# Mixed Precision Iterative Refinement 

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FACULTY
of mathematics
AND PHYSICS
Charles University

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## Floating Point Formats

$$
(-1)^{\text {sign }} \times 2^{(\text {exponent-offset) }} \times 1 \text {. fraction }
$$

IEEE double (FP64)


| IEEE half (FP16) | exponent ( 5 bits) fraction (10 bits) |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | $\begin{gathered} \text { size } \\ \text { (bits) } \end{gathered}$ | range | $u$ | perf. (NVIDIA H100) |
|  | exponent ( 8 bits) fraction ( 7 bits) | FP64 | 64 | $10^{ \pm 308}$ | $1 \times 10^{-16}$ | 60 Tflops/s |
| bfloat16 |  | FP32 | 32 | $10^{ \pm 38}$ | $6 \times 10^{-8}$ | 1 Pflop/s |
|  |  | FP16 | 16 | $10^{ \pm 5}$ | $5 \times 10^{-4}$ | 2 Pflops/s |
| FP8 $\left\{\begin{array}{l}\left\\|\left\\|\left\\|\\|^{-1}\right.\right.\right. \\ \text { e5m2 }\end{array}\right.$ |  | bfloat16 | 16 | $10^{ \pm 38}$ | $4 \times 10^{-3}$ |  |
|  |  | FP8-e5m2 | 8 | $10^{ \pm 5}$ | $3 \times 10^{-1}$ | 4 Pflops/s |
|  |  | FP8-e4m3 | 8 | $10^{ \pm 2}$ | $1 \times 10^{-1}$ |  |

## Hardware Support for Multiprecision Computation

Use of low precision in machine learning has driven emergence of lowprecision capabilities in hardware:

- Half precision (FP16) defined as storage format in 2008 IEEE standard
- ARM NEON: SIMD architecture, instructions for $8 \times 16$-bit, $4 \times 32$-bit, $2 \times 64$-bit
- AMD Radeon Instinct MI25 GPU, 2017:
- single: 12.3 TFLOPS, half: 24.6 TFLOPS
- NVIDIA Tesla P100, 2016: native ISA support for 16-bit FP arithmetic
- NVIDIA Tesla V100, 2017: tensor cores for half precision;
$4 \times 4$ matrix multiply in one clock cycle
- double: 7 TFLOPS, half+tensor: 112 TFLOPS (16x!)
- Google's Tensor processing unit (TPU)
- NVIDIA A100, 2020: tensor cores with multiple supported precisions: FP16, FP64, Binary, INT4, INT8, bfloat16
- NVIDIA H100, 2022: now with quarter-precision (FP8) tensor cores
- Exascale supercomputers: Expected extensive support for reduced-precision arithmetic (Frontier: FP64, FP32, FP16, bfloat16, INT8, INT4)


## Mixed precision in NLA

- BLAS: cuBLAS, MAGMA, [Agullo et al. 2009], [Abdelfattah et al., 2019], [Haidar et al., 2018]
- Iterative refinement:
- Long history: [Wilkinson, 1963], [Moler, 1967], [Stewart, 1973], .
- More recently: [Langou et al., 2006], [C., Higham, 2017], [C., Higham, 2018], [C., Higham, Pranesh, 2020], [Amestoy et al., 2021]
- Matrix factorizations: [Haidar et al., 2017], [Haidar et al., 2018], [Haidar et al., 2020], [Abdelfattah et al., 2020]
- Eigenvalue problems: [Dongarra, 1982], [Dongarra, 1983], [Tisseur, 2001], [Davies et al., 2001], [Petschow et al., 2014], [Alvermann et al., 2019]
- Sparse direct solvers: [Buttari et al., 2008]
- Orthogonalization: [Yamazaki et al., 2015]
- Multigrid: [Tamstorf et al., 2020], [Richter et al., 2014], [Sumiyoshi et al., 2014], [Ljungkvist, Kronbichler, 2017, 2019]
- (Preconditioned) Krylov subspace methods: [Emans, van der Meer, 2012], [Yamagishi, Matsumura, 2016], [C., Gergelits, Yamazaki, 2021], [Clark, 2019], [Anzt et al., 2019], [Clark et al., 2010], [Gratton et al., 2020], [Arioli, Duff, 2009], [Hogg, Scott, 2010]

For survey and references, see [Abdelfattah et al., IJHPC, 2021]

## HPL-Al Benchmark

- Supercomputers traditionally ranked by performance on high-performance LINPACK (HPL) benchmark
- Solves dense $A x=b$ via Gaussian elimination with partial pivoting
- HPL-AI: Like HPL, solves dense $A x=b$, results still to double precision accuracy
- But achieves this via mixed-precision GMRES-based iterative refinement


## HPL-Al Benchmark

## June 2022

| Rank | Site | Computer | Cores | HPL-AI (Eflop/s) | TOP500 Rank | HPL Rmax (Eflop/s) | Speedup |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 1 | DOE/SC/ORNL, USA | Frontier | $8,730,112$ | 6.861 | 1 | 1.102 | 6.2 |
| 2 | RIKEN, Japan | Fugaku | $7,630,848$ | 2.000 | 2 | 0.4420 | 4.5 |
| 3 | DOE/SC/ORNL, USA | Summit | $2,414,592$ | 1.411 | 4 | 0.1486 | 9.5 |
| 4 | NVIDIA, USA | Selene | 555,520 | 0.630 | 8 | 0.0630 | 9.9 |
| 5 | DOE/SC/LBNL, USA | Perlmutter | 761,856 | 0.590 | 7 | 0.0709 | 8.3 |
| 6 | FZJ, Germany | JUWELS BM | 449,280 | 0.470 | 11 | 0.0440 | 10.0 |
| 7 | University of Florida, USA | HiPerGator | 138,880 | 0.170 | 34 | 0.0170 | 9.9 |
| 8 | SberCloud, Russia | Christofari Neo | 98,208 | 0.123 | 47 | 0.0120 | 10.3 |
| 9 | DOE/SC/ANL, USA | Polaris | 259,840 | 0.114 | 14 | 0.0238 | 4.8 |
| 10 | ITC, Japan | Wisteria | 368,640 | 0.100 | 20 | 0.0220 | 4.5 |

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## Iterative Refinement for $A x=b$

Iterative refinement: well-established method for improving an approximate solution to $A x=b$
$A$ is $n \times n$ and nonsingular; $u$ is unit roundoff
Solve $A x_{0}=b$ by LU factorization
for $i=0$ : maxit

$$
r_{i}=b-A x_{i}
$$

Solve $A d_{i}=r_{i} \quad$ via $d_{i}=U^{-1}\left(L^{-1} r_{i}\right)$
$x_{i+1}=x_{i}+d_{i}$

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$A$ is $n \times n$ and nonsingular; $u$ is unit roundoff
Solve $A x_{0}=b$ by LU factorization
(in precision $u$ )
for $i=0$ : maxit

$$
\begin{array}{ll}
r_{i}=b-A x_{i} & \\
\text { Solve } A d_{i}=r_{i} & \text { (in precision } \left.u^{2}\right) \\
x_{i+1}=x_{i}+d_{i} & \\
\text { via } d_{i}=U^{-1}\left(L^{-1} r_{i}\right) & \text { (in precision } u \text { ) } \\
\text { (in precision } u \text { ) }
\end{array}
$$

$$
\begin{array}{ll}
\text { "Traditional" } & \begin{array}{l}
\text { (high-precision } \\
\text { residual computation) }
\end{array}
\end{array}
$$

[Wilkinson, 1948] (fixed point), [Moler, 1967] (floating point)

## Iterative Refinement for $A x=b$

$$
\kappa_{\infty}(A)=\left\|A^{-1}\right\|_{\infty}\|A\|_{\infty}
$$

As long as $\kappa_{\infty}(A) \leq u^{-1}$,

- relative forward error is $O(u)$
- relative normwise and componentwise backward errors are $O(u)$

Solve $A x_{0}=b$ by LU factorization
(in precision $u$ )
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$$
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\text { (in precision } u^{2} \text { ) } \\
\text { (in precision } u \text { ) } \\
\text { (insion } u \text { ) }
\end{array}
$$

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Solve $A x_{0}=b$ by LU factorization
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\left.r_{i}=b-A x_{i} \quad \quad \text { (in precision } u\right)
$$

$$
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\end{array}
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## "Fixed-Precision"

[Jankowski and Woźniakowski, 1977], [Skeel, 1980], [Higham, 1991]

## Iterative Refinement for $A x=b$

As long as $\kappa_{\infty}(A) \leq u^{-1}$,

$$
\left.\operatorname{cond}(A, x)=\left\|\left|A^{-1}\right||A||x|\right\|_{\infty} /\|x\|_{\infty}\right)
$$

- relative forward error is $O(u) \operatorname{cond}(\boldsymbol{A}, \boldsymbol{x})$
- relative normwise and componentwise backward errors are $O(u)$

Solve $A x_{0}=b$ by LU factorization
(in precision $u$ )
for $i=0$ : maxit

$$
r_{i}=b-A x_{i}
$$

(in precision $u$ )

$$
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[Jankowski and Woźniakowski, 1977], [Skeel, 1980], [Higham, 1991]

## Iterative Refinement for $A x=b$

Solve $A x_{0}=b$ by LU factorization
(in precision $u^{1 / 2}$ ) for $i=0$ : maxit

$$
r_{i}=b-A x_{i} \quad(\text { in precision } u)
$$

$$
\begin{array}{ll}
\text { Solve } A d_{i}=r_{i} & \text { via } d_{i}=U^{-1}\left(L^{-1} r_{i}\right) \\
x_{i+1}=x_{i}+d_{i} & \text { (in precision } u) \\
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"Low-precision factorization"
[Langou et al., 2006], [Arioli and Duff, 2009], [Hogg and Scott, 2010], [Abdelfattah et al., 2016]

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## Iterative Refinement for $A x=b$

3-precision iterative refinement [C. and Higham, 2018]
$u_{f}=$ factorization precision, $u=$ working precision, $u_{r}=$ residual precision

$$
u_{f} \geq u \geq u_{r}
$$

Solve $A x_{0}=b$ by LU factorization
for $i=0$ : maxit

$$
r_{i}=b-A x_{i}
$$

(in precision $u_{r}$ )
Solve $A d_{i}=r_{i}$
(in precision $u_{s}$ )
$x_{i+1}=x_{i}+d_{i}$
(in precision $u$ )
$u_{s}$ is the effective precision of the solve, with $u \leq u_{s} \leq u_{f}$

## Key Aspects of Analysis: Tighter Upper Bounds

Obtain tighter upper bounds:

Typical bounds used in analysis: $\left\|A\left(x-\hat{x}_{i}\right)\right\|_{\infty} \leq\|A\|_{\infty}\left\|x-\hat{x}_{i}\right\|_{\infty}$

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For a stable refinement scheme, in early stages we expect

$$
\frac{\left\|r_{i}\right\|}{\|A\|\left\|\hat{x}_{i}\right\|} \approx u \ll \frac{\left\|x-\hat{x}_{i}\right\|}{\|x\|} \longrightarrow \mu_{i} \ll 1
$$

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$$

But close to convergence,

$$
\left\|r_{i}\right\| \approx\|A\|\left\|x-\hat{x}_{i}\right\| \longrightarrow \mu_{i} \approx 1
$$

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Allow for general solver:
Let $u_{s}$ be the effective precision of the solve, with $u \leq u_{s} \leq u_{f}$

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Assume computed solution $\hat{d}_{i}$ to $A d_{i}=\hat{r}_{i}$ satisfies:

1. $\hat{d}_{i}=\left(I+u_{s} E_{i}\right) d_{i}, \quad u_{s}\left\|E_{i}\right\|_{\infty}<1$
$\rightarrow$ normwise relative forward error is bounded by multiple of $u_{s}$ and is less than 1

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u_{s}\left\|E_{i}\right\|_{\infty} \leq 3 n u_{f}\left\|\left|A^{-1}\|\hat{L}\| \widehat{U}\right|\right\|_{\infty}
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$$
u_{s}\left\|E_{i}\right\|_{\infty} \leq 3 n u_{f}\left\|| A ^ { - 1 } | \left|\hat{L}\|\widehat{U} \mid\|_{\infty}\right.\right.
$$

by multiple of $u_{s}$ and is less than 1
2. $\left\|\hat{r}_{i}-A \hat{d}_{i}\right\|_{\infty} \leq u_{s}\left(c_{1}\|A\|_{\infty}\left\|\hat{d}_{i}\right\|_{\infty}+c_{2}\left\|\hat{r}_{i}\right\|_{\infty}\right)$
$\rightarrow$ normwise relative backward error is at most $\max \left(c_{1}, c_{2}\right) u_{s}$

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$\rightarrow$ normwise relative backward error is at most $\max \left(c_{1}, c_{2}\right) u_{s}$

$$
\max \left(c_{1}, c_{2}\right) u_{S} \leq \frac{3 n u_{f}\|\mid \hat{L}\| \widehat{U}\| \|_{\infty}}{\|A\|_{\infty}}
$$

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\max \left(c_{1}, c_{2}\right) u_{s} \leq \frac{3 n u_{f}\left\|\left|\hat{L}\|\widehat{U} \mid\|_{\infty}\right.\right.}{\|A\|_{\infty}}
$$

3. $\left|\hat{r}_{i}-A \hat{d}_{i}\right| \leq u_{s} G_{i}\left|\hat{d}_{i}\right|$
$\rightarrow$ componentwise relative backward error is bounded by a multiple of $u_{s}$
$E_{i}, c_{1}, c_{2}$, and $G_{i}$ depend on $A, \hat{r}_{i}, n$, and $u_{s}$

## Key Aspects of Analysis: Effective Solve Precision

Allow for general solver:
Let $u_{s}$ be the effective precision of the solve, with $u \leq u_{s} \leq u_{f}$
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u_{s}\left\|E_{i}\right\|_{\infty} \leq 3 n u_{f}\left\|\left|A ^ { - 1 } \left\|\left|\hat{L}\|\widehat{U} \mid\|_{\infty}\right.\right.\right.\right.
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by multiple of $u_{s}$ and is less than 1
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3. $\left|\hat{r}_{i}-A \hat{d}_{i}\right| \leq u_{s} G_{i}\left|\hat{d}_{i}\right|$
$\rightarrow$ componentwise relative backward error is bounded by a multiple of $u_{s}$

$$
u_{s}\left\|G_{i}\right\|_{\infty} \leq 3 n u_{f}\|\hat{L}\| \widehat{U} \mid \|_{\infty}
$$

$E_{i}, c_{1}, c_{2}$, and $G_{i}$ depend on $A, \hat{r}_{i}, n$, and $u_{s}$

## Key Aspects of Analysis: Effective Solve Precision

Allow for general solver:
Let $u_{s}$ be the effective precision of the solve, with $u \leq u_{s} \leq u_{f}$

Assume computed solution $\hat{d}_{i}$ to $A d_{i}=\hat{r}_{i}$ satisfies:
example: LU solve:

$$
u_{s}=u_{f}
$$

1. $\quad \hat{d}_{i}=\left(I+u_{s} E_{i}\right) d_{i}, \quad u_{s}\left\|E_{i}\right\|_{\infty}<1$
$\rightarrow$ normwise relative forward error is bounded by multiple of $u_{s}$ and is less than 1
2. $\left\|\hat{r}_{i}-A \hat{d}_{i}\right\|_{\infty} \leq u_{s}\left(c_{1}\|A\|_{\infty}\left\|\hat{d}_{i}\right\|_{\infty}+c_{2}\left\|\hat{r}_{i}\right\|_{\infty}\right)$
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$$
\max \left(c_{1}, c_{2}\right) u_{s} \leq \frac{3 n u_{f}\left\|\left|\hat{L}\|\widehat{U} \mid\|_{\infty}\right.\right.}{\|A\|_{\infty}}
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3. $\left|\hat{r}_{i}-A \hat{d}_{i}\right| \leq u_{s} G_{i}\left|\hat{d}_{i}\right|$
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$E_{i}, c_{1}, c_{2}$, and $G_{i}$ depend on $A, \hat{r}_{i}, n$, and $u_{s}$

$$
u_{s}\left\|G_{i}\right\|_{\infty} \leq 3 n u_{f}\|\hat{L}\| \widehat{U} \mid \|_{\infty}
$$

## Forward Error for IR3

- Three precisions:
- $u_{f}$ : factorization precision
- $u$ : working precision
- $u_{r}$ : residual computation precision

$$
\begin{aligned}
\kappa_{\infty}(A) & =\left\|A^{-1}\right\|_{\infty}\|A\|_{\infty} \\
\operatorname{cond}(A) & =\left\|\left|A^{-1}\|A \mid\|_{\infty}\right.\right. \\
\operatorname{cond}(A, x) & =\left\|\left|A^{-1}\right||A||x|\right\|_{\infty} /\|x\|_{\infty}
\end{aligned}
$$

## Forward Error for IR3

- Three precisions:
- $u_{f}$ : factorization precision
- u: working precision
- $u_{r}$ : residual computation precision

$$
\begin{aligned}
\kappa_{\infty}(A) & =\left\|A^{-1}\right\|_{\infty}\|A\|_{\infty} \\
\operatorname{cond}(A) & =\left\|\left|A^{-1}\|A \mid\|_{\infty}\right.\right. \\
\operatorname{cond}(A, x) & =\left\|\left|A^{-1}\right||A|\right\| x \mid\left\|_{\infty} /\right\| x \|_{\infty}
\end{aligned}
$$

## Theorem [C. and Higham, SISC 40(2), 2018]

For IR in precisions $u_{f} \geq u \geq u_{r}$ and effective solve precision $u_{s}$, if

$$
\phi_{i} \equiv 2 u_{s} \min \left(\operatorname{cond}(A), \kappa_{\infty}(A) \mu_{i}\right)+u_{s}\left\|E_{i}\right\|_{\infty}
$$

is less than 1 , then the forward error is reduced on the $i$ th iteration by a factor $\approx \phi_{i}$ until an iterate $\hat{x}_{i}$ is produced for which

$$
\frac{\left\|x-\hat{x}_{i}\right\|_{\infty}}{\|x\|_{\infty}} \lesssim 4 N u_{r} \operatorname{cond}(A, x)+u
$$

where $N$ is the maximum number of nonzeros per row in $A$.

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$$

where $N$ is the maximum number of nonzeros per row in $A$.
Analogous traditional bounds: $\phi_{i} \equiv 3 n u_{f} \kappa_{\infty}(A)$

## Normwise Backward Error for IR3

## Theorem [C. and Higham, SISC 40(2), 2018]

For $\operatorname{IR}$ in precisions $u_{f} \geq u \geq u_{r}$ and effective solve precision $u_{s}$, if

$$
\phi_{i} \equiv\left(c_{1} \kappa_{\infty}(A)+c_{2}\right) u_{s}
$$

is less than 1 , then the residual is reduced on the $i$ th iteration by a factor $\approx \phi_{i}$ until an iterate $\hat{x}_{i}$ is produced for which

$$
\left\|b-A \hat{x}_{i}\right\|_{\infty} \lesssim N u\left(\|b\|_{\infty}+\|A\|_{\infty}\left\|\hat{x}_{i}\right\|_{\infty}\right),
$$

where $N$ is the maximum number of nonzeros per row in $A$.

## IR3: Summary

Standard (LU-based) IR in three precisions ( $u_{s}=u_{f}$ ) Half $\approx 10^{-4}$, Single $\approx 10^{-8}$, Double $\approx 10^{-16}$, Quad $\approx 10^{-34}$

|  |  |  |  | Backward error |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |

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|  |  |  |  |  | Backward error |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| LP fact. | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
|  | H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
|  | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
|  | S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
|  | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
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| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
|  | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| Fixed | S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
|  | H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
|  | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
|  | S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |

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|  |  |  |  |  | Backward error |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
|  | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
|  | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| Fixed | S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| Trad. | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
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|  |  |  |  |  | Backward error |  |  |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| New | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| New | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| Fixed | S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| Trad. | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
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|  |  |  |  |  | Backward error |  |  |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $\boldsymbol{u}$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| New | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | H | D | D | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| New | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
|  |  | Sixed | S | S | S | $10^{8}$ | $10^{-8}$ |
| $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |  |  |  |  |  |  |
| Trad. | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| New | S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |

$\Rightarrow$ Benefit of IR3 vs. "LP fact.": no $\operatorname{cond}(A, x)$ term in forward error

## IR3: Summary

Standard (LU-based) IR in three precisions ( $u_{s}=u_{f}$ ) Half $\approx 10^{-4}$, Single $\approx 10^{-8}$, Double $\approx 10^{-16}$, Quad $\approx 10^{-34}$

|  |  |  |  |  |  | Backward error |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\boldsymbol{u}_{\boldsymbol{f}}$ | $\boldsymbol{u}$ | $\boldsymbol{u}_{\boldsymbol{r}}$ | $\max \kappa_{\infty}(A)$ | norm | $\operatorname{comp}$ | Forward error |
| LP fact. | H | S | S | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| New | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
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| New | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
|  | Sixed | S | S | S | $10^{8}$ | $10^{-8}$ | $10^{-8}$ |
| Trad. | S | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $\operatorname{cond}(A, x) \cdot 10^{-8}$ |
| LP fact. | S | D | D | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $\operatorname{cond}(A, x) \cdot 10^{-16}$ |
| New | S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |

$\Rightarrow$ Benefit of IR3 vs. traditional IR: As long as $\kappa_{\infty}(A) \leq 10^{4}$, can use lower precision factorization w/no loss of accuracy!

$$
\begin{aligned}
& A=\text { gallery }(' r a n d s v d ', 100,1 e 3) \\
& b=\operatorname{randn}(100,1) \\
& \boldsymbol{\kappa}_{\infty}(\boldsymbol{A}) \approx 1 e 4
\end{aligned}
$$

Standard (LU-based) IR with $\boldsymbol{u}_{\boldsymbol{f}}$ : single, $\boldsymbol{u}$ : double, $\boldsymbol{u}_{r}$ : quad



$$
\begin{aligned}
& A=\text { gallery }(' r a n d s v d ', 100,1 e 7) \\
& b=\operatorname{randn}(100,1) \\
& \boldsymbol{\kappa}_{\infty}(\boldsymbol{A}) \approx 7 \mathrm{e} 7
\end{aligned}
$$

Standard (LU-based) IR with $\boldsymbol{u}_{\boldsymbol{f}}$ : single, $\boldsymbol{u}$ : double, $\boldsymbol{u}_{r}$ : quad



$$
\begin{aligned}
& A=\text { gallery }(' r a n d s v d ', 100,1 e 9) \\
& \mathrm{b}=\operatorname{randn}(100,1) \\
& \boldsymbol{\kappa}_{\infty}(\boldsymbol{A}) \approx 2 \mathrm{e} 10
\end{aligned}
$$

Standard (LU-based) IR with $\boldsymbol{u}_{\boldsymbol{f}}$ : single, $\boldsymbol{u}$ : double, $\boldsymbol{u}_{r}$ : quad



```
A = gallery('randsvd', 100, 1e9)
b = randn (100,1)
\kappa
```

Standard (LU-based) IR with $\boldsymbol{u}_{f}$ : single, $\boldsymbol{u}$ : double, $\boldsymbol{u}_{r}$ : quad



$$
\begin{aligned}
& \mathrm{A}=\text { gallery('randsvd', } 100,1 \mathrm{e} 9) \\
& \mathrm{b}=\operatorname{randn}(100,1) \\
& \boldsymbol{\kappa}_{\infty}(\boldsymbol{A}) \approx 2 \mathbf{e} 10
\end{aligned}
$$

Standard (LU-based) IR with $\boldsymbol{u}_{f}$ : double, $\boldsymbol{u}$ : double, $\boldsymbol{u}_{r}$ : quad



## GMRES-Based Iterative Refinement

- Observation [Rump, 1990]: if $\hat{L}$ and $\widehat{U}$ are computed LU factors of $A$ in precision $\boldsymbol{u}_{f}$, then

$$
\kappa_{\infty}\left(\widehat{U}^{-1} \hat{L}^{-1} A\right) \approx 1+\kappa_{\infty}(A) u_{f}
$$

even if $\kappa_{\infty}(A) \gg u_{f}^{-1}$.

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GMRES-IR [C. and Higham, SISC 39(6), 2017]

- To compute the updates $d_{i}$, apply GMRES to



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- To compute the updates $d_{i}$, apply GMRES to


Solve $A x_{0}=b$ by LU factorization
for $i=0$ : maxit

$$
r_{i}=b-A x_{i}
$$

Solve $A d_{i}=r_{i}$ via GMRES on $\tilde{A} d_{i}=\tilde{r}_{i}$
$x_{i+1}=x_{i}+d_{i}$

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$x_{i+1}=x_{i}+d_{i}$

```
A = gallery('randsvd', 100, 1e9, 2)
b = randn (100,1)
\kappa\infty
```

Standard (LU-based) IR with $u_{f}$ : single, $\boldsymbol{u}$ : double, $u_{r}$ : quad



```
A = gallery('randsvd', 100, 1e9, 2)
b = randn (100,1)
```

$$
\kappa_{\infty}(A) \approx 2 \mathrm{e} 10, \operatorname{cond}(A, x) \approx 5 \mathrm{e} 9, \kappa_{\infty}(\tilde{A}) \approx 2 \mathrm{e} 4
$$



Number of GMRES iterations: $(2,3)$

## GMRES-IR: Summary

GMRES-IR: Solve for $d_{i}$ via GMRES on $U^{-1} L^{-1} A d_{i}=U^{-1} L^{-1} r_{i}$

|  | GMRES-based IR in three precisions ( $u_{s}=u$ ) |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | Backward error |  | Forward error |
|  |  |  |  |  | norm | comp |  |
| LU-IR | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| GMRES-IR | H | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LU-IR | S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| GMRES-IR | S | D | Q | $10^{16}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| LU-IR | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| GMRES-IR | H | D | Q | $10^{12}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |

$\Rightarrow$ With GMRES-IR, lower precision factorization will work for higher $\kappa_{\infty}(A)$

## GMRES-IR: Summary

GMRES-IR: Solve for $d_{i}$ via GMRES on $U^{-1} L^{-1} A d_{i}=U^{-1} L^{-1} r_{i}$

| GMRES-based IR in three precisions ( $u_{s}=u$ ) |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | Backward error |  | Forward error |
|  | $u_{f}$ | $u$ | $u_{r}$ | $\max \kappa_{\infty}(A)$ | norm | comp |  |
| LU-IR | H | S | D | $10^{4}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| GMRES-IR | H | S | D | $10^{8}$ | $10^{-8}$ | $10^{-8}$ | $10^{-8}$ |
| LU-IR | S | D | Q | $10^{8}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| GMRES-IR | S | D | Q | $10^{16}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| LU-IR | H | D | Q | $10^{4}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
| GMRES-IR | H) | D | Q | $10^{12}$ | $10^{-16}$ | $10^{-16}$ | $10^{-16}$ |
|  |  |  |  |  | - | $\kappa_{\infty}(A$ | $u^{-1 / 2} u_{f}^{-1}$ |

$\Rightarrow$ As long as $\kappa_{\infty}(A) \leq 10^{12}$, can use half precision factorization and still obtain double precision accuracy!

## Comments and Caveats I

- Convergence tolerance $\tau$ for GMRES?
- Smaller $\tau \rightarrow$ more GMRES iterations, potentially fewer refinement steps
- Larger $\tau \rightarrow$ fewer GMRES iterations, potentially more refinement steps


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- What about overflow, underflow, subnormal numbers?
- Sophisticated scaling methods can help avoid this
- "Squeezing a Matrix into Half Precision, with an Application to Solving Linear Systems" [Higham, Pranesh, Zounon, 2019]


## Comments and Caveats II

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- e.g., if (normal) $\tilde{A}$ still has cluster of eigenvalues near origin, GMRES can stagnate until $n^{\text {th }}$ iteration, regardless of $\kappa_{\infty}(A)$ [Liesen and Tichý, 2004]
- Potential remedies: deflation, Krylov subspace recycling [C., Oktay, 2022], using additional preconditioner


## Performance Results (MAGMA)

- [Haidar, Tomov, Dongarra, Higham, 2018]

(b) Matrix of type 4: clustered singular values, $\sigma_{i}=\left(1, \cdots, 1, \frac{1}{\text { cond }}\right)$.


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- Depending on conditioning of $A$, applying $\tilde{A}$ to a vector must be done accurately (precision $u^{2}$ ) in each GMRES iteration
- Recent development of 5-precision GMRES-IR algorithm [Amestoy et al., 2021]
- For GMRES entirely in precision $u$,

$$
\kappa_{\infty}(A) \leq u^{-1 / 2} u_{f}^{-1} \rightarrow \kappa_{\infty}(A) \leq u^{-1 / 3} u_{f}^{-2 / 3}
$$

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- If $A$ is ill conditioned and LU factorization is performed in very low precision, it can be a poor preconditioner
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$$

- Why GMRES?
- Theoretical purposes: existing analysis and proof of backward stability [Paige, Rozložník, Strakoš, 2006]
- In practice, use any solver you want!


## GMRES-IR in Libraries and Applications

- MAGMA: Dense linear algebra routines for heterogeneous/hybrid architectures

```
magma / src / dxgesv_gmres_gpu.cpp
```

$\square$

```
DSGESV or DHGESV expert interface.
```

    It computes the solution to a real system of linear equations
        \(A * X=B, A^{* *} T * X=B\), or \(A^{* * H} * X=B\),
    where \(A\) is an \(N-b y-N\) matrix and \(X\) and \(B\) are \(N\)-by-NRHS matrices.
    the accomodate the Single Precision DSGESV and the Half precision dhgesv API.
    precision and iterative refinement solver are specified by facto_type, solver_type.
    For other API parameter please refer to the corresponding dsgesv or dhgesv.
    - NVIDIA's cuSOLVER Library
2.2.1.6. cusolverIRSRefinement_t

The cusolverIRSRefinement_t type indicates which solver type would be used for the specific cusolver function. Most of our experimentation shows that CUSOLVER_IRS_REFINE_GMRES is the best option.

| CUSOLVER_IRS_REFINE_GMRES | GMRES (Generalized Minimal Residual) based <br> iterative refinement solver. In recent study, the <br> GMRES method has drawn the scientific <br> community attention for its ability to be used as <br> refinement solver that outperforms the classical <br> iterative refinement method. based on our <br> experimentation, we recommend this setting. |
| :--- | :--- |

- In production codes: FK6D/ASGarD code (Oak Ridge National Lab, USA) for tokomak containment problem


## Least Squares Iterative Refinement

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck,1967): Least squares problem can be written as a linear system with square matrix of size $(m+n)$;

$$
\left[\begin{array}{cc}
I & A \\
A^{T} & 0
\end{array}\right]\left[\begin{array}{l}
r \\
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- Refinement proceeds as follows:

1. Compute "residuals"

$$
\left[\begin{array}{l}
f_{i} \\
g_{i}
\end{array}\right]=\left[\begin{array}{l}
b \\
0
\end{array}\right]-\left[\begin{array}{cc}
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$$

2. Solve for corrections

$$
\left[\begin{array}{cc}
I & A \\
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\end{array}\right]\left[\begin{array}{c}
\Delta r_{i} \\
\Delta x_{i}
\end{array}\right]=\left[\begin{array}{l}
f_{i} \\
g_{i}
\end{array}\right]
$$

3. Update "solution":

$$
\left[\begin{array}{l}
r_{i+1} \\
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\end{array}\right] \begin{aligned}
& \text { Results for 3-precision } \\
& \text { IR for linear systems } \\
& \text { also applies to least } \\
& \text { squares problems! }
\end{aligned} \quad \tilde{x}_{i+1}=\tilde{x}_{i}+d_{i}
$$

See [C., Higham, Pranesh, 2020]

## GMRES-IR with Inexact Preconditioners

- Existing analyses of GMRES-IR assume we use full LU factors
- In practice, often want to use approximate preconditioners (ILU, SPAI, etc.)


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- [Amestoy et al., 2022]
- Analysis of block low-rank (BLR) LU within GMRES-IR
- Analysis of use of static pivoting in LU within GMRES-IR


## GMRES-IR with Inexact Preconditioners

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- [Amestoy et al., 2022]
- Analysis of block low-rank (BLR) LU within GMRES-IR
- Analysis of use of static pivoting in LU within GMRES-IR
- [C., Khan, 2022]
- Analysis of sparse approximate inverse (SPAI) preconditioners within GMRES-IR


## SPAI Preconditioners

Goal: Construct sparse matrix $M \approx A^{-1}$ (for survey see [Benzi, 2002])
Approach of [Grote, Huckle, 1997]: Construct columns $m_{k}$ of $M$ dynamically

Given matrix $A$, initial sparsity structure $J$, and tolerance $\varepsilon$ For each column $k$ :

Compute QR factorization of submatrix of $A$ defined by $J$
Use QR factorization to solve $\min _{m_{k}}\left\|e_{k}-A m_{k}\right\|_{2}$
If $\left\|r_{k}\right\|_{2}=\left\|e_{k}-A m_{k}\right\|_{2} \leq \varepsilon$ break;
Else
add select nonzeros to $J$, repeat.

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Benefits: Highly parallelizable
But construction can still be costly, esp. for large-scale problems
[Gao, Chen, He, 2021], [Chao, 2001], [Benzi, Tůma, 1999], [He, Yin, Gao, 2020]

## SPAI Preconditioners in Low Precision

What is the effect of using low precision in SPAI construction?

Notes and assumptions:

- We will assume that the SPAI construction is performed in some precision $u_{f}$
- We will denote quantities computed in finite precision with hats
- In our application, we want a left preconditioner, so we will run the algorithm on $A^{T}$ and set $M \leftarrow M^{T}$.
- We will assume that the QR factorization of the submatrix of $A^{T}$ is computed fully using HouseholderQR/TSQR


## SPAI Preconditioners in Low Precision

Two interesting questions:

1. Assuming we impose no maximum sparsity pattern on $\widehat{M}$, under what constraint on $u_{f}$ can we guarantee that $\left\|\hat{r}_{k}\right\|_{2} \leq \varepsilon$, with $\hat{r}_{k}=f l_{u_{f}}\left(e_{k}-\right.$ $A^{T} \widehat{m}_{k}^{T}$ ) for the computed $\widehat{m}_{k}^{T}$ ?

## SPAI Preconditioners in Low Precision

Two interesting questions:

1. Assuming we impose no maximum sparsity pattern on $\widehat{M}$, under what constraint on $\boldsymbol{u}_{f}$ can we guarantee that $\left\|\hat{r}_{k}\right\|_{2} \leq \varepsilon$, with $\hat{r}_{k}=f l_{u_{f}}\left(e_{k}-\right.$ $A^{T} \widehat{m}_{k}^{T}$ ) for the computed $\widehat{m}_{k}^{T}$ ?
2. Assume that when $M$ is computed in exact arithmetic, we quit as soon as $\left\|r_{k}\right\| \leq \varepsilon$. For $\widehat{M}$ computed in precision $u_{f}$ with the same sparsity pattern as $M$, what is $\left\|e_{k}-A^{T} \widehat{m}_{k}^{T}\right\|_{2}$ ?

## SPAI Preconditioning in Low Precision

Using standard rounding error analysis and perturbation results for LS problems, we have

$$
\left\|\hat{r}_{k}\right\|_{2} \leq n^{3} u_{f}\left\|\left|e_{k}\right|+\left|A^{T}\right|\left|\widehat{m}_{k}^{T}\right|\right\|_{2} .
$$

So in order to guarantee we eventually reach a solution with $\left\|\hat{r}_{k}\right\|_{2} \leq \varepsilon$, we need

$$
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$$

$\rightarrow$ problem must not be so ill-conditioned WRT $\boldsymbol{u}_{f}$ that we incur an error greater than $\varepsilon$ just computing the residual

## SPAI Preconditioning in Low Precision

Can turn this into the looser but more descriptive a priori bound:

$$
\operatorname{cond}_{2}\left(A^{T}\right) \lesssim \varepsilon \boldsymbol{u}_{f}^{-1},
$$

where $\operatorname{cond}_{2}\left(A^{T}\right)=\left\|\left|A^{-T}\right|\left|A^{T}\right|\right\|_{2}$.

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Another view: with a given matrix $A$ and a given precision $\boldsymbol{u}_{f}$, one must set $\boldsymbol{\varepsilon}$ such that

$$
\varepsilon \geq u_{f} \operatorname{cond}_{2}\left(A^{T}\right)
$$

Confirms intuition: The more approximate the inverse, the lower the precision we can use.

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Confirms intuition: The more approximate the inverse, the lower the precision we can use.

Resulting bounds for $\widehat{M}$ :

$$
\left\|I-A^{T} \widehat{M}^{T}\right\|_{F} \leq 2 \sqrt{n} \varepsilon, \quad\|I-\widehat{M} A\|_{\infty} \leq 2 n \varepsilon
$$

## Size of SPAI Preconditioner in Low Precision

How does precision used affect the number of nonzeros in $\widehat{M}$ ?


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## Second Question

Assume that when $M$ is computed in exact arithmetic, we quit as soon as $\left\|r_{k}\right\| \leq \varepsilon$. For $\widehat{M}$ computed in precision $\boldsymbol{u}_{\boldsymbol{f}}$ with the same sparsity pattern as $M$, what is $\left\|e_{k}-A^{T} \widehat{m}_{k}^{T}\right\|_{2}$ ?

## Second Question

Assume that when $M$ is computed in exact arithmetic, we quit as soon as $\left\|r_{k}\right\| \leq \varepsilon$. For $\widehat{M}$ computed in precision $\boldsymbol{u}_{\boldsymbol{f}}$ with the same sparsity pattern as $M$, what is $\left\|e_{k}-A^{T} \widehat{m}_{k}^{T}\right\|_{2}$ ?

In this case, we obtain the bound

$$
\|I-\widehat{M} A\|_{\infty} \leq n\left(\varepsilon+n^{7 / 2} u_{f} \kappa_{\infty}(A)\right)
$$

$\rightarrow$ If $\kappa_{\infty}(A) \gg \varepsilon \boldsymbol{u}_{f}^{-1}$, then computed $\widehat{M}$ with same sparsity structure as $M$ can be of much lower quality.

## SPAI-GMRES-IR

## SPAI-GMRES-IR

To compute the updates $d_{i}$, apply GMRES to $\widehat{M} A d_{i}=\widehat{M} r_{i}$

Solve $\widehat{M} A x_{0}=\widehat{M} b$ for $i=0$ : maxit

$$
r_{i}=b-A x_{i}
$$

Solve $A d_{i}=r_{i} \quad$ via GMRES on $\widehat{M} A d_{i}=\widehat{M} r_{i}$

$$
x_{i+1}=x_{i}+d_{i}
$$

## Low Precision SPAI within GMRES-IR

Using $\widehat{M}$ computed in precision $\boldsymbol{u}_{\boldsymbol{f}}$, for the preconditioned system $\tilde{A}=\widehat{M} A$,

$$
\kappa_{\infty}(\tilde{A}) \leqq(1+2 n \varepsilon)^{2} .
$$

steam3

saylr1


## Low Precision SPAI within GMRES-IR

To guarantee that both SPAI construction will complete and the GMRESbased iterative refinement scheme will converge, we must have roughly

$$
n u_{f} \operatorname{cond}_{2}\left(A^{T}\right) \lesssim n \varepsilon \lesssim u^{-1 / 2} .
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\hat{M} \text { can be } \\
\text { constructed }
\end{array}} \lesssim u^{-1 / 2} .
$$

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If $\varepsilon$ satisfies these constraints, then the constraints on condition number for forward and backward errors to converge are the same as for GMRES-IR with full LU factorization.

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If $\varepsilon$ satisfies these constraints, then the constraints on condition number for forward and backward errors to converge are the same as for GMRES-IR with full LU factorization.

Compared to GMRES-IR with full LU factorization, in general expect slower convergence, but much sparser preconditioner.

## SPAI-GMRES-IR Example

Matrix: steam1, $n=240, \mathrm{nnz}=2,248, \kappa_{\infty}(A)=3 \cdot 10^{7}, \operatorname{cond}\left(A^{T}\right)=3 \cdot 10^{3}$


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$\left(u_{f}, u, u_{r}\right)=($ single, double, quad $)$


$$
\mathrm{nnz}(L+U)=13,765
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$\left(u_{f}, u, u_{r}\right)=($ single, double, quad $)$

$n n z(L+U)=13,765$

$n n z(M)=2,248$

Is there a point in using precision higher than that dictated by $u_{f} \operatorname{cond}{ }_{2}\left(A^{T}\right) \leq \varepsilon$ ? Matrix: bfwa782, $n=782, \mathrm{nnz}=7514, \kappa_{\infty}(A)=7 \cdot 10^{3}, \operatorname{cond}\left(A^{T}\right)=1 \cdot 10^{3}$

$$
\left(u_{f}, u, u_{r}\right)=(\text { half, single, double })
$$

| Preconditioner | $\kappa_{\infty}(\tilde{A})$ | Precond. nnz | GMRES-IR steps/iteration |
| :--- | :---: | :---: | :--- |
| SPAI $(\varepsilon=0.2)$ | $2.1 e+02$ | 28053 | $67(31,36)$ |
| SPAI $(\varepsilon=0.5)$ | $9.7 e+02$ | 7528 | $153(71,82)$ |

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| :--- | :---: | :---: | :--- |
| SPAI $(\varepsilon=0.2)$ | $2.2 e+02$ | 26801 | $69(32,37)$ |
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## Related and Current Work

- Multistage mixed precision iterative refinement [Oktay, C., 2021] If IR not converging, first try changing the solver before increasing precision
- Low-precision randomized preconditioners [C., Daužickaité, 2022] Single-pass Nyström can be run in precision $u_{p} \approx \frac{\lambda_{k+1}}{\sqrt{n} \lambda_{1}}$ without affecting the quality of limited memory preconditioner.
- Low-precision in ILU-type preconditioners What can we prove?


## Summary and Takeaway

- We now have a multi-precision ecosystem
- Huge opportunities for using mixed precision in matrix computations
- But also big challenges!


## Thank You!

## carson@karlin.mff.cuni.cz

 www.karlin.mff.cuni.cz/ ${ }^{\text {carson/ }}$
## Key Analysis Innovations I

$$
\begin{gathered}
\left\|r_{i}\right\|_{2}=\mu_{i}^{(2)}\|A\|_{2}\left\|x-\hat{x}_{i}\right\|_{2} \\
x-\hat{x}_{i}=V \Sigma^{-1} U^{T} r_{i}=\sum_{j=1}^{n} \frac{\left(u_{j}^{T} r_{i}\right) v_{j}}{\sigma_{j}} \quad\left(A=U \Sigma V^{T}\right)
\end{gathered}
$$

## Key Analysis Innovations I

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\left\|x-\hat{x}_{i}\right\|_{2}^{2} \geq \sum_{j=n+1-k}^{n} \frac{\left(u_{j}^{T} r_{i}\right)^{2}}{\sigma_{j}^{2}} \geq \frac{1}{\sigma_{n+1-k}^{2}} \sum_{j=n+1-k}^{n}\left(u_{j}^{T} r_{i}\right)^{2}=\frac{\left\|P_{k} r_{i}\right\|_{2}^{2}}{\sigma_{n+1-k}^{2}} \\
\quad \text { where } P_{k}=U_{k} U_{k}^{T}, U_{k}=\left[u_{n+1-k}, \ldots, u_{n}\right]
\end{gathered}
$$

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x-\hat{x}_{i}=V \Sigma^{-1} U^{T} r_{i}=\sum_{j=1}^{n} \frac{\left(u_{j}^{T} r_{i}\right) v_{j}}{\sigma_{j}} \quad\left(A=U \Sigma V^{T}\right) \\
\left\|x-\hat{x}_{i}\right\|_{2}^{2} \geq \sum_{j=n+1-k}^{n} \frac{\left(u_{j}^{T} r_{i}\right)^{2}}{\sigma_{j}^{2}} \geq \frac{1}{\sigma_{n+1-k}^{2}} \sum_{j=n+1-k}^{n}\left(u_{j}^{T} r_{i}\right)^{2}=\frac{\left\|P_{k} r_{i}\right\|_{2}^{2}}{\sigma_{n+1-k}^{2}} \\
\text { where } P_{k}=U_{k} U_{k}^{T}, U_{k}=\left[u_{n+1-k}, \ldots, u_{n}\right] \\
\mu_{i}^{(2)} \leq \frac{\left\|r_{i}\right\|_{2}}{\left\|P_{k} r_{i}\right\|_{2}} \frac{\sigma_{n+1-k}}{\sigma_{1}}
\end{gathered}
$$

## Key Analysis Innovations I

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- $\mu_{i}^{(2)} \ll 1$ if $r_{i}$ contains significant component in $\operatorname{span}\left(U_{k}\right)$ for any $k$ s.t. $\sigma_{n+1-k} \approx \sigma_{n}$
- In that case, $x-\hat{x}_{i}$ is not "typical", i.e., it contains large components in right singular vectors corresponding to small singular values of $A$
- Wilkinson (1977), comment in unpublished manuscript: $\mu_{i}^{(2)}$ increases with $i$


## Performance Results (MAGMA)

- [Haidar, Tomov, Dongarra, Higham, 2018]

(b) Matrix of type 4: clustered singular values, $\sigma_{i}=\left(1, \cdots, 1, \frac{1}{\text { cond }}\right)$.


## Randomized Limited Memory Preconditioners

Let $A \in \mathbb{R}^{n \times n}$ be a symmetric positive semidefinite matrix. Want to solve

$$
(A+\mu I) x=b
$$

where $\mu \geq 0$ is set so that $A+\mu I$ is positive definite. Assume $A$ has rapidly decreasing eigenvalues or cluster of large eigenvalues.

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Want to solve using PCG using spectral limited memory preconditioner [Gratton, Sartenaer, Tshimanga, 2011], [Tshimanga et al., 2008]:

$$
\begin{aligned}
P & =I-U U^{T}+\frac{1}{\alpha+\mu} U(\Theta+\mu I) U^{T} \\
P^{-1} & =I-U U^{T}+(\alpha+\mu) U(\Theta+\mu I)^{-1} U^{T}
\end{aligned}
$$

where columns of $U \in \mathbb{R}^{n \times k}$ are $k$ approximate eigenvectors of $A$ and $U^{T} U=I, \Theta$ is diagonal with approximations to eigenvalues of $A$, and $\alpha \geq 0$.

Used in data assimilation [Laloyaux et al., 2018], [Mogensen, Alonso Balmaseda, Weaver, 2012], [Moore et al., 2011], [Daužickaité, Lawless, Scott, van Leeuwen, 2021]

## Randomized Nyström Approximation

Want to compute a rank- $k$ approximation $A \approx U \Theta U^{T}$ via the randomized Nyström method.

Nyström approximation:

$$
A_{N}=(A Q)\left(Q^{T} A Q\right)^{+}(A Q)^{T}
$$

where $Q$ is an $n \times k$ sampling matrix (random projection).

## Randomized Nyström Approximation

In the case that $A$ is very large, matrix-matrix products with $A$ are the bottleneck.

This motivates the single-pass version of the Nyström method.
Stabilized Single-Pass Nyström method [Tropp et al., 2017]
Given sym. PSD matrix $A$, target rank $k$
$G=\operatorname{randn}(n, k)$
$[Q, \sim]=\operatorname{qr}(G, 0)$
$\boldsymbol{Y}=\boldsymbol{A} \boldsymbol{Q}$
Compute shift $v ; Y_{v}=Y+v Q$
$B=Q^{T} Y_{v}$
$C=\operatorname{chol}\left(\left(B+B^{T}\right) / 2\right)$
Solve $F=Y_{v} / C$
$[U, \Sigma, \sim]=\operatorname{svd}(F, 0)$
$\Theta=\max \left(0, \Sigma^{2}-v I\right)$

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$\Theta=\max \left(0, \Sigma^{2}-v I\right)$

Can we further reduce the cost of the matrix-matrix product with $A$ by using low precision?

## Error Bounds

$$
\left\|A-\hat{A}_{N}\right\|_{2}=\left\|A-A_{N}+A_{N}-\hat{A}_{N}\right\|_{2} \leq \underbrace{\left\|A-A_{N}\right\|_{2}}_{\begin{array}{c}
\text { exact } \\
\text { approximation }
\end{array}}+\underbrace{\left\|A_{N}-\hat{A}_{N}\right\|_{2}}_{\begin{array}{c}
\text { finite precision } \\
\text { error }
\end{array}}
$$

## Error Bounds

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\left\|A-\hat{A}_{N}\right\|_{2}=\left\|A-A_{N}+A_{N}-\hat{A}_{N}\right\|_{2} \leq \underbrace{A-A_{N} \|_{2}}_{\substack{\text { exact } \\ \text { approximation } \\ \text { error }}}+\underbrace{\left\|A_{N}-\hat{A}_{N}\right\|_{2}}_{\substack{\text { finite precision } \\ \text { error }}}
$$

Deterministic bound [Gittens, Mahoney, 2016]:

$$
\left\|A-A_{N}\right\|_{2} \leq \lambda_{k+1}+\left\|\Sigma_{2}^{1 / 2} U_{2}^{T} Q\left(U_{1} Q\right)^{+}\right\|_{2}^{2}
$$

with $A=\left[\begin{array}{ll}U_{1} & U_{2}\end{array}\right]\left[\begin{array}{ll}\Sigma_{1} & \\ & \Sigma_{2}\end{array}\right]\left[\begin{array}{ll}U_{1} & U_{2}\end{array}\right]^{T}$.

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Expected value bound [Frangella, Tropp, Udell, 2021]:

$$
\mathbb{E}\left\|A-A_{N}\right\|_{2} \leq \min _{2 \leq p \leq k-2}\left(\left(1+\frac{2(k-p)}{p-1}\right) \lambda_{k-p+1}+\frac{2 e^{2} k}{p^{2}-1} \sum_{j=k-p+1}^{n} \lambda_{j}\right)
$$

where $\lambda_{i} \geq \lambda_{i+1}$ are the eigenvalues of $A$.

## Finite Precision Error Bound

Finite precision error: $A_{N}-\hat{A}_{N}$

Assumptions:

- $A$ is stored in precision $u_{p}$ and matrix-matrix product $A Q$ is computed in precision $u_{p}$
- All other quantities stored and computed in precision $u \ll u_{p}$


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Interpretation: $\left\|A_{N}-\hat{A}_{N}\right\|_{2} \gtrsim\left\|A-A_{N}\right\|_{2}$ when

$$
\frac{\lambda_{k+1}}{\lambda_{1}} \lesssim \sqrt{n} u_{p}
$$

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The more approximate the Interpretation: $\left\|A_{N}-\hat{A}_{N}\right\|_{2} \gtrsim\left\|A-A_{N}\right\|_{2}$ when low-rank representation, the lower the precision we can use!

$$
\frac{\lambda_{k+1}}{\lambda_{1}} \lesssim \sqrt{n} u_{p}
$$

## Condition Number Bounds

Let $E=A-A_{N}, \varepsilon=A_{N}-\hat{A}_{N}$, and assume $(A+\mu I)$ is SPD.

Let

$$
\widehat{P}^{-1}=I-\widehat{U} \widehat{U}^{T}+\left(\hat{\lambda}_{k}+\mu\right) \widehat{U}(\widehat{\Theta}+\mu I)^{-1} \widehat{U}^{T}
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be the LMP preconditioner constructed using the mixed precision Nyström approximation $\hat{A}_{N}=\widehat{U} \widehat{\Theta} \widehat{U}^{T}$.

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be the LMP preconditioner constructed using the mixed precision Nyström approximation $\hat{A}_{N}=\widehat{U} \widehat{\widehat{U}} \widehat{U}^{T}$.

Then

$$
\max \left\{1, \frac{\hat{\lambda}_{k}+\mu-\|\varepsilon\|_{2}}{\mu+\lambda_{\min }(A)}\right\} \leq \kappa\left(\hat{P}^{-1 / 2}(A+\mu I) \hat{P}^{-1 / 2}\right) \leq 1+\frac{\hat{\lambda}_{k}+\|E\|_{2}+2\|\varepsilon\|_{2}}{\mu-\|\varepsilon\|_{2}}
$$

where the upper bound holds if $\mu>\|\mathcal{E}\|_{2}$.
Regardless of this constraint, if $A$ is positive definite, then

$$
\kappa\left(\hat{P}^{-1 / 2}(A+\mu I) \hat{P}^{-1 / 2}\right) \leq\left(\hat{\lambda}_{k}+\mu+\|E\|_{2}+\|\varepsilon\|_{2}\right)\left(\frac{1}{\hat{\lambda}_{k}+\mu}+\frac{\|\varepsilon\|_{2}+1}{\lambda_{\min }(A)+\mu}\right) .
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be the LMP preconditioner constructed using the mixed precision Nyström approximation $\hat{A}_{N}=\widehat{U} \widehat{\Theta} \widehat{U}^{T}$.

Then If $\mathcal{E}=0$, reduces to bounds of [Frangella, Tropp, Udell, 2021] for exact case.

$$
\max \left\{1, \frac{\hat{\lambda}_{k}+\mu-\|\varepsilon\|_{2}}{\mu+\lambda_{\min }(A)}\right\} \leq \kappa\left(\hat{P}^{-1 / 2}(A+\mu I) \hat{P}^{-1 / 2}\right) \leq 1+\frac{\hat{\lambda}_{k}+\|E\|_{2}+2\|\varepsilon\|_{2}}{\mu-\|\varepsilon\|_{2}}
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$$

## Numerical Experiment

Matrix: bcsstm07, $n=420$


$$
\begin{aligned}
& \lambda_{k+1} / \lambda_{1} \\
& \sqrt{n} u_{p}, u_{p}=\text { half } \\
& \sqrt{n} u_{p}, u_{p}=\text { single }
\end{aligned}
$$

## Numerical Experiment

Matrix: bcsstm07, $n=420$

exact

$$
\begin{aligned}
& \text { mixed, } u_{p}=\text { half } \\
& \text { mixed, } u_{p}=\text { single } \\
& \text { mixed, } u_{p}=\text { double }
\end{aligned}
$$

mean finite prec. error, $\left\|A_{N}-\hat{A}_{N}\right\|_{2}$


## Numerical Experiment


unpreconditioned exact
$\square$ mixed, $u_{p}=$ half
$\square$ mixed, $u_{p}=$ single
$\square$ mixed, $u_{p}=$ double

PCG iteration count


## GMRES-IR for Least Squares

- Similar to the linear system case, we can use a lower precision factorization for even more ill-conditioned problems if we improve the effective precision of the solver
- Again, don't want to compute an LU factorization of the augmented system
- How can we use existing QR factors to construct an effective and inexpensive preconditioner for the augmented system?
- Note that augmented system is a saddle-point system; lots of existing work (block diagonal, triangular, constraint-based, ... )


## GMRES-IR for Least Squares

- Ex: block diagonal preconditioner ([Murphy, Golub, Wathen, 2000], [Ipsen, 2001])

$$
\left[\begin{array}{cc}
\alpha I & 0 \\
0 & \frac{1}{\alpha} \hat{R}^{T} \hat{R}
\end{array}\right]=\left[\begin{array}{cc}
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- Assuming QR factorization is exact,

$$
M_{2}^{-1} M_{1}^{-1} \tilde{A}=\left[\begin{array}{cc}
I & \frac{1}{\alpha} A \\
\alpha \hat{R}^{-1} \hat{R}^{-T} A^{T} & 0
\end{array}\right]
$$

is nonsymmetric, diagonalizable, with eigenvalues $\left\{1, \frac{1}{2}(1 \pm \sqrt{5})\right\}$.

- However, condition number can still be quite large; unsuitable for proving backward stability of GMRES


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- However, condition number can still be quite large; unsuitable for proving backward stability of GMRES
- If we take split preconditioner

$$
M_{1}^{-1} \tilde{A} M_{2}^{-1}=\left[\begin{array}{cc}
I & A \hat{R} \\
\hat{R}^{-T} A^{T} & 0
\end{array}\right]
$$

we will have a well-conditioned system

- However, split-preconditioned GMRES is not backward stable
- Potentially useful in practice, not but in theory


## GMRES-IR for Least Squares

- One option:

$$
M=\left[\begin{array}{cc}
\alpha I & \hat{Q}_{1} \hat{R} \\
\hat{R}^{T} \widehat{Q}_{1}^{T} & 0
\end{array}\right]
$$

- Then we can prove that for the left-preconditioned system,

$$
\kappa\left(M^{-1} \tilde{A}\right) \leq\left(1+u_{f} c \kappa(A)\right)^{2}
$$

where $c=O\left(m^{2}\right)$, where we note this bound is pessimistic.

- Thus even if $\kappa(A) \gg \boldsymbol{u}_{f}^{-1}$, the preconditioned system can still be reasonably well conditioned


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- GMRES run on $\tilde{A}$ with left-preconditioner $M$ gives

$$
u_{s}\left\|E_{i}\right\|_{\infty} \equiv u f(m+n) \kappa_{\infty}\left(M^{-1} \tilde{A}\right)
$$

where $f$ is a quadratic polynomial

## GMRES-IR for Least Squares

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where $f$ is a quadratic polynomial

- So for GMRES-based LSIR, $u_{s} \equiv u$; expect convergence of forward error when $\kappa_{\infty}(A)<u^{-1 / 2} \boldsymbol{u}_{f}^{-1}$
[C., Higham, Pranesh, SISC 2020]

