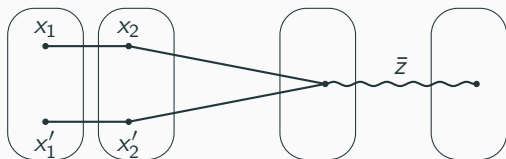
- $(x_1, x_2) \sim (x'_1, x'_2) :\Leftrightarrow \exists \bar{z} R(x_1, x_2, \bar{z}) \wedge R(x'_1, x'_2, \bar{z})$
- $\sim \in \text{Con}(\text{pr}_{1,2} R)$, $\mathbf{A}_{1,2} := (\text{pr}_{1,2} R) / \sim$

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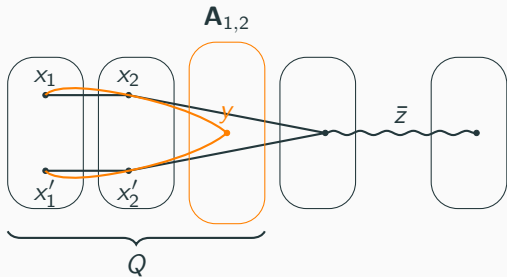


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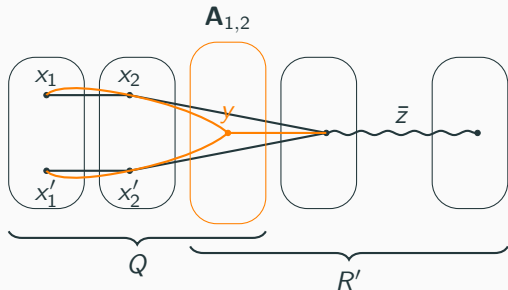
$$(x_1, x_2, y) \in Q :\Leftrightarrow \\ y = (x_1, x_2)/\sim$$

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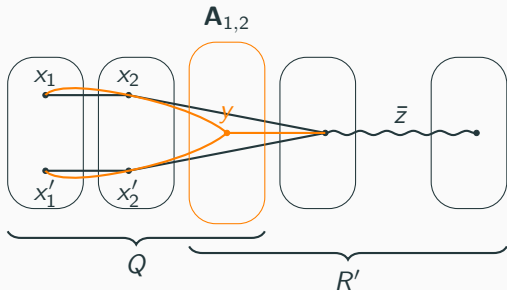
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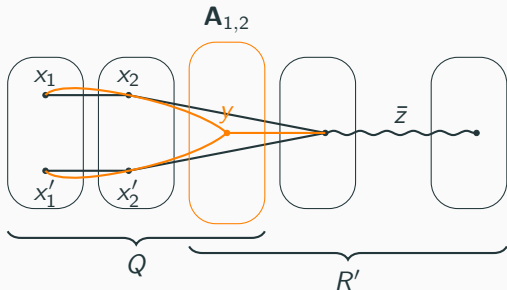
$$R(x_1, x_2, x_3, \dots, x_n) \Leftrightarrow \exists y \in A_{1,2} Q(x_1, x_2, y) \wedge R'(y, x_3, \dots, x_n).$$

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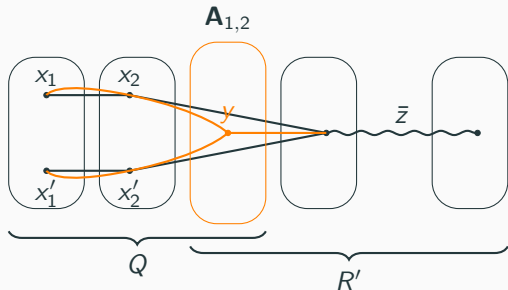
Problem: in general $\mathbf{A}_{1,2} \neq \mathbf{A}$

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Problem: in general $\mathbf{A}_{1,2} \neq \mathbf{A}$

But: R critical $\Rightarrow \mathbf{A}_{1,2}$ is SI \Rightarrow bounded by residual finiteness. \square

Application:
Subpower Membership Problem

Subpower Membership Problem

\mathbf{A} ... finite algebra

SMP(\mathbf{A})

INPUT: $\bar{a}_1, \dots, \bar{a}_k, \bar{b} \in A^n$

DECIDE: Is $\bar{b} \in \text{Sg}_{\mathbf{A}^n}(\bar{a}_1, \dots, \bar{a}_k)$?

Question (IMMVW'10): Is $\text{SMP}(\mathbf{A}) \in \text{P}$ for \mathbf{A} with few subpowers?

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Observation

$\bar{b} \notin \text{Sg}_{\mathbf{A}^n}(\bar{a}_1, \dots, \bar{a}_k) \Leftrightarrow \exists \text{pp-fma.} \psi : \neg \psi(\bar{b}) \wedge \psi(\bar{a}_1) \wedge \dots \wedge \psi(\bar{a}_k).$

A has short pp-definitions $\Rightarrow \text{SMP}(\mathbf{A}) \in \text{coNP}.$

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Theorem (BMS'19)

- $\text{SMP}(\mathbf{A}) \in \text{NP}$ if \mathbf{A} has few subpowers

(weak) Conjecture $\Rightarrow \text{SMP}(\mathbf{A}) \in \text{NP} \cap \text{coNP}.$

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Theorem (BMS'19)

- $\text{SMP}(\mathbf{A}) \in \text{NP}$ if \mathbf{A} has few subpowers
- $\text{SMP}(\mathbf{A}) \in \text{P}$ if further $\text{HSP}(\mathbf{A})$ is residually finite.

(weak) Conjecture $\Rightarrow \text{SMP}(\mathbf{A}) \in \text{NP} \cap \text{coNP}.$

Thank you for your attention!

Any questions?