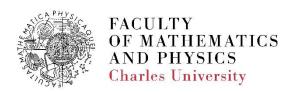
High-Performance Mixed Precision Numerical Linear Algebra

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Scientific Computing and Numerics (SCAN) Seminar
Cornell University



Exascale Computing: The Modern Space Race

- "Exascale": 10¹⁸ floating point operations per second
 - with maximum energy consumption around 20-40 MWatts
- Large investment in HPC worldwide



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Technical challenges at all levels

hardware to algorithms to applications

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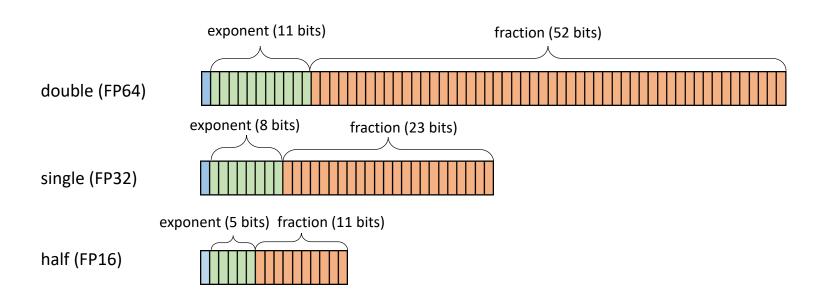
Hardware Support for Multiprecision Computation

Use of low precision in machine learning has driven emergence of low-precision capabilities in hardware:

- Half precision (FP16) defined as storage format in 2008 IEEE standard
- ARM NEON: SIMD architecture, instructions for 8x16-bit, 4x32-bit, 2x64-bit
- AMD Radeon Instinct MI25 GPU, 2017:
 - single: 12.3 TFLOPS, half: 24.6 TFLOPS
- NVIDIA Tesla P100, 2016: native ISA support for 16-bit FP arithmetic
- NVIDIA Tesla V100, 2017: tensor cores for half precision;
 - 4x4 matrix multiply in one clock cycle
 - double: 7 TFLOPS, half+tensor: 112 TFLOPS (16x!)
- NVIDIA A100, 2020: tensor cores with multiple supported precisions: FP16, FP64, Binary, INT4, INT8, bfloat16
- Google's Tensor processing unit (TPU): quantizes 32-bit FP computations into 8-bit integer arithmetic
- Future exascale supercomputers: (~2021) Expected extensive support for reduced-precision arithmetic (32/16/8-bit)

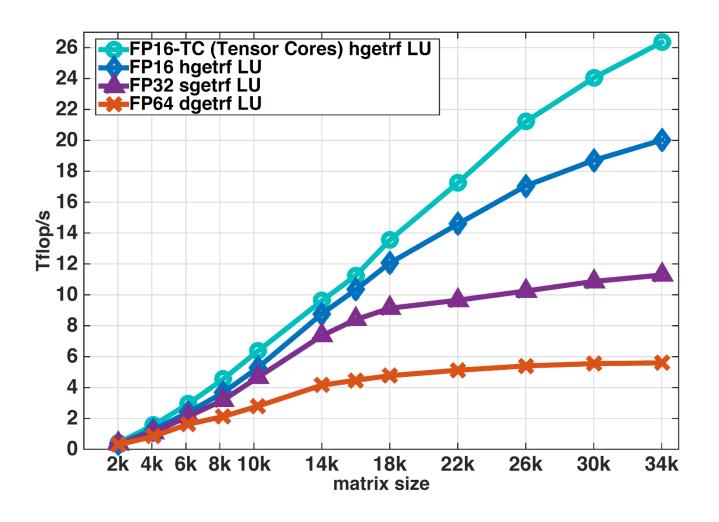
IEEE Floating Point Formats

$$(-1)^{\text{sign}} \times 2^{(\text{exponent-offset})} \times 1$$
. fraction



	u	x_{\min}^s	x_{\min}	x_{max}
fp16	4.88×10^{-4}	5.96×10^{-8}	6.10×10^{-5}	6.55×10^{4}
fp32	5.96×10^{-8}	1.40×10^{-45}	1.18×10^{-38}	3.40×10^{38}
fp64	1.11×10^{-16}	4.94×10^{-324}	2.22×10^{-308}	1.80×10^{308}

Performance of LU factorization on an NVIDIA V100 GPU



[Haidar, Tomov, Dongarra, Higham, 2018]

Mixed Precision Capabilities on Supercomputers

From TOP500:

June 2020

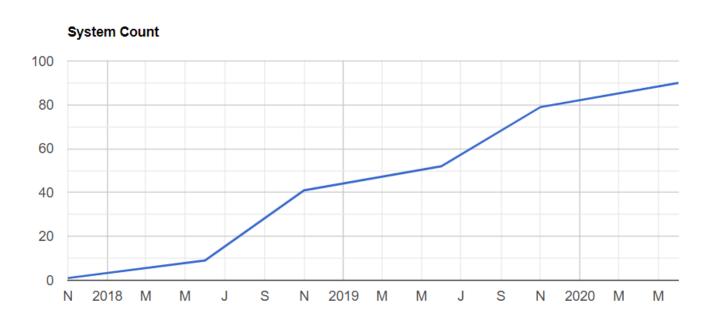
	Accelerator/Co-Processor	Count	System Share (%)	Rmax (GFlops)	Rpeak (GFlops)	Cores
1	NVIDIA Tesla V100	90	18	249,989,467	447,350,553	5,215,012
2	NVIDIA Tesla P100	17	3.4	63,536,390	99,682,523	1,350,772
3	NVIDIA Tesla V100 SXM2	13	2.6	63,358,310	100,091,657	1,218,784
4	NVIDIA Volta GV100	4	0.8	269,439,000	362,564,722	4,408,096
5	Nvidia Volta V100	3	0.6	24,246,000	33,871,440	451,136

June 2018

	Accelerator/Co-Processor	Count	System Share (%)	Rmax (GFlops)	Rpeak (GFlops)	Cores
1	NVIDIA Tesla P100	54	10.8	92,474,566	157,608,577	2,141,088
2	NVIDIA Tesla V100	9	1.8	10,769,400	21,164,239	242,016
3	NVIDIA Tesla K80	7	1.4	10,834,746	17,986,269	341,390
4	NVIDIA Tesla K40	6	1.2	11,062,790	18,034,192	246,920
5	NVIDIA Tesla P100 NVLink	5	1	11,716,000	18,611,235	211,268

Mixed Precision Capabilities on Supercomputers

ACCELERATOR/CO-PROCESSOR / NVIDIA TESLA V100



List	Count	System Share (%)	Rmax (GFlops)	Rpeak (GFlops)	Cores
Jun 2020	90	18	249,989,467	447,350,553	5,215,012
Nov 2019	79	15.8	171,425,790	316,852,935	3,676,828
Jun 2019	52	10.4	99,558,070	196,283,310	2,331,252
Nov 2018	41	8.2	65,671,290	136,467,348	1,645,796
Jun 2018	9	1.8	10,769,400	21,164,239	242,016
Nov 2017	1	0.2	1,070,000	1,819,752	22,440

- When will victory be declared?
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 - Solving dense Ax = b using Gaussian elimination with partial pivoting in double precision (FP64)

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- HPL benchmark is typically a compute-bound problem ("BLAS-3")
- Not a good indication of performance for a large number of applications!
 - Lots of remaining work even after exascale performance is achieved
 - Has led to incorporation of other benchmarks into the TOP500 ranking
 - e.g., HPCG: Solving sparse Ax = b iteratively using the conjugate gradient method

- HPL doesn't make use of modern mixed precision hardware
- We can already achieve "exaflop" performance today if we allow for mixed precision computations



https://www.olcf.ornl.gov/2018/06/08/genomics-code-exceeds-exaops-on-summit-supercomputer/

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=>HPL-AI: A new mixed precision benchmark

Iterative refinement: well-established method for improving an approximate solution to Ax = b

A is $n \times n$ and nonsingular; u is unit roundoff

Solve
$$Ax_0 = b$$
 by LU factorization for $i = 0$: maxit
$$r_i = b - Ax_i$$
 Solve $Ad_i = r_i$ via $d_i = U^{-1}(L^{-1}r_i)$ $x_{i+1} = x_i + d_i$

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Solve Ax_0 = b by LU factorization (in precision u) for i = 0: maxit  r_i = b - Ax_i \qquad \qquad \text{(in precision } u^2\text{)}  Solve Ad_i = r_i \qquad \text{via } d_i = U^{-1}(L^{-1}r_i) \qquad \text{(in precision } u\text{)}  x_{i+1} = x_i + d_i \qquad \qquad \text{(in precision } u\text{)}
```

"Traditional" (high-precision residual computation)

[Wilkinson, 1948] (fixed point), [Moler, 1967] (floating point)

$$\kappa_{\infty}(A) = \|A^{-1}\|_{\infty} \|A\|_{\infty}$$

As long as $\kappa_{\infty}(A) \leq u^{-1}$,

- relative forward error is O(u)
- relative normwise and componentwise backward errors are O(u)

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"Fixed-Precision"

[Jankowski and Woźniakowski, 1977], [Skeel, 1980], [Higham, 1991]

 $cond(A, x) = \| |A^{-1}||A||x| \|_{\infty} / \|x\|_{\infty}$

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"Fixed-Precision"

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Solve Ax_0 = b by LU factorization (in precision u^{1/2}) for i = 0: maxit  r_i = b - Ax_i \qquad \qquad \text{(in precision } u \text{)}  Solve Ad_i = r_i \qquad \text{via } d_i = U^{-1}(L^{-1}r_i) \qquad \text{(in precision } u \text{)}  x_{i+1} = x_i + d_i \qquad \qquad \text{(in precision } u \text{)}
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"Low-precision factorization"

[Langou et al., 2006], [Arioli and Duff, 2009], [Hogg and Scott, 2010], [Abdelfattah et al., 2016]

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⇒ 3-precision iterative refinement

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New analysis generalizes existing types of IR:

[C. and Higham, SIAM SISC 40(2), 2018]

Traditional	$u_f = u$, $u_r = u^2$
Fixed precision	$u_f = u = u_r$
Lower precision factorization	$u_f^2 = u = u_r$

(and improves upon existing analyses in some cases)

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(and improves upon existing analyses in some cases)

Enables new types of IR: (half, single, double), (half, single, quad),
 (half, double, quad), etc.

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Typical bounds used in analysis: $||A(x - \hat{x}_i)||_{\infty} \le ||A||_{\infty} ||x - \hat{x}_i||_{\infty}$

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For a stable refinement scheme, in early stages we expect

$$\frac{\|r_i\|}{\|A\|\|\hat{x}_i\|} \approx u \ll \frac{\|x - \hat{x}_i\|}{\|x\|} \longrightarrow \mu_i \ll 1$$

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But close to convergence,

$$||r_i|| \approx ||A|| ||x - \hat{x}_i|| \longrightarrow \mu_i \approx 1$$

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example: LU solve:

$$\frac{\mathbf{u}_{s}}{\|E_{i}\|_{\infty}} \leq 3n \frac{\mathbf{u}_{f}}{\|A^{-1}\|\widehat{L}\|\widehat{U}\|_{\infty}}$$

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- 2. $\|\hat{r}_i A\hat{d}_i\|_{\infty} \le u_s(c_1\|A\|_{\infty}\|\hat{d}_i\|_{\infty} + c_2\|\hat{r}_i\|_{\infty})$ \rightarrow normwise relative backward error is at most $\max(c_1, c_2) u_s$

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- 3. $\left|\hat{r}_i A\hat{d}_i\right| \leq \mathbf{u}_s G_i |\hat{d}_i|$
 - ightarrow componentwise relative backward error is bounded by a multiple of u_{s}

 $|Au_i| \leq u_s u_i |u_i|$

example: LU solve:

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$$\max(c_1, c_2) \mathbf{u}_s \leq \frac{3n\mathbf{u}_f \||\hat{L}||\hat{U}|\|_{\infty}}{\|A\|_{\infty}}$$

 E_i, c_1, c_2 , and G_i depend on A, \hat{r}_i , n, and u_s

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$$(u_s = u_f)$$

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Forward Error for IR3

- Three precisions:
 - u_f : factorization precision
 - *u*: working precision
 - u_r : residual computation precision

$$\kappa_{\infty}(A) = \|A^{-1}\|_{\infty} \|A\|_{\infty}$$

$$\operatorname{cond}(A) = \| |A^{-1}||A| \|_{\infty}$$

$$\operatorname{cond}(A, x) = \| |A^{-1}||A||x| \|_{\infty} / \|x\|_{\infty}$$

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Theorem [C. and Higham, SISC 40(2), 2018]

For IR in precisions $u_f \ge u \ge u_r$ and effective solve precision u_s , if

$$\phi_i \equiv 2\mathbf{u_s} \min(\operatorname{cond}(A), \kappa_{\infty}(A)\mu_i) + \mathbf{u_s} ||E_i||_{\infty}$$

is less than 1, then the forward error is reduced on the *i*th iteration by a factor $\approx \phi_i$ until an iterate \hat{x}_i is produced for which

$$\frac{\|x - \hat{x}_i\|_{\infty}}{\|x\|_{\infty}} \lesssim 4Nu_r \operatorname{cond}(A, x) + u,$$

where N is the maximum number of nonzeros per row in A.

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Analogous traditional bounds: $\phi_i \equiv 3n\mathbf{u_f}\kappa_{\infty}(A)$

Normwise Backward Error for IR3

Theorem [C. and Higham, SISC 40(2), 2018]

For IR in precisions $u_f \ge u \ge u_r$ and effective solve precision u_s , if

$$\phi_i \equiv (c_1 \kappa_{\infty}(A) + c_2) \mathbf{u}_s$$

is less than 1, then the residual is reduced on the *i*th iteration by a factor $\approx \phi_i$ until an iterate \hat{x}_i is produced for which

$$||b - A\hat{x}_i||_{\infty} \lesssim N\mathbf{u}(||b||_{\infty} + ||A||_{\infty}||\hat{x}_i||_{\infty}),$$

where N is the maximum number of nonzeros per row in A.

				Backwai	d error	
u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
Н	S	S	10 ⁴	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
S	S	S	10 ⁸	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
S	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
S	D	D	108	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
S	D	Q	108	10^{-16}	10^{-16}	10^{-16}

					Backwar	d error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10^{4}	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
	S	S	S	10 ⁸	10^{-8}	10^{-8}	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	S	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
LP fact.	S	D	D	10 ⁸	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	S	D	Q	10 ⁸	10^{-16}	10^{-16}	10^{-16}

					Backwai	d error	
	u_f	u	u_r	$\max \kappa_\infty(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10 ⁴	10-8	10-8	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
Fixed	S	S	S	10 ⁸	10^{-8}	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	S	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
LP fact.	S	D	D	10 ⁸	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	S	D	Q	10 ⁸	10^{-16}	10^{-16}	10^{-16}

					Backwai	rd error	
	u_f	u	u_r	$\max \kappa_\infty(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10^{4}	10^{-8}	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
Fixed	S	S	S	108	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
LP fact.	S	D	D	10 ⁸	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
	S	D	Q	108	10^{-16}	10^{-16}	10^{-16}

					Backwai	d error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10^{4}	10^{-8}	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
New	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
New	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
Fixed	S	S	S	10 ⁸	10^{-8}	10^{-8}	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Trad.	S	S	D	108	10^{-8}	10^{-8}	10^{-8}
LP fact.	S	D	D	108	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
New	S	D	Q	10 ⁸	10^{-16}	10^{-16}	10^{-16}

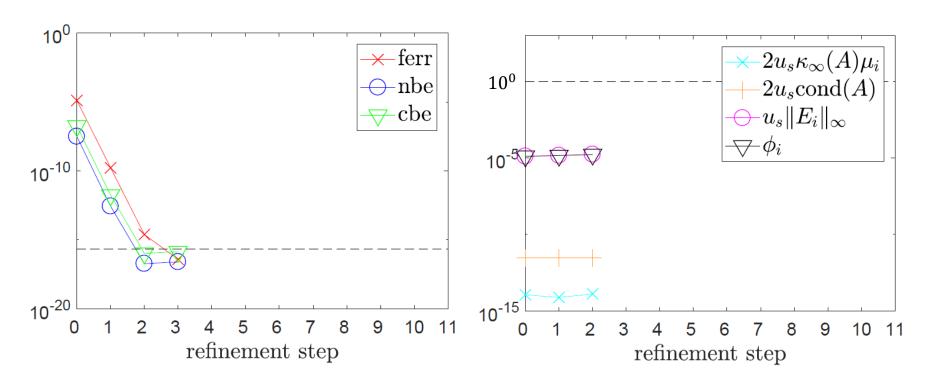
					Backwar	d error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10^{4}	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$
New	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$\operatorname{cond}(A, x) \cdot 10^{-16}$
New	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
New Fixed	H S	D S	Q S	10 ⁴ 10 ⁸	10 ⁻¹⁶ 10 ⁻⁸	10 ⁻¹⁶ 10 ⁻⁸	10^{-16} cond(A, x) · 10^{-8}
Į.							
Fixed	S	S	S	10 ⁸	10 ⁻⁸	10 ⁻⁸	$\operatorname{cond}(A, x) \cdot 10^{-8}$

 $[\]Rightarrow$ Benefit of IR3 vs. "LP fact.": no cond(A, x) term in forward error

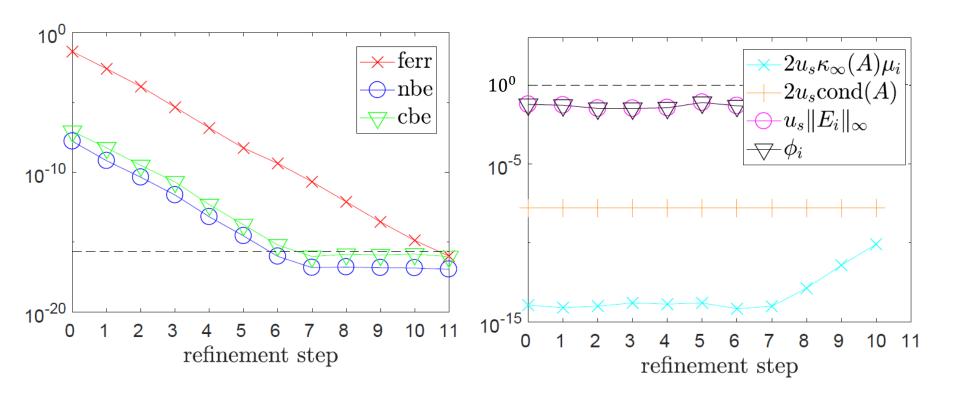
					Backwai	d error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10^{4}	10 ⁻⁸	10-8	$\operatorname{cond}(A, x) \cdot 10^{-8}$
New	Н	S	D	10^{4}	10^{-8}	10^{-8}	10^{-8}
LP fact.	Н	D	D	10^{4}	10^{-16}	10^{-16}	$cond(A, x) \cdot 10^{-16}$
New	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
Fixed	S	S	S	10 ⁸	10 ⁻⁸	10^{-8}	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 ⁸	10^{-8}	10 ⁻⁸	10^{-8}
LP fact.	S	D	D	10 ⁸	10^{-16}	10^{-16}	$cond(A, x) \cdot 10^{-16}$
New	S	D	Q	10 ⁸	10^{-16}	10^{-16}	10^{-16}

 $[\]Rightarrow$ Benefit of IR3 vs. traditional IR: As long as $\kappa_{\infty}(A) \leq 10^4$, can use lower precision factorization w/no loss of accuracy!

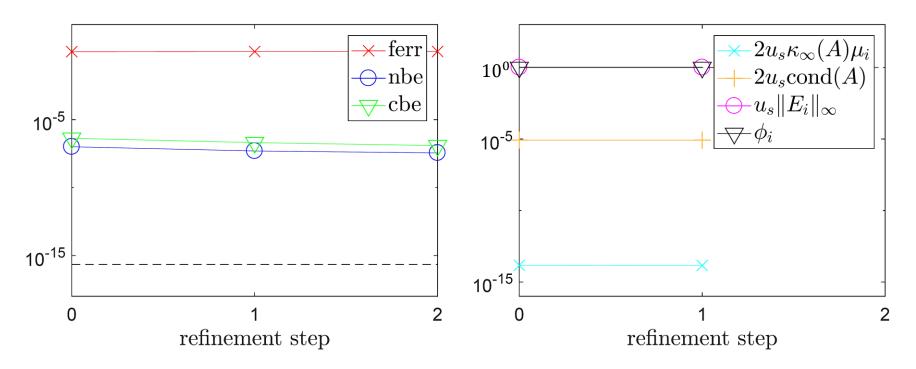
A = gallery('randsvd', 100, 1e3)
b = randn(100,1)
$$\kappa_{\infty}(A) \approx 1e4$$



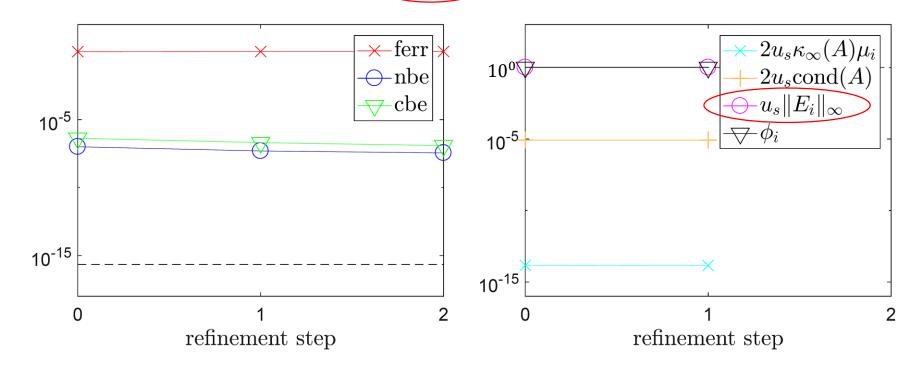
A = gallery('randsvd', 100, 1e7)
b = randn(100,1)
$$\kappa_{\infty}(A) \approx 7e7$$



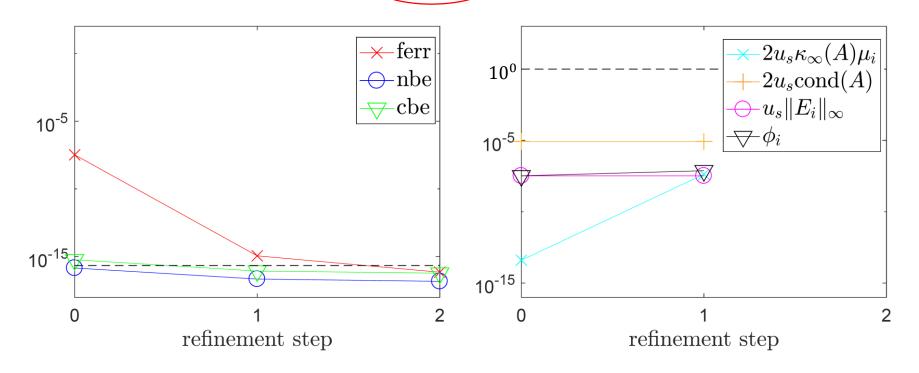
A = gallery('randsvd', 100, 1e9)
b = randn(100,1)
$$\kappa_{\infty}(A) \approx 2e10$$



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• Observation [Rump, 1990]: if \hat{L} and \hat{U} are computed LU factors of A in precision $\boldsymbol{u_f}$, then $\kappa_{\infty} \big(\widehat{U}^{-1} \widehat{L}^{-1} A \big) \approx 1 + \kappa_{\infty}(A) \boldsymbol{u_f},$

even if
$$\kappa_{\infty}(A) \gg u_f^{-1}$$
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GMRES-IR [C. and Higham, SISC 39(6), 2017]

• To compute the updates d_i , apply GMRES to $\widehat{U}^{-1}\widehat{L}^{-1}Ad_i=\widehat{U}^{-1}\widehat{L}^{-1}r_i$

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GMRES-IR [C. and Higham, SISC 39(6), 2017]

 $\hat{I} = \hat{I}^{-1} \hat{I}^{-1} A d_i = \hat{I}^{-1} \hat{I}^{-1} r_i$

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Solve $Ax_0 = b$ by LU factorization for i = 0: maxit $r_i = b - Ax_i$ Solve $Ad_i = r_i$ via GMRES on $\tilde{A}d_i = \tilde{r}_i$ $x_{i+1} = x_i + d_i$

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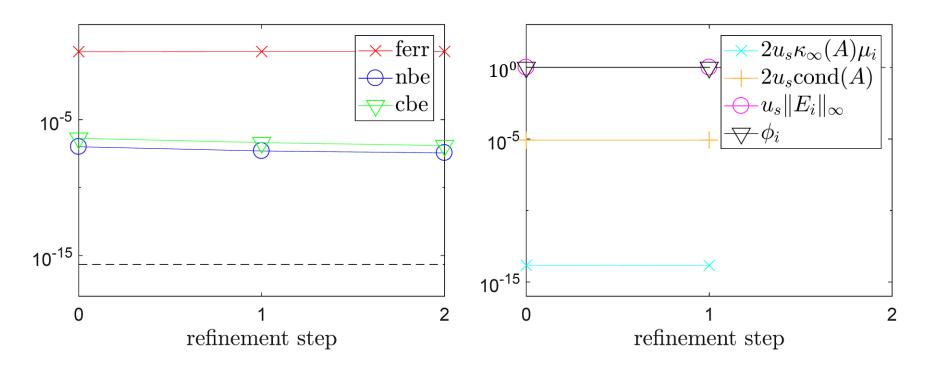
GMRES-IR [C. and Higham, SISC 39(6), 2017]

 $\tilde{Q}^{-1}\hat{L}^{-1}Ad_{i} = \tilde{Q}^{-1}\hat{L}^{-1}r_{i}$

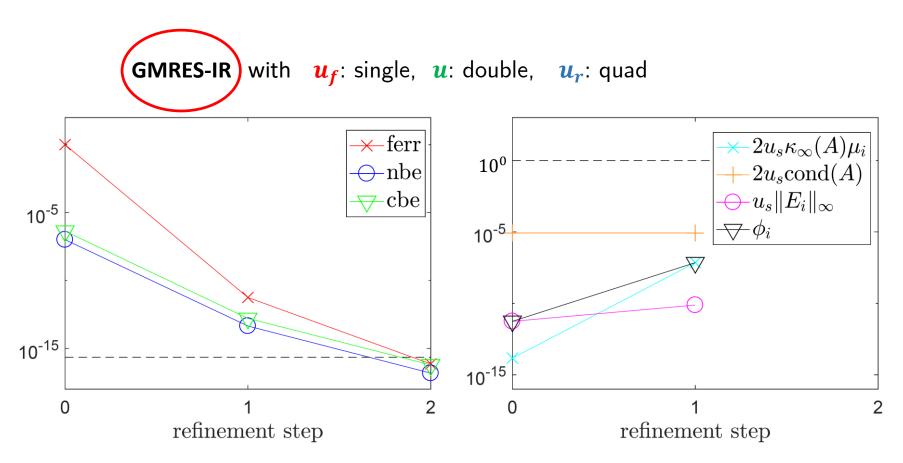
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A = gallery('randsvd', 100, 1e9, 2) b = randn(100,1) $\kappa_{\infty}(A) \approx 2e10$, $\operatorname{cond}(A, x) \approx 5e9$



A = gallery('randsvd', 100, 1e9, 2)
b = randn(100,1)
$$\kappa_{\infty}(A) \approx 2e10$$
, $\operatorname{cond}(A, x) \approx 5e9$, $\kappa_{\infty}(\tilde{A}) \approx 2e4$



Number of GMRES iterations: (2,3)

Benefits of GMRES-IR:

					Backwa	rd error	
	u_f	u	u_r	$\max \kappa_\infty(A)$	norm	comp	Forward error
LU-IR	Н	S	D	10 ⁴	10 ⁻⁸	10 ⁻⁸	10 ⁻⁸
GMRES-IR	Н	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
LU-IR	S	D	Q	10 ⁸	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	S	D	Q	10^{16}	10^{-16}	10^{-16}	10^{-16}
LU-IR	Н	D	Q	10^{4}	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	Н	D	Q	10^{12}	10^{-16}	10^{-16}	10^{-16}

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GMRES-IR	S	D	Q	10^{16}	10^{-16}	10^{-16}	10^{-16}
LU-IR	Н	D	Q	10 ⁴	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	Н	D	Q	10^{12}	10^{-16}	10^{-16}	10^{-16}

 \Rightarrow With GMRES-IR, low precision factorization will work for higher $\kappa_{\infty}(A)$

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LU-IR	Н	D	Q	10 ⁴	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	Н	D	Q	10^{12}	10^{-16}	10^{-16}	10^{-16}

 \Rightarrow With GMRES-IR, lower precision factorization will work for higher $\kappa_{\infty}(A)$ $\kappa_{\infty}(A) \leq u^{-1/2} u_f^{-1}$

Benefits of GMRES-IR:

					Backwa	rd error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LU-IR	Н	S	D	10 ⁴	10 ⁻⁸	10 ⁻⁸	10 ⁻⁸
GMRES-IR	Н	S	D	10 ⁸	10^{-8}	10^{-8}	10^{-8}
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GMRES-IR	S	D	Q	10^{16}	10^{-16}	10^{-16}	10^{-16}
LU-IR	Н	D	Q	10 ⁴	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	H	D	Q	10^{12}	10^{-16}	10^{-16}	10^{-16}

 $[\]Rightarrow$ As long as $\kappa_{\infty}(A) \leq 10^{12}$, can use half precision factorization and still obtain double precision accuracy!

Benefits of GMRES-IR:

					Backwa	rd error	
	u_f	u	u_r	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LU-IR	Н	S	D	10 ⁴	10 ⁻⁸	10 ⁻⁸	10 ⁻⁸
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LU-IR	Н	D	Q	10 ⁴	10^{-16}	10^{-16}	10^{-16}
GMRES-IR	H	D	Q	10^{12}	10^{-16}	10^{-16}	10^{-16}

Try IR3! MATLAB codes available at:

https://github.com/eccarson/ir3

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 - Smaller $\tau \to \text{more GMRES}$ iterations, potentially fewer refinement steps
 - Larger $\tau \rightarrow$ fewer GMRES iterations, potentially more refinement steps

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- What about overflow, underflow, subnormal numbers?
 - Sophisticated scaling methods can help avoid this
 - "Squeezing a Matrix into Half Precision, with an Application to Solving Linear Systems" [Higham, Pranesh, Zounon, 2019]

• Convergence rate of GMRES?

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 - If A is ill conditioned and LU factorization is performed in very low precision, it can be a poor preconditioner
 - e.g., if \tilde{A} still has cluster of eigenvalues near origin, GMRES can stagnate until n^{th} iteration, regardless of $\kappa_{\infty}(A)$ [Liesen and Tichý, 2004]
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 - Potential remedies: deflation, Krylov subspace recycling, using additional preconditioner
- Depending on conditioning of A, applying \tilde{A} to a vector must be done accurately (precision u^2) in each GMRES iteration
- Why GMRES?
 - Theoretical purposes: existing analysis and proof of backward stability [Paige, Rozložník, Strakoš, 2006]
 - In practice, use any solver you want!

GMRES-IR in Libraries and Applications

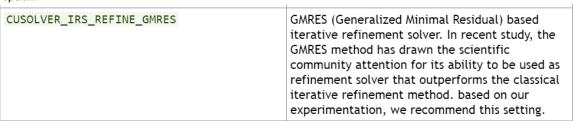
MAGMA: Dense linear algebra routines for heterogeneous/hybrid architectures

```
magma / src / dxgesv gmres gpu.cpp
128
129
          DSGESV or DHGESV expert interface.
130
          It computes the solution to a real system of linear equations
             A * X = B, A^{**T} * X = B, or A^{**H} * X = B,
131
          where A is an N-by-N matrix and X and B are N-by-NRHS matrices.
132
133
          the accomodate the Single Precision DSGESV and the Half precision dhgesv API.
134
          precision and iterative refinement solver are specified by facto type, solver type.
135
          For other API parameter please refer to the corresponding dsgesv or dhgesv.
```

NVIDIA's CuSolver Library

2.2.1.6. cusolverIRSRefinement_t

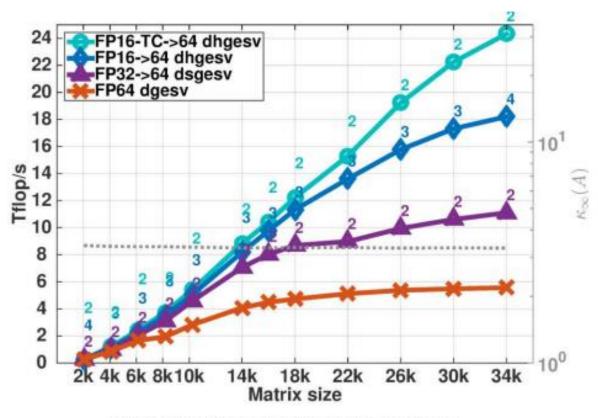
The cusolverIRSRefinement_t type indicates which solver type would be used for the specific cusolver function. Most of our experimentation shows that CUSOLVER_IRS_REFINE_GMRES is the best option.



 In production codes: FK6D/ASGarD code (Oak Ridge National Lab, USA) for tokomak containment problem

Performance Results (MAGMA)

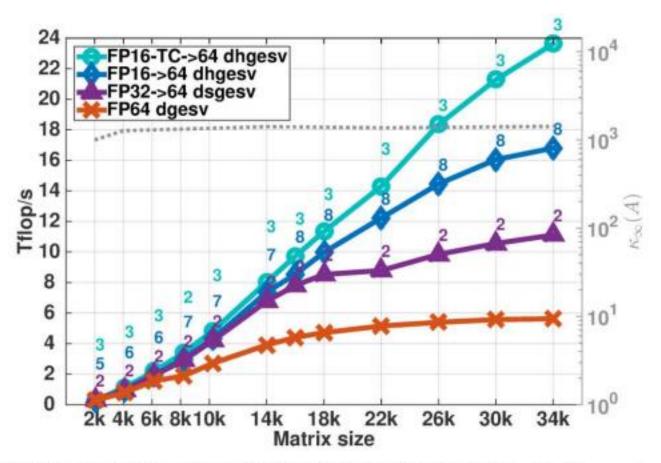
- [Haidar, Tomov, Dongarra, Higham, 2018]
- 2-precision GMRES-IR approach ($u = u_r$) on NVIDIA V100
- IR run to FP64 accuracy, max 400 iterations in GMRES
- Tflops/s measured as $(2n^3/3)$ /time



(a) Matrix of type 1: diagonally dominant.

Performance Results (MAGMA)

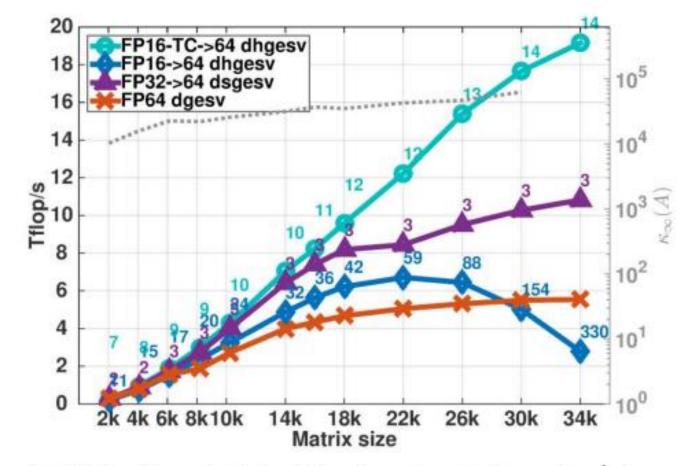
[Haidar, Tomov, Dongarra, Higham, 2018]



(a) Matrix of type 3: positive λ with clustered singular values, $\sigma_i = (1, \dots, 1, \frac{1}{cond})$.

Performance Results (MAGMA)

[Haidar, Tomov, Dongarra, Higham, 2018]



(b) Matrix of type 4: clustered singular values, $\sigma_i = (1, \dots, 1, \frac{1}{cond})$.

Performance Results

• [Haidar, Tomov, Dongarra, Higham, 2018]

Performance for Matrices from SuiteSparse

name	Description	size	$\kappa_{\infty}(A)$	dgesv	dsgesv		dhgesv		dhgesv-TC	
				time(s)	# iter	time (s)	# iter	time (s)	# iter	time (s)
em192	radar design	26896	106	5.70	3	3.11	40	5.21	10	2.05
appu	NASA app benchmark	14000	104	0.43	2	0.27	7	0.24	4	0.19
ns3Da	3D Navier Stokes	20414	$7.6 \ 10^3$	1.12	2	0.69	6	0.54	4	0.43
nd6k	ND problem set	18000	$3.5 \ 10^2$	0.81	2	0.45	4	0.36	3	0.30
nd12k	ND problem set	36000	$4.3 \ 10^2$	5.36	2	2.75	3	1.76	3	1.31

HPL-Al Benchmark

- HPL/LINPACK benchmark has been used in TOP500 since the 90s
 - Double precision, dense Ax=b using GEPP
 - Not necessarily indicative of application performance, especially for ML/Al applications
 - Doesn't take advantage of low-precision hardware!

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- HPL-AI benchmark (2019)
 - Highlights confluence of HPC+AI workloads
 - Like HPL, solves dense Ax=b, results still to double precision accuracy
 - Achieves this via mixed-precision GMRES-IR
 - may be implemented in a way that takes advantage of the current and upcoming devices for accelerating AI workloads

HPL-Al Benchmark Performance

• So far, two machines have reported results for HPL-AI (#1 and #2 on TOP500)

HPL-Al Results (June 2020):

- Summit: 550 PETAFLOP/s (vs. 148.6 PETAFLOP/s on HPL)
- Fugaku: 1.42 EXAFLOP/s (vs. 415.5 PETAFLOP/s on HPL)





HPL-Al Benchmark

• In the future, HPL-AI will gain same status as benchmarks that complement HPL, like HPCG, Graph500, Green500

More information: https://icl.bitbucket.io/hpl-ai/

Reference implementation: https://bitbucket.org/icl/hpl-ai/src/

Extension to Least Squares Problems

Want to solve

$$\min_{x} \|b - Ax\|_2$$

where $A \in \mathbb{R}^{m \times n}$ (m > n) has rank n

Commonly solved using QR factorization:

$$A = QR = [Q_1, Q_2] \begin{bmatrix} U \\ 0 \end{bmatrix}$$

where Q is an $m \times m$ orthogonal matrix and U is upper triangular.

$$x = U^{-1}Q_1^T b$$
, $||b - Ax||_2 = ||Q_2^T b||_2$

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• As in linear system case, for ill-conditioned problems, iterative refinement often needed to improve accuracy and stability

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck,1967): Least squares problem can be written as a linear system with square matrix of size (m + n):

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$$

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- Refinement proceeds as follows:
- 1. Compute "residuals"

$$\begin{bmatrix} f_i \\ g_i \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} - \begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r_i \\ x_i \end{bmatrix} = \begin{bmatrix} b - r_i - Ax_i \\ -A^T r_i \end{bmatrix}$$

2. Solve for corrections

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix} = \begin{bmatrix} f_i \\ g_i \end{bmatrix}$$

3. Update "solution":

$$\begin{bmatrix} r_{i+1} \\ x_{i+1} \end{bmatrix} = \begin{bmatrix} r_i \\ x_i \end{bmatrix} + \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix}$$

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck,1967): Least squares problem can be written as a linear system with square matrix of size (m + n):

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} \qquad \tilde{A}\tilde{x} = \tilde{b}$$

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- 1. Compute "residuals"

$$\begin{bmatrix} f_i \\ g_i \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} - \begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r_i \\ x_i \end{bmatrix} = \begin{bmatrix} b - r_i - Ax_i \\ -A^T r_i \end{bmatrix}$$

2. Solve for corrections

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix} = \begin{bmatrix} f_i \\ g_i \end{bmatrix}$$

Update "solution":

$$\begin{bmatrix} r_{i+1} \\ x_{i+1} \end{bmatrix} = \begin{bmatrix} r_i \\ x_i \end{bmatrix} + \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix}$$

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck, 1967): Least squares problem can be written as a linear system with square matrix of size (m + n):

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Results for 3-precision IR for linear systems also applies to least squares problems

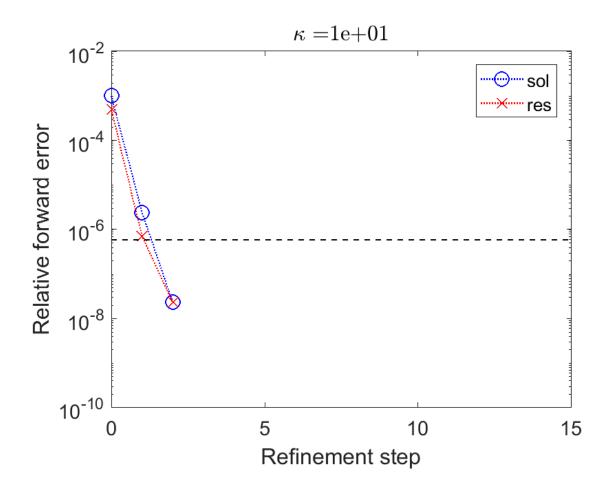
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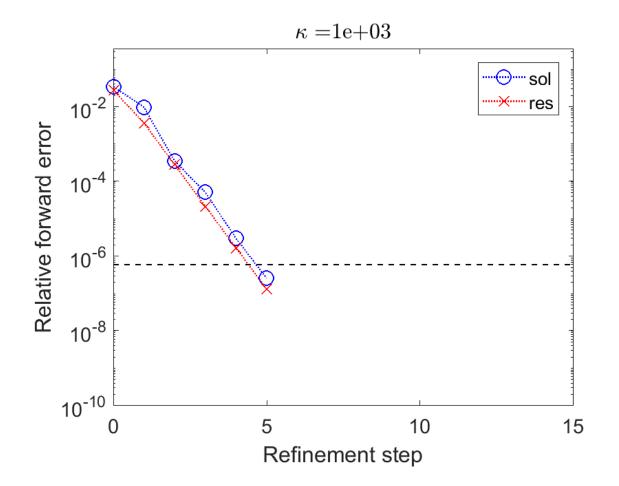
A = gallery('randsvd', [100, 10], kappa,3)
b = randn(100,1); b = b./norm(b)

Standard (QR-based) least squares IR with u_f : half, u: single, u_r : double



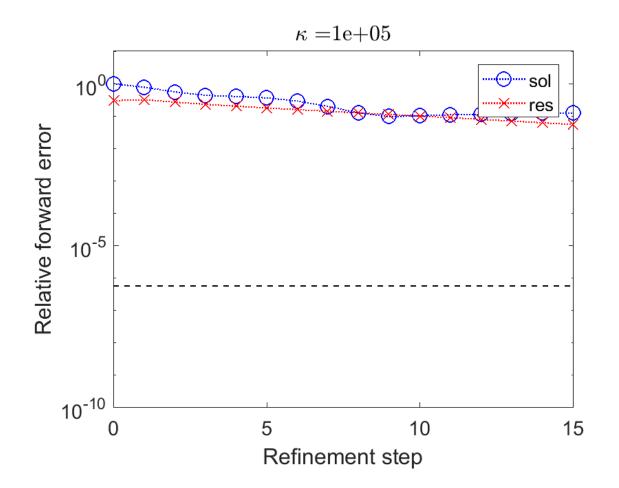
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- Similar to the linear system case, we can use a lower precision factorization for even more ill-conditioned problems if we **improve the effective precision of the solver**
- Again, don't want to compute an LU factorization of the augmented system
- How can we use existing QR factors to construct an effective and inexpensive preconditioner for the augmented system?
- Note that augmented system is a saddle-point system; lots of existing work (block diagonal, triangular, constraint-based, ...)

Ex: block diagonal preconditioner ([Murphy, Golub, Wathen, 2000], [Ipsen, 2001])

$$\begin{bmatrix} \alpha I & 0 \\ 0 & \frac{1}{\alpha} \hat{R}^T \hat{R} \end{bmatrix} = \begin{bmatrix} \sqrt{\alpha} I & 0 \\ 0 & \frac{1}{\sqrt{\alpha}} \hat{R}^T \end{bmatrix} \begin{bmatrix} \sqrt{\alpha} I & 0 \\ 0 & \frac{1}{\sqrt{\alpha}} \hat{R} \end{bmatrix} \equiv M_1 M_2$$

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Assuming QR factorization is exact,

$$M_2^{-1} M_1^{-1} \tilde{A} = \begin{vmatrix} I & \frac{1}{\alpha} A \\ \alpha \hat{R}^{-1} \hat{R}^{-T} A^T & 0 \end{vmatrix}$$

is nonsymmetric, diagonalizable, with eigenvalues $\{1, \frac{1}{2}(1 \pm \sqrt{5})\}$.

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- However, condition number can still be quite large; unsuitable for proving backward stability of GMRES
- If we take split preconditioner

$$M_1^{-1}\tilde{A}M_2^{-1} = \begin{bmatrix} I & A\hat{R} \\ \hat{R}^{-T}A^T & 0 \end{bmatrix}$$

we will have a well-conditioned system

- However, split-preconditioned GMRES is not backward stable
- Potentially useful in practice, not but in theory

• One option:

$$M = \begin{bmatrix} \alpha I & \hat{Q}_1 \hat{R} \\ \hat{R}^T \hat{Q}_1^T & 0 \end{bmatrix}$$

Then we can prove that for the left-preconditioned system,

$$\kappa(M^{-1}\tilde{A}) \le (1 + \mathbf{u_f} c \kappa(A))^2$$

where $c = O(m^2)$, where we note this bound is pessimistic.

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• So for GMRES-based LSIR, $u_s \equiv u$; expect convergence of forward error when $\kappa_{\infty}(A) < u^{-1/2} u_f^{-1}$ [C., Higham, Pranesh, SISC 2020]

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- "A Survey of Numerical Methods Utilizing Mixed Precision Arithmetic" https://arxiv.org/abs/2007.06674

Thank You!

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