# Mixed Precision Iterative Refinement

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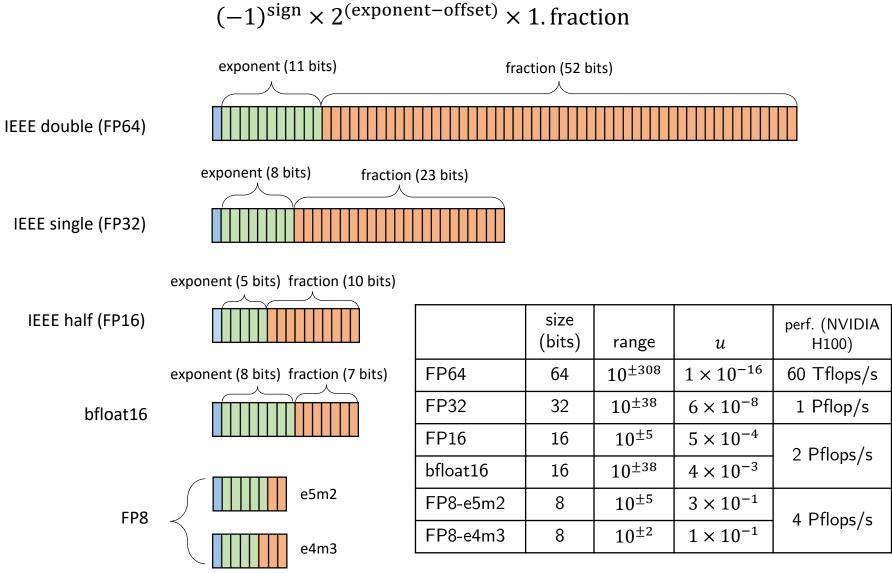
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Collaborators: Nicholas J. Higham (Manchester), Srikara Pranesh (V-Labs), Noaman Khan (Charles Univ.)

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## Floating Point Formats



# Hardware Support for Multiprecision Computation

Use of low precision in machine learning has driven emergence of lowprecision capabilities in hardware:

- Half precision (FP16) defined as storage format in 2008 IEEE standard
- ARM NEON: SIMD architecture, instructions for 8x16-bit, 4x32-bit, 2x64-bit
- AMD Radeon Instinct MI25 GPU, 2017:
  - single: 12.3 TFLOPS, half: 24.6 TFLOPS
- NVIDIA Tesla P100, 2016: native ISA support for 16-bit FP arithmetic
- NVIDIA Tesla V100, 2017: tensor cores for half precision;

4x4 matrix multiply in one clock cycle

- double: 7 TFLOPS, half+tensor: 112 TFLOPS (16x!)
- Google's Tensor processing unit (TPU)
- NVIDIA A100, 2020: tensor cores with multiple supported precisions: FP16, FP64, Binary, INT4, INT8, bfloat16
- NVIDIA H100, 2022: now with quarter-precision (FP8) tensor cores
- Exascale supercomputers: Expected extensive support for reduced-precision arithmetic (Frontier: FP64, FP32, FP16, bfloat16, INT8, INT4)

### Mixed precision in NLA

- BLAS: cuBLAS, MAGMA, [Agullo et al. 2009], [Abdelfattah et al., 2019], [Haidar et al., 2018]
- Iterative refinement:
  - Long history: [Wilkinson, 1963], [Moler, 1967], [Stewart, 1973], ...
  - More recently: [Langou et al., 2006], [C., Higham, 2017], [C., Higham, 2018], [C., Higham, Pranesh, 2020], [Amestoy et al., 2021]
- Matrix factorizations: [Haidar et al., 2017], [Haidar et al., 2018], [Haidar et al., 2020], [Abdelfattah et al., 2020]
- Eigenvalue problems: [Dongarra, 1982], [Dongarra, 1983], [Tisseur, 2001], [Davies et al., 2001], [Petschow et al., 2014], [Alvermann et al., 2019]
- Sparse direct solvers: [Buttari et al., 2008]
- Orthogonalization: [Yamazaki et al., 2015]
- Multigrid: [Tamstorf et al., 2020], [Richter et al., 2014], [Sumiyoshi et al., 2014], [Ljungkvist, Kronbichler, 2017, 2019]
- (Preconditioned) Krylov subspace methods: [Emans, van der Meer, 2012], [Yamagishi, Matsumura, 2016], [C., Gergelits, Yamazaki, 2021], [Clark, 2019], [Anzt et al., 2019], [Clark et al., 2010], [Gratton et al., 2020], [Arioli, Duff, 2009], [Hogg, Scott, 2010]

For survey and references, see [Abdelfattah et al., IJHPC, 2021]

- Supercomputers traditionally ranked by performance on high-performance LINPACK (HPL) benchmark
  - Solves dense Ax = b via Gaussian elimination with partial pivoting
- HPL-AI: Like HPL, solves dense Ax = b, results still to double precision accuracy
  - But achieves this via mixed-precision GMRES-based iterative refinement

#### June 2022

Rank	Site	Computer	Cores	HPL-AI (Eflop/s)	TOP500 Rank	HPL Rmax (Eflop/s)	Speedup
1	DOE/SC/ORNL, USA	Frontier	8,730,112	6.861	1	1.102	6.2
2	RIKEN, Japan	Fugaku	7,630,848	2.000	2	0.4420	4.5
3	DOE/SC/ORNL, USA	Summit	2,414,592	1.411	4	0.1486	9.5
4	NVIDIA, USA	Selene	555,520	0.630	8	0.0630	9.9
5	DOE/SC/LBNL, USA	Perlmutter	761,856	0.590	7	0.0709	8.3
6	FZJ, Germany	JUWELS BM	449,280	0.470	11	0.0440	10.0
7	University of Florida, USA	HiPerGator	138,880	0.170	34	0.0170	9.9
8	SberCloud, Russia	Christofari Neo	98,208	0.123	47	0.0120	10.3
9	DOE/SC/ANL, USA	Polaris	259,840	0.114	14	0.0238	4.8
10	ITC, Japan	Wisteria	368,640	0.100	20	0.0220	4.5

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Iterative refinement: well-established method for improving an approximate solution to Ax = b

A is  $n \times n$  and nonsingular; u is unit roundoff

Solve  $Ax_0 = b$  by LU factorization for i = 0: maxit  $r_i = b - Ax_i$ Solve  $Ad_i = r_i$  via  $d_i = U^{-1}(L^{-1}r_i)$  $x_{i+1} = x_i + d_i$ 

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"Traditional"

(high-precision residual computation)

[Wilkinson, 1948] (fixed point), [Moler, 1967] (floating point)

 $\kappa_{\infty}(A) = \|A^{-1}\|_{\infty} \|A\|_{\infty}$ 

As long as  $\kappa_{\infty}(A) \leq u^{-1}$ ,

- relative forward error is O(u)
- relative normwise and componentwise backward errors are O(u)

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#### "Fixed-Precision"

[Jankowski and Woźniakowski, 1977], [Skeel, 1980], [Higham, 1991]

 $cond(A, x) = || |A^{-1}||A||x| ||_{\infty} / ||x||_{\infty}$ 

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Solve 
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[Langou et al., 2006], [Arioli and Duff, 2009], [Hogg and Scott, 2010], [Abdelfattah et al., 2016]

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3-precision iterative refinement [C. and Higham, 2018]  $u_f$  = factorization precision, u = working precision,  $u_r$  = residual precision  $u_f \ge u \ge u_r$ 

Solve 
$$Ax_0 = b$$
 by LU factorization(in precision  $u_f$ )for  $i = 0$ : maxit(in precision  $u_r$ ) $r_i = b - Ax_i$ (in precision  $u_r$ )Solve  $Ad_i = r_i$ (in precision  $u_s$ ) $x_{i+1} = x_i + d_i$ (in precision  $u$ )

 $u_s$  is the *effective precision* of the solve, with  $u \leq u_s \leq u_f$ 

Obtain tighter upper bounds:

Typical bounds used in analysis:  $||A(x - \hat{x}_i)||_{\infty} \le ||A||_{\infty} ||x - \hat{x}_i||_{\infty}$ 

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Define  $\mu_i$ :  $||A(x - \hat{x}_i)||_{\infty} = \mu_i ||A||_{\infty} ||x - \hat{x}_i||_{\infty}$ 

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For a stable refinement scheme, in early stages we expect

$$\frac{\|r_i\|}{\|A\|\|\hat{x}_i\|} \approx u \ll \frac{\|x - \hat{x}_i\|}{\|x\|} \longrightarrow \mu_i \ll 1$$

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But close to convergence,

$$||r_i|| \approx ||A|| ||x - \hat{x}_i|| \longrightarrow \mu_i \approx 1$$

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Assume computed solution  $\hat{d}_i$  to  $Ad_i = \hat{r}_i$  satisfies:

1.  $\hat{d}_i = (I + \mathbf{u}_s E_i)d_i$ ,  $\mathbf{u}_s ||E_i||_{\infty} < 1$ 

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 $\mathbf{u}_{s} \| E_{i} \|_{\infty} \leq 3n \mathbf{u}_{f} \| |A^{-1}| |\hat{L}| |\hat{U}| \|_{\infty}$ 

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2. 
$$\|\hat{r}_i - A\hat{d}_i\|_{\infty} \le u_s(c_1 \|A\|_{\infty} \|\hat{d}_i\|_{\infty} + c_2 \|\hat{r}_i\|_{\infty})$$
  
→ normwise relative backward error is at most  $\max(c_1, c_2) u_s$ 

	exa	mple	e: Ll	Jsc	lve:		
$\boldsymbol{u}_{c} \  E_{i}$	ill a	< 3n	$ \mathbf{u}_{\mathbf{f}}  $	$ A^{-1} $	-   <u>}</u>	$ \widehat{U}   $	

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 $\mathbf{u}_{\mathbf{s}} \| E_i \|_{\infty} \le 3n \mathbf{u}_{\mathbf{f}} \| |A^{-1}| |\hat{L}| |\hat{U}| \|_{\infty}$ 

example: LU solve:

max(c, c) u	$3n\mathbf{u}_{f} \   \hat{L}   \hat{U}  \ _{\infty}$
$\max(c_1, c_2)  \boldsymbol{u}_s$	$\leq   A  _{\infty}$

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$$\|\hat{r}_i - A\hat{d}_i\|_{\infty} \leq u_s(c_1 \|A\|_{\infty} \|\hat{d}_i\|_{\infty} + c_2 \|\hat{r}_i\|_{\infty})$$
  
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3. 
$$\left|\hat{r}_i - A\hat{d}_i\right| \le \mathbf{u}_s G_i |\hat{d}_i|$$

 $\rightarrow\,$  componentwise relative backward error is bounded by a multiple of  $u_{s}$ 

 $E_i, c_1, c_2$ , and  $G_i$  depend on A,  $\hat{r}_i, n$ , and  $u_s$ 

 $\mathbf{u}_{\mathbf{s}} \| E_i \|_{\infty} \le 3n \mathbf{u}_{\mathbf{f}} \| |A^{-1}| |\hat{L}| |\hat{U}| \|_{\infty}$ 

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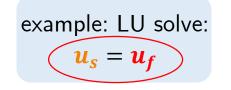
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# Forward Error for IR3

- Three precisions:
  - $u_f$ : factorization precision
  - *u*: working precision
  - $u_r$ : residual computation precision

 $\kappa_{\infty}(A) = ||A^{-1}||_{\infty} ||A||_{\infty}$   $\operatorname{cond}(A) = |||A^{-1}||A||_{\infty}$  $\operatorname{cond}(A, x) = |||A^{-1}||A||x||_{\infty} / ||x||_{\infty}$ 

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#### Theorem [C. and Higham, SISC 40(2), 2018]

For IR in precisions  $u_f \ge u \ge u_r$  and effective solve precision  $u_s$ , if

 $\phi_i \equiv 2 \mathbf{u}_s \min(\operatorname{cond}(A), \kappa_\infty(A)\mu_i) + \mathbf{u}_s ||E_i||_\infty$ 

is less than 1, then the forward error is reduced on the *i*th iteration by a factor  $\approx \phi_i$  until an iterate  $\hat{x}_i$  is produced for which

$$\frac{\|x - \hat{x}_i\|_{\infty}}{\|x\|_{\infty}} \lesssim 4N\boldsymbol{u}_r \operatorname{cond}(A, x) + \boldsymbol{u},$$

where N is the maximum number of nonzeros per row in A.

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where N is the maximum number of nonzeros per row in A.

Analogous traditional bounds:  $\phi_i \equiv 3n u_f \kappa_{\infty}(A)$ 

### Normwise Backward Error for IR3

#### Theorem [C. and Higham, SISC 40(2), 2018]

For IR in precisions  $u_f \ge u \ge u_r$  and effective solve precision  $u_s$ , if

$$\phi_i \equiv (c_1 \kappa_\infty(A) + c_2) \mathbf{u}_s$$

is less than 1, then the residual is reduced on the *i*th iteration by a factor  $\approx \phi_i$  until an iterate  $\hat{x}_i$  is produced for which

 $\|b - A\hat{x}_i\|_{\infty} \leq N\boldsymbol{u}(\|b\|_{\infty} + \|A\|_{\infty}\|\hat{x}_i\|_{\infty}),$ 

where N is the maximum number of nonzeros per row in A.

				Backwai	rd error	
<b>u</b> <sub>f</sub>	u	$u_r$	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
Н	S	S	104	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Н	S	D	104	$10^{-8}$	$10^{-8}$	$10^{-8}$
Н	D	D	104	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
Н	D	Q	104	$10^{-16}$	$10^{-16}$	$10^{-16}$
S	S	S	10 <sup>8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
S	S	D	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	$10^{-8}$
S	D	D	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
S	D	Q	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	$10^{-16}$

					Backwar	rd error	
	<b>u</b> <sub>f</sub>	u	$u_r$	$\max \kappa_\infty(A)$	norm	comp	Forward error
LP fact.	Н	S	S	104	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	104	$10^{-8}$	$10^{-8}$	$10^{-8}$
LP fact.	Н	D	D	10 <sup>4</sup>	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
	Н	D	Q	104	$10^{-16}$	$10^{-16}$	$10^{-16}$
	S	S	S	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	S	S	D	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	$10^{-8}$
LP fact.	S	D	D	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
	S	D	Q	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>

					Backwai	rd error	
	<b>u</b> <sub>f</sub>	u	$u_r$	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	104	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	104	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LP fact.	Н	D	D	104	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
	Н	D	Q	104	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>
Fixed	S	S	S	10 <sup>8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	$cond(A, x) \cdot 10^{-8}$
	S	S	D	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LP fact.	S	D	D	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
	S	D	Q	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>

					Backward error		
	<b>u</b> <sub>f</sub>	u	<b>u</b> <sub>r</sub>	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	104	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
	Н	S	D	104	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LP fact.	Н	D	D	104	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
	Н	D	Q	104	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>
Fixed	S	S	S	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$cond(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 <sup>8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>
LP fact.	S	D	D	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	$cond(A, x) \cdot 10^{-16}$
	S	D	Q	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>

# IR3: Summary

Standard (LU-based) IR in three precisions  $(u_s = u_f)$ Half  $\approx 10^{-4}$ , Single  $\approx 10^{-8}$ , Double  $\approx 10^{-16}$ , Quad  $\approx 10^{-34}$ 

					Backwai	rd error	
	<b>u</b> <sub>f</sub>	u	<i>u</i> <sub>r</sub>	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	104	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
New	н	S	D	10 <sup>4</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>
LP fact.	Н	D	D	104	$10^{-16}$	10 <sup>-16</sup>	$cond(A, x) \cdot 10^{-16}$
New	н	D	Q	10 <sup>4</sup>	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>
Fixed	S	S	S	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LP fact.	S	D	D	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	$cond(A, x) \cdot 10^{-16}$
New	S	D	Q	10 <sup>8</sup>	10 <sup>-16</sup>	10 <sup>-16</sup>	10 <sup>-16</sup>

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Standard (LU-based) IR in three precisions  $(u_s = u_f)$ Half  $\approx 10^{-4}$ , Single  $\approx 10^{-8}$ , Double  $\approx 10^{-16}$ , Quad  $\approx 10^{-34}$ 

					Backwai	rd error	
	<b>u</b> <sub>f</sub>	u	<i>u</i> <sub>r</sub>	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	10 <sup>4</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	$\operatorname{cond}(A, x) \cdot 10^{-8}$
New	Н	S	D	10 <sup>4</sup>	$10^{-8}$	$10^{-8}$	$10^{-8}$
LP fact.	Н	D	D	10 <sup>4</sup>	10 <sup>-16</sup>	10 <sup>-16</sup>	$cond(A, x) \cdot 10^{-16}$
New	Н	D	Q	10 <sup>4</sup>	$10^{-16}$	$10^{-16}$	$10^{-16}$
Fixed	S	S	S	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	$cond(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	$10^{-8}$
LP fact.	S	D	D	10 <sup>8</sup>	10 <sup>-16</sup>	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
New	S	D	Q	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$10^{-16}$

 $\Rightarrow$  Benefit of IR3 vs. "LP fact.": no cond(A, x) term in forward error

# IR3: Summary

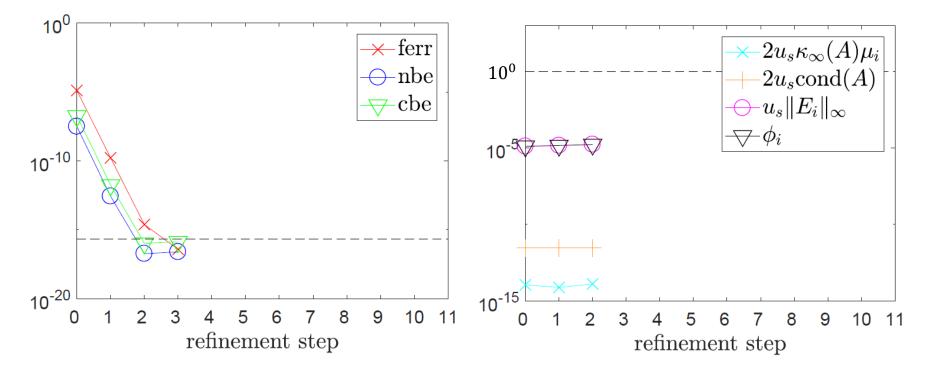
Standard (LU-based) IR in three precisions  $(u_s = u_f)$ Half  $\approx 10^{-4}$ , Single  $\approx 10^{-8}$ , Double  $\approx 10^{-16}$ , Quad  $\approx 10^{-34}$ 

					Backwai	rd error	
	<b>u</b> <sub>f</sub>	u	$u_r$	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LP fact.	Н	S	S	$10^{4}$	$10^{-8}$	$10^{-8}$	$cond(A, x) \cdot 10^{-8}$
New	н	S	D	104	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LP fact.	Н	D	D	$10^{4}$	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
New	Н	D	Q	$10^{4}$	$10^{-16}$	$10^{-16}$	$10^{-16}$
Fixed	S	S	S	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	$cond(A, x) \cdot 10^{-8}$
Trad.	S	S	D	10 <sup>8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>
LP fact.	S	D	D	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$cond(A, x) \cdot 10^{-16}$
New	S	D	Q	10 <sup>8</sup>	$10^{-16}$	$10^{-16}$	$10^{-16}$

⇒ Benefit of IR3 vs. traditional IR: As long as  $\kappa_{\infty}(A) \leq 10^4$ , can use lower precision factorization w/no loss of accuracy!

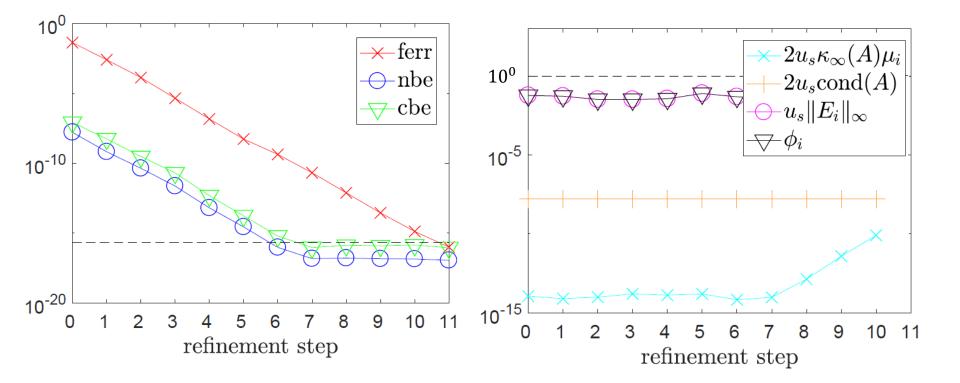
```
A = gallery('randsvd', 100, 1e3)
b = randn(100,1)
```

 $\kappa_\infty(A) pprox$  1e4



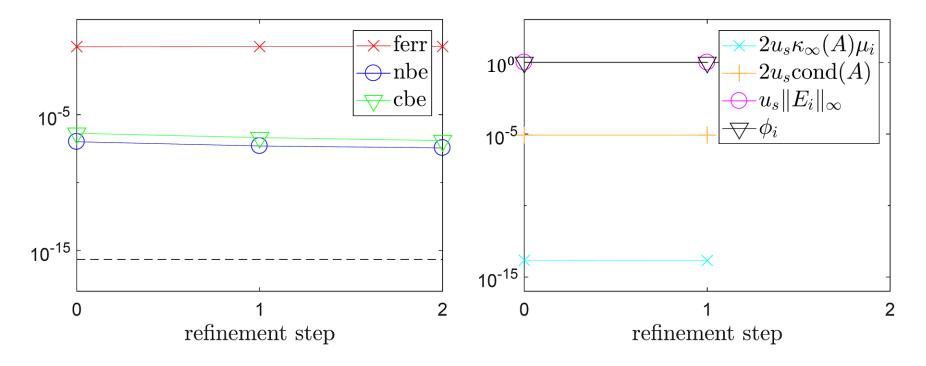
```
A = gallery('randsvd', 100, 1e7)
b = randn(100,1)
```

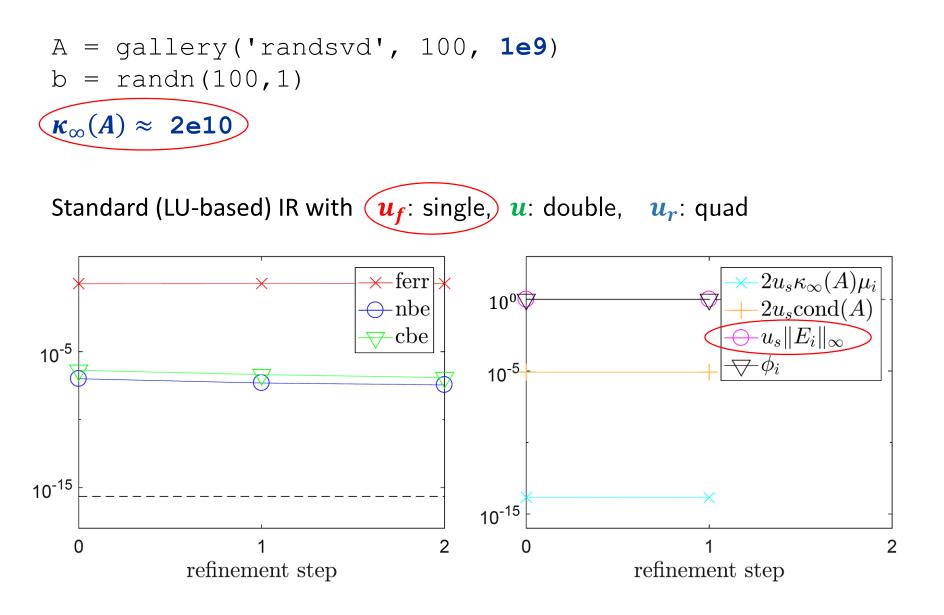
 $\kappa_{\infty}(A) \approx$  7e7

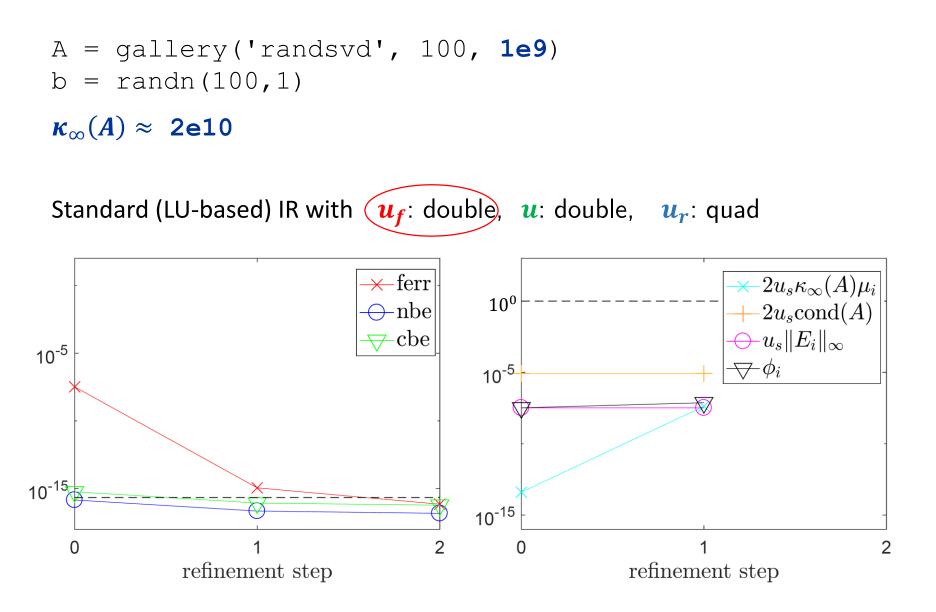


```
A = gallery('randsvd', 100, 1e9)
b = randn(100,1)
```

 $\kappa_{\infty}(A) \approx$  2e10







• Observation [Rump, 1990]: if  $\hat{L}$  and  $\hat{U}$  are computed LU factors of A in precision  $\boldsymbol{u}_{f}$ , then  $\kappa_{\infty}(\hat{U}^{-1}\hat{L}^{-1}A) \approx 1 + \kappa_{\infty}(A)\boldsymbol{u}_{f}$ ,

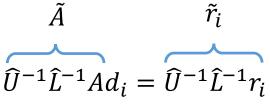
even if  $\kappa_{\infty}(A) \gg u_f^{-1}$ .

Observation [Rump, 1990]: if  $\hat{L}$  and  $\hat{U}$  are computed LU factors of A in precision  $\boldsymbol{u}_{f}$ , then •  $\kappa_{\infty}(\widehat{U}^{-1}\widehat{L}^{-1}A) \approx 1 + \kappa_{\infty}(A)\boldsymbol{u}_{\boldsymbol{f}},$ 

even if  $\kappa_{\infty}(A) \gg u_f^{-1}$ .

GMRES-IR [C. and Higham, SISC 39(6), 2017]

• To compute the updates  $d_i$ , apply GMRES to  $\widehat{U}^{-1}\widehat{L}^{-1}Ad_i = \widehat{U}^{-1}\widehat{L}^{-1}r_i$ 



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Â

 $r_i$ 

even if  $\kappa_{\infty}(A) \gg u_f^{-1}$ .

GMRES-IR [C. and Higham, SISC 39(6), 2017]

• To compute the updates  $d_i$ , apply GMRES to  $\hat{U}^{-1}\hat{L}^{-1}Ad_i = \hat{U}^{-1}\hat{L}^{-1}r_i$ 

Solve  $Ax_0 = b$  by LU factorization for i = 0: maxit  $r_i = b - Ax_i$ Solve  $Ad_i = r_i$  via GMRES on  $\tilde{A}d_i = \tilde{r}_i$  $x_{i+1} = x_i + d_i$ 

 Observation [Rump, 1990]: if L̂ and Û̂ are computed LU factors of A in precision u<sub>f</sub>, then
 κ<sub>∞</sub>(Û<sup>-1</sup>L̂<sup>-1</sup>A) ≈ 1 + κ<sub>∞</sub>(A)u<sub>f</sub>,

even if  $\kappa_{\infty}(A) \gg u_f^{-1}$ .

GMRES-IR [C. and Higham, SISC 39(6), 2017]

• To compute the updates  $d_i$ , apply GMRES to  $\hat{U}^{-1}\hat{L}^{-1}Ad_i = \hat{U}^{-1}\hat{L}^{-1}r_i$ 

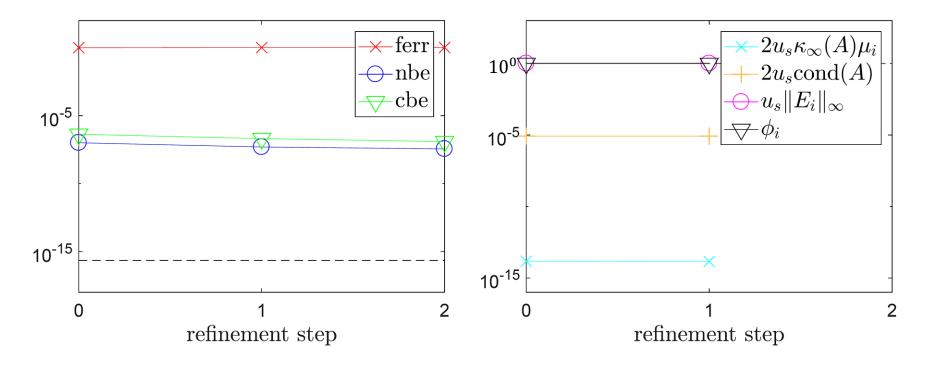
Solve 
$$Ax_0 = b$$
 by LU factorization  
for  $i = 0$ : maxit  
 $r_i = b - Ax_i$   
Solve  $Ad_i = r_i$  via GMRES on  $\tilde{A}d_i = \tilde{r}_i$   
 $x_{i+1} = x_i + d_i$ 

Â

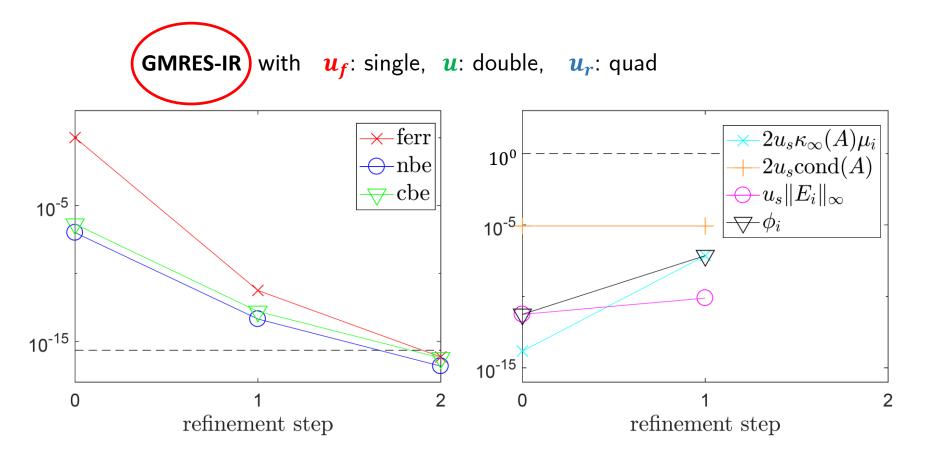
 $r_i$ 

```
A = gallery('randsvd', 100, 1e9, 2)
b = randn(100,1)
```

```
\kappa_{\infty}(A) \approx 2e10, \operatorname{cond}(A, x) \approx 5e9
```



 $\kappa_{\infty}(A) \approx$  2e10,  $\operatorname{cond}(A, x) \approx$  5e9,  $\kappa_{\infty}(\tilde{A}) \approx$  2e4



Number of GMRES iterations: (2,3)

# GMRES-IR: Summary

GMRES-IR: Solve for  $d_i$  via GMRES on  $U^{-1}L^{-1}Ad_i = U^{-1}L^{-1}r_i$ 

	UI UI				precision		
					Backwa	rd error	
	<b>u</b> <sub>f</sub>	u	<i>u</i> <sub>r</sub>	$\max \kappa_\infty(A)$	norm	comp	Forward error
LU-IR	Н	S	D	104	10 <sup>-8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>
GMRES-IR	Н	S	D	10 <sup>8</sup>	$10^{-8}$	$10^{-8}$	10 <sup>-8</sup>
LU-IR	S	D	Q	10 <sup>8</sup>	$10^{-16}$	10 <sup>-16</sup>	10 <sup>-16</sup>
GMRES-IR	S	D	Q	10 <sup>16</sup>	$10^{-16}$	$10^{-16}$	10 <sup>-16</sup>
LU-IR	Н	D	Q	104	$10^{-16}$	$10^{-16}$	10 <sup>-16</sup>
GMRES-IR	Н	D	Q	10 <sup>12</sup>	$10^{-16}$	$10^{-16}$	$10^{-16}$

GMRES-based IR in three precisions  $(u_s = u)$ 

 $\Rightarrow$ With GMRES-IR, lower precision factorization will work for higher  $\kappa_{\infty}(A)$ 

# GMRES-IR: Summary

GMRES-IR: Solve for  $d_i$  via GMRES on  $U^{-1}L^{-1}Ad_i = U^{-1}L^{-1}r_i$ 

					Backwa	rd error	
	u <sub>f</sub>	u	<b>u</b> <sub>r</sub>	$\max \kappa_{\infty}(A)$	norm	comp	Forward error
LU-IR	Н	S	D	104	10 <sup>-8</sup>	10 <sup>-8</sup>	10 <sup>-8</sup>
GMRES-IR	Н	S	D	10 <sup>8</sup>	$10^{-8}$	10 <sup>-8</sup>	$10^{-8}$
LU-IR	S	D	Q	10 <sup>8</sup>	10 <sup>-16</sup>	10 <sup>-16</sup>	10 <sup>-16</sup>
GMRES-IR	S	D	Q	10 <sup>16</sup>	$10^{-16}$	$10^{-16}$	$10^{-16}$
LU-IR	Н	D	Q	104	$10^{-16}$	10 <sup>-16</sup>	$10^{-16}$
GMRES-IR	Н	D	Q	10 <sup>12</sup>	$10^{-16}$	10 <sup>-16</sup>	$10^{-16}$
						$\rightarrow \kappa_{\infty}(A)$	$\leq u^{-1/2} u_f^{-1}$

GMRES-based IR in three precisions  $(u_s = u)$ 

⇒ As long as  $\kappa_{\infty}(A) \leq 10^{12}$ , can use half precision factorization and still obtain double precision accuracy!

- Convergence tolerance  $\tau$  for GMRES?
  - Smaller  $\tau \rightarrow$  more GMRES iterations, potentially fewer refinement steps
  - Larger  $\tau \rightarrow$  fewer GMRES iterations, potentially more refinement steps

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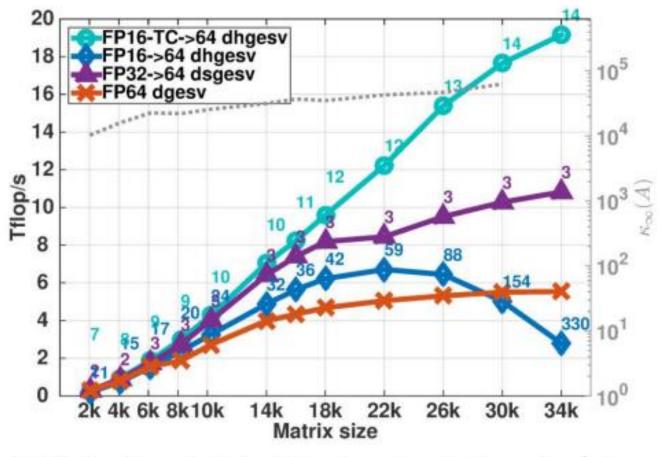
- What about overflow, underflow, subnormal numbers?
  - Sophisticated scaling methods can help avoid this
    - "Squeezing a Matrix into Half Precision, with an Application to Solving Linear Systems" [Higham, Pranesh, Zounon, 2019]

• Convergence rate of GMRES?

- Convergence rate of GMRES?
  - If A is ill conditioned and LU factorization is performed in very low precision, it can be a poor preconditioner
    - e.g., if (normal)  $\tilde{A}$  still has cluster of eigenvalues near origin, GMRES can stagnate until  $n^{\text{th}}$  iteration, regardless of  $\kappa_{\infty}(A)$  [Liesen and Tichý, 2004]
  - Potential remedies: deflation, Krylov subspace recycling [C., Oktay, 2022], using additional preconditioner

# Performance Results (MAGMA)

• [Haidar, Tomov, Dongarra, Higham, 2018]



(b) Matrix of type 4: clustered singular values,  $\sigma_i = (1, \dots, 1, \frac{1}{cond})$ .

- Convergence rate of GMRES?
  - If A is ill conditioned and LU factorization is performed in very low precision, it can be a poor preconditioner
    - e.g., if (normal)  $\tilde{A}$  still has cluster of eigenvalues near origin, GMRES can stagnate until  $n^{\text{th}}$  iteration, regardless of  $\kappa_{\infty}(A)$  [Liesen and Tichý, 2004]
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  - Potential remedies: deflation, Krylov subspace recycling [C., Oktay, 2022], using additional preconditioner
- Depending on conditioning of A, applying  $\tilde{A}$  to a vector must be done accurately (precision  $u^2$ ) in each GMRES iteration
  - Recent development of 5-precision GMRES-IR algorithm [Amestoy et al., 2021]
    - For GMRES entirely in precision **u**,

$$\kappa_{\infty}(A) \leq u^{-1/2} u_{f}^{-1} \rightarrow \kappa_{\infty}(A) \leq u^{-1/3} u_{f}^{-2/3}$$

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  - If A is ill conditioned and LU factorization is performed in very low precision, it can be a poor preconditioner
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$$\kappa_{\infty}(A) \leq u^{-1/2} u_f^{-1} \to \kappa_{\infty}(A) \leq u^{-1/3} u_f^{-2/3}$$

- Why GMRES?
  - Theoretical purposes: existing analysis and proof of backward stability [Paige, Rozložník, Strakoš, 2006]
  - In practice, use any solver you want!

# **GMRES-IR** in Libraries and Applications

• MAGMA: Dense linear algebra routines for heterogeneous/hybrid architectures

<b>D</b> m	agma / src / dxgesv_gmres_gpu.cpp
128	
129	DSGESV or DHGESV expert interface.
130	It computes the solution to a real system of linear equations
131	$A * X = B$ , $A^{**T} * X = B$ , or $A^{**H} * X = B$ ,
132	where A is an N-by-N matrix and X and B are N-by-NRHS matrices.
133	the accomodate the Single Precision DSGESV and the Half precision dhgesv API.
134	precision and iterative refinement solver are specified by facto_type, solver_type.
135	For other API parameter please refer to the corresponding dsgesv or dhgesv.

• NVIDIA's cuSOLVER Library

#### 2.2.1.6. cusolverIRSRefinement\_t

The cusolverIRSRefinement\_t type indicates which solver type would be used for the specific cusolver function. Most of our experimentation shows that CUSOLVER\_IRS\_REFINE\_GMRES is the best option.

CUSOLVER_IRS_REFINE_GMRES GMRES (Generalized Minimal Residual) based iterative refinement solver. In recent study, the GMRES method has drawn the scientific community attention for its ability to be used as refinement solver that outperforms the classical iterative refinement method. based on our experimentation, we recommend this setting.

 In production codes: FK6D/ASGarD code (Oak Ridge National Lab, USA) for tokomak containment problem

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck,1967): Least squares problem can be written as a linear system with square matrix of size (m + n):

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$$

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- Refinement proceeds as follows:
- 1. Compute "residuals"

$$\begin{bmatrix} f_i \\ g_i \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} - \begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r_i \\ x_i \end{bmatrix} = \begin{bmatrix} b - r_i - Ax_i \\ -A^T r_i \end{bmatrix}$$

2. Solve for corrections

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix} = \begin{bmatrix} f_i \\ g_i \end{bmatrix}$$

3. Update "solution":

$$\begin{bmatrix} r_{i+1} \\ x_{i+1} \end{bmatrix} = \begin{bmatrix} r_i \\ x_i \end{bmatrix} + \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix}$$

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- 1. Compute "residuals"

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$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix} = \begin{bmatrix} f_i \\ g_i \end{bmatrix}$$

3. Update "solution":

$$\begin{bmatrix} r_{i+1} \\ x_{i+1} \end{bmatrix} = \begin{bmatrix} r_i \\ x_i \end{bmatrix} + \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix}$$

 $\tilde{A}d_i = \tilde{r}_i$ 

 $\tilde{x}_{i+1} = \tilde{x}_i + d_i$ 

- For inconsistent systems, must simultaneously refine both solution and residual
- (Björck,1967): Least squares problem can be written as a linear system with square matrix of size (m + n):

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} \qquad \tilde{A}\tilde{x} = \tilde{b}$$

- Refinement proceeds as follows:
- 1. Compute "residuals"

$$\begin{bmatrix} f_i \\ g_i \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix} - \begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} r_i \\ x_i \end{bmatrix} = \begin{bmatrix} b - r_i - Ax_i \\ -A^T r_i \end{bmatrix}$$

2. Solve for corrections

$$\begin{bmatrix} I & A \\ A^T & 0 \end{bmatrix} \begin{bmatrix} \Delta r_i \\ \Delta x_i \end{bmatrix} = \begin{bmatrix} f_i \\ g_i \end{bmatrix}$$

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Results for 3-precision IR for linear systems also applies to least squares problems!

 $\tilde{x}_{i+1} = \tilde{x}_i + d_i$ 

 $\tilde{A}d_i = \tilde{r}_i$ 

 $\tilde{r}_i = \tilde{b} - \tilde{A}\tilde{x}_i$ 

See [C., Higham, Pranesh, 2020]

# GMRES-IR with Inexact Preconditioners

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- [Amestoy et al., 2022]
  - Analysis of block low-rank (BLR) LU within GMRES-IR
  - Analysis of use of static pivoting in LU within GMRES-IR

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- [Amestoy et al., 2022]
  - Analysis of block low-rank (BLR) LU within GMRES-IR
  - Analysis of use of **static pivoting** in LU within GMRES-IR
- [C., Khan, 2022]
  - Analysis of sparse approximate inverse (SPAI) preconditioners within GMRES-IR

## **SPAI** Preconditioners

Goal: Construct sparse matrix  $M \approx A^{-1}$  (for survey see [Benzi, 2002])

Approach of [Grote, Huckle, 1997]: Construct columns  $m_k$  of M dynamically

```
Given matrix A, initial sparsity structure J, and tolerance \varepsilon
For each column k:
Compute QR factorization of submatrix of A defined by J
Use QR factorization to solve \min_{m_k} ||e_k - Am_k||_2
If ||r_k||_2 = ||e_k - Am_k||_2 \le \varepsilon
break;
Else
add select nonzeros to J, repeat.
```

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Else
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```

Benefits: Highly parallelizable

But construction can still be costly, esp. for large-scale problems [Gao, Chen, He, 2021], [Chao, 2001], [Benzi, Tůma, 1999], [He, Yin, Gao, 2020]

## SPAI Preconditioners in Low Precision

What is the effect of using low precision in SPAI construction?

Notes and assumptions:

- We will assume that the SPAI construction is performed in some precision  $u_f$
- We will denote quantities computed in finite precision with hats
- In our application, we want a left preconditioner, so we will run the algorithm on  $A^T$  and set  $M \leftarrow M^T$ .
- We will assume that the QR factorization of the submatrix of  $A^T$  is computed fully using HouseholderQR/TSQR

Two interesting questions:

1. Assuming we impose no maximum sparsity pattern on  $\widehat{M}$ , under what constraint on  $\boldsymbol{u}_{f}$  can we guarantee that  $\|\hat{r}_{k}\|_{2} \leq \boldsymbol{\varepsilon}$ , with  $\hat{r}_{k} = f l_{\boldsymbol{u}_{f}}(e_{k} - A^{T} \widehat{m}_{k}^{T})$  for the computed  $\widehat{m}_{k}^{T}$ ?

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- 2. Assume that when M is computed in exact arithmetic, we quit as soon as  $||r_k|| \le \varepsilon$ . For  $\widehat{M}$  computed in precision  $u_f$  with the same sparsity pattern as M, what is  $||e_k A^T \widehat{m}_k^T||_2$ ?

Using standard rounding error analysis and perturbation results for LS problems, we have

$$\|\hat{r}_{k}\|_{2} \leq n^{3} \boldsymbol{u}_{f} \||e_{k}| + |A^{T}| \|\hat{m}_{k}^{T}\|\|_{2}.$$

So in order to guarantee we eventually reach a solution with  $\|\hat{r}_k\|_2 \leq \pmb{\varepsilon},$  we need

 $n^{3} \boldsymbol{u_{f}} \| |\boldsymbol{e}_{k}| + |\boldsymbol{A}^{T}| \left\| \widehat{\boldsymbol{m}}_{k}^{T} \right\|_{2} \leq \boldsymbol{\varepsilon}.$ 

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 $\rightarrow$  problem must not be so ill-conditioned WRT  $u_f$  that we incur an error greater than  $\epsilon$  just computing the residual

Can turn this into the looser but more descriptive a priori bound:

 $\operatorname{cond}_2(A^T) \leq \varepsilon u_f^{-1},$ 

where  $\operatorname{cond}_2(A^T) = |||A^{-T}||A^T|||_2$ .

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Another view: with a given matrix A and a given precision  $u_f$ , one must set  $\varepsilon$  such that

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Confirms intuition: The more approximate the inverse, the lower the precision we can use.

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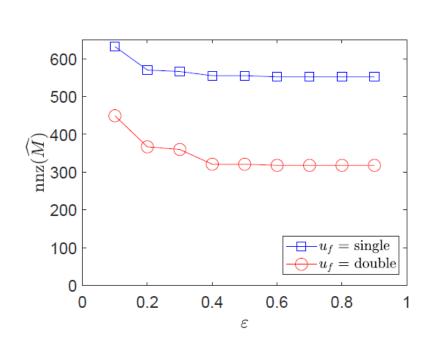
Confirms intuition: The more approximate the inverse, the lower the precision we can use.

Resulting bounds for  $\widehat{M}$ :

$$\left\|I - A^T \widehat{M}^T\right\|_F \le 2\sqrt{n}\varepsilon, \qquad \left\|I - \widehat{M}A\right\|_{\infty} \le 2n\varepsilon$$

#### Size of SPAI Preconditioner in Low Precision

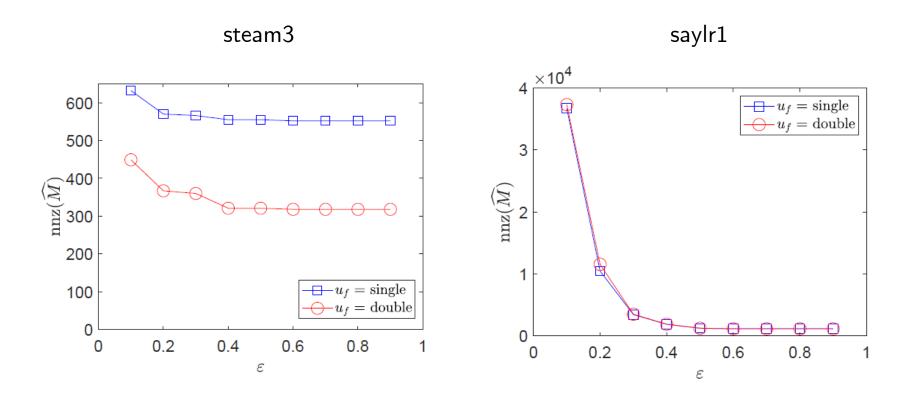
How does precision used affect the number of nonzeros in  $\widehat{M}$ ?



steam3

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# Second Question

Assume that when M is computed in exact arithmetic, we quit as soon as  $||r_k|| \leq \varepsilon$ . For  $\widehat{M}$  computed in precision  $u_f$  with the same sparsity pattern as M, what is  $||e_k - A^T \widehat{m}_k^T||_2$ ?

### Second Question

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In this case, we obtain the bound

$$\left\|I - \widehat{M}A\right\|_{\infty} \leq n\left(\boldsymbol{\varepsilon} + n^{7/2}\boldsymbol{u_f}\kappa_{\infty}(A)\right).$$

 $\rightarrow$  If  $\kappa_{\infty}(A) \gg \varepsilon u_{f}^{-1}$ , then computed  $\widehat{M}$  with same sparsity structure as M can be of much lower quality.

#### SPAI-GMRES-IR

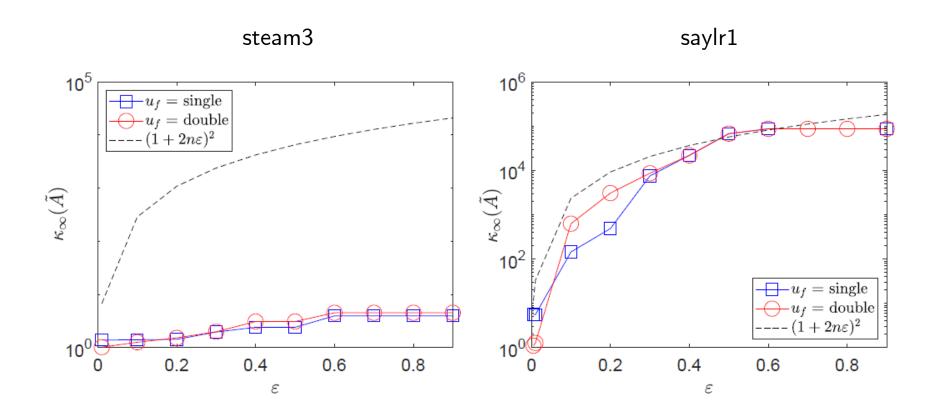
#### **SPAI-GMRES-IR**

To compute the updates  $d_i$ , apply GMRES to  $\widehat{M}Ad_i = \widehat{M}r_i$ 

Solve  $\widehat{M}Ax_0 = \widehat{M}b$ for i = 0: maxit  $r_i = b - Ax_i$ Solve  $Ad_i = r_i$  via GMRES on  $\widehat{M}Ad_i = \widehat{M}r_i$  $x_{i+1} = x_i + d_i$ 

Using  $\widehat{M}$  computed in precision  $\boldsymbol{u}_{f}$ , for the preconditioned system  $\widetilde{A} = \widehat{M}A$ ,

 $\kappa_{\infty}(\tilde{A}) \lesssim (1+2n\varepsilon)^2.$ 



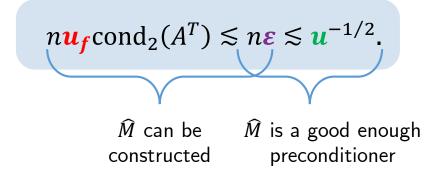
To guarantee that both SPAI construction will complete and the GMRESbased iterative refinement scheme will converge, we must have roughly

 $n \mathbf{u}_{\mathbf{f}} \operatorname{cond}_2(A^T) \leq n \boldsymbol{\varepsilon} \leq \boldsymbol{u}^{-1/2}.$ 

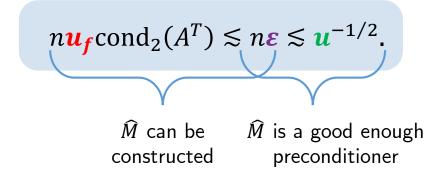
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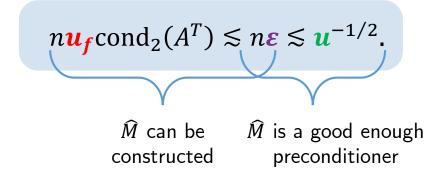


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If  $\varepsilon$  satisfies these constraints, then the constraints on condition number for forward and backward errors to converge are the same as for GMRES-IR with full LU factorization.

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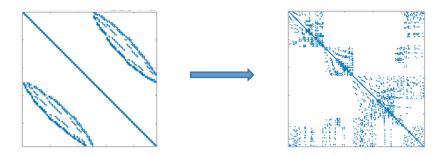


If  $\varepsilon$  satisfies these constraints, then the constraints on condition number for forward and backward errors to converge are the same as for GMRES-IR with full LU factorization.

Compared to GMRES-IR with full LU factorization, in general expect slower convergence, but much sparser preconditioner.

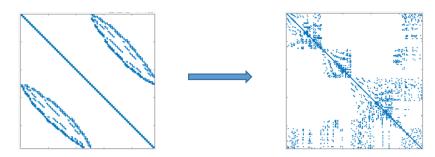
# SPAI-GMRES-IR Example

Matrix: steam1, n = 240, nnz = 2,248,  $\kappa_{\infty}(A) = 3 \cdot 10^7$ , cond $(A^T) = 3 \cdot 10^3$ 



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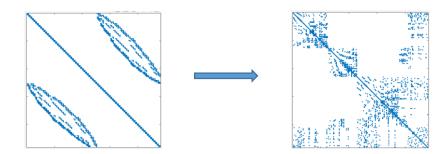


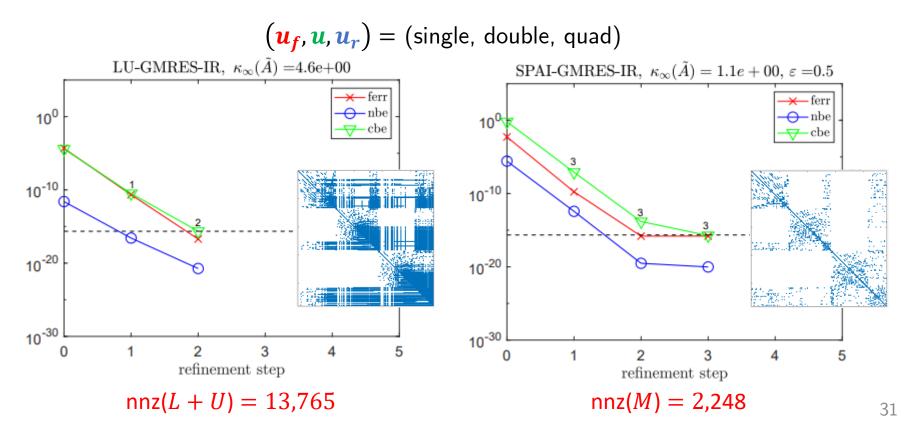
 $(\mathbf{u_f}, \mathbf{u}, \mathbf{u_r}) = (\text{single, double, quad})$ LU-GMRES-IR,  $\kappa_{\infty}(\tilde{A}) = 4.6e + 00$ <del>×</del>ferr 10<sup>0</sup> nbe cbe 10<sup>-10</sup> - 1 C 10<sup>-20</sup> 10<sup>-30</sup> 2 3 1 4 5 0 refinement step nnz(L + U) = 13,765

31

#### SPAI-GMRES-IR Example

Matrix: steam1, n = 240, nnz = 2,248,  $\kappa_{\infty}(A) = 3 \cdot 10^7$ , cond $(A^T) = 3 \cdot 10^3$ 





Is there a point in using precision higher than that dictated by  $u_f \operatorname{cond}_2(A^T) \leq \varepsilon$ ? Matrix: bfwa782, n = 782, nnz = 7514,  $\kappa_{\infty}(A) = 7 \cdot 10^3$ ,  $\operatorname{cond}(A^T) = 1 \cdot 10^3$ 

Preconditioner	$\kappa_\infty( ilde A)$	Precond. nnz	GMRES-IR steps/iteration
SPAI ( $\boldsymbol{\varepsilon} = 0.2$ )	2.1e + 02	28053	67 (31, 36)
SPAI ( $\boldsymbol{\varepsilon} = 0.5$ )	9.7 <i>e</i> + 02	7528	153 (71, 82)

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 $(\mathbf{u}_f, \mathbf{u}, \mathbf{u}_r) = (\text{single}, \text{ single}, \text{ double})$ 

Preconditioner	$\kappa_\infty( ilde A)$	Precond. nnz	GMRES-IR steps/iteration
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# Related and Current Work

- Multistage mixed precision iterative refinement
   [Oktay, C., 2021]
   If IR not converging, first try changing the solver before increasing precision
- Low-precision randomized preconditioners
   [C., Daužickaitė, 2022]

Single-pass Nyström can be run in precision  $u_p \approx \frac{\lambda_{k+1}}{\sqrt{n}\lambda_1}$  without affecting the quality of limited memory preconditioner.

• Low-precision in ILU-type preconditioners What can we prove?

- We now have a multi-precision ecosystem
- Huge opportunities for using mixed precision in matrix computations
- But also big challenges!

# Thank You!

carson@karlin.mff.cuni.cz www.karlin.mff.cuni.cz/~carson/

$$\|r_i\|_2 = \mu_i^{(2)} \|A\|_2 \|x - \hat{x}_i\|_2$$
$$x - \hat{x}_i = V \Sigma^{-1} U^T r_i = \sum_{j=1}^n \frac{(u_j^T r_i) v_j}{\sigma_j} \qquad (A = U \Sigma V^T)$$

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$$\|x - \hat{x}_{i}\|_{2}^{2} \ge \sum_{j=n+1-k}^{n} \frac{(u_{j}^{T}r_{i})^{2}}{\sigma_{j}^{2}} \ge \frac{1}{\sigma_{n+1-k}^{2}} \sum_{j=n+1-k}^{n} (u_{j}^{T}r_{i})^{2} = \frac{\|P_{k}r_{i}\|_{2}^{2}}{\sigma_{n+1-k}^{2}}$$

where  $P_k = U_k U_k^T$ ,  $U_k = [u_{n+1-k}, ..., u_n]$ 

$$\begin{aligned} \|r_{i}\|_{2} &= \mu_{i}^{(2)} \|A\|_{2} \|x - \hat{x}_{i}\|_{2} \\ x - \hat{x}_{i} &= V\Sigma^{-1}U^{T}r_{i} &= \sum_{j=1}^{n} \frac{(u_{j}^{T}r_{i})v_{j}}{\sigma_{j}} \qquad (A = U\Sigma V^{T}) \\ \|x - \hat{x}_{i}\|_{2}^{2} &\geq \sum_{j=n+1-k}^{n} \frac{(u_{j}^{T}r_{i})^{2}}{\sigma_{j}^{2}} \geq \frac{1}{\sigma_{n+1-k}^{2}} \sum_{j=n+1-k}^{n} (u_{j}^{T}r_{i})^{2} &= \frac{\|P_{k}r_{i}\|_{2}^{2}}{\sigma_{n+1-k}^{2}} \\ \text{where } P_{k} = U_{k}U_{k}^{T}, U_{k} = [u_{n+1-k}, \dots, u_{n}] \\ \mu_{i}^{(2)} &\leq \frac{\|r_{i}\|_{2}}{\|P_{k}r_{i}\|_{2}} \frac{\sigma_{n+1-k}}{\sigma_{1}} \end{aligned}$$

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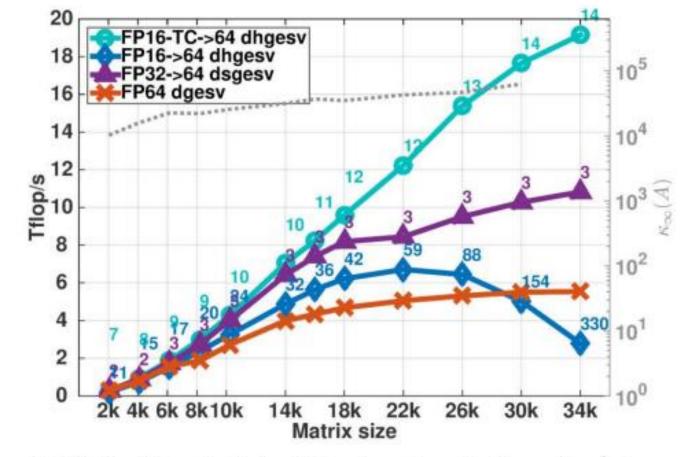
•  $\mu_i^{(2)} \ll 1$  if  $r_i$  contains significant component in span $(U_k)$  for any k s.t.  $\sigma_{n+1-k} \approx \sigma_n$ 

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- $\mu_i^{(2)} \ll 1$  if  $r_i$  contains significant component in span $(U_k)$  for any k s.t.  $\sigma_{n+1-k} \approx \sigma_n$
- In that case,  $x \hat{x}_i$  is not "typical", i.e., it contains large components in right singular vectors corresponding to small singular values of A
- Wilkinson (1977), comment in unpublished manuscript:  $\mu_i^{(2)}$  increases with *i*

# Performance Results (MAGMA)

• [Haidar, Tomov, Dongarra, Higham, 2018]



(b) Matrix of type 4: clustered singular values,  $\sigma_i = (1, \dots, 1, \frac{1}{cond})$ .

### Randomized Limited Memory Preconditioners

Let  $A \in \mathbb{R}^{n \times n}$  be a symmetric positive semidefinite matrix. Want to solve

 $(A + \mu I)x = b$ 

where  $\mu \ge 0$  is set so that  $A + \mu I$  is positive definite. Assume A has rapidly decreasing eigenvalues or cluster of large eigenvalues.

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Want to solve using PCG using **spectral limited memory preconditioner** [Gratton, Sartenaer, Tshimanga, 2011], [Tshimanga et al., 2008]:

$$P = I - UU^T + \frac{1}{\alpha + \mu}U(\Theta + \mu I)U^T$$
$$P^{-1} = I - UU^T + (\alpha + \mu)U(\Theta + \mu I)^{-1}U^T$$

where columns of  $U \in \mathbb{R}^{n \times k}$  are k approximate eigenvectors of A and  $U^T U = I$ ,  $\Theta$  is diagonal with approximations to eigenvalues of A, and  $\alpha \ge 0$ .

Used in data assimilation [Laloyaux et al., 2018], [Mogensen, Alonso Balmaseda, Weaver, 2012], [Moore et al., 2011], [Daužickaitė, Lawless, Scott, van Leeuwen, 2021]

### Randomized Nyström Approximation

Want to compute a rank-k approximation  $A \approx U \Theta U^T$  via the randomized Nyström method.

Nyström approximation:

$$A_N = (AQ)(Q^T AQ)^+ (AQ)^T$$

where Q is an  $n \times k$  sampling matrix (random projection).

#### Randomized Nyström Approximation

In the case that A is very large, matrix-matrix products with A are the bottleneck.

This motivates the single-pass version of the Nyström method.

Stabilized Single-Pass Nyström method [Tropp et al., 2017]

```
Given sym. PSD matrix A, target rank k

G = \operatorname{randn}(n, k)

[Q, \sim] = \operatorname{qr}(G, 0)

Y = AQ

Compute shift \nu; Y_{\nu} = Y + \nu Q

B = Q^T Y_{\nu}

C = \operatorname{chol}((B + B^T)/2)

Solve F = Y_{\nu}/C

[U, \Sigma, \sim] = \operatorname{svd}(F, 0)

\Theta = \max(0, \Sigma^2 - \nu I)
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```

Can we further reduce the cost of the matrix-matrix product with A by using low precision?

#### Error Bounds

$$\|A - \hat{A}_N\|_2 = \|A - A_N + A_N - \hat{A}_N\|_2 \le \|A - A_N\|_2 + \|A_N - \hat{A}_N\|_2$$

exact approximation error finite precision

error

### Error Bounds

$$\begin{split} \left\|A - \hat{A}_{N}\right\|_{2} &= \left\|A - A_{N} + A_{N} - \hat{A}_{N}\right\|_{2} \leq \left\|A - A_{N}\right\|_{2} + \left\|A_{N} - \hat{A}_{N}\right\|_{2} \\ & \text{exact} & \text{finite precision} \\ & \text{approximation} & \text{error} \\ & \text{error} \end{split}$$

$$\begin{aligned} \text{Deterministic bound [Gittens, Mahoney, 2016]:} \\ & \left\|A - A_{N}\right\|_{2} \leq \lambda_{k+1} + \left\|\sum_{2}^{1/2} U_{2}^{T} Q(U_{1}Q)^{+}\right\|_{2}^{2} \end{aligned}$$

$$\text{with } A &= \begin{bmatrix}U_{1} \ U_{2}\end{bmatrix}\begin{bmatrix}\Sigma_{1} \\ & \sum_{2}\end{bmatrix} \begin{bmatrix}U_{1} \ U_{2}\end{bmatrix}^{T}. \end{aligned}$$

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Expected value bound [Frangella, Tropp, Udell, 2021]:

$$\mathbb{E}\|A - A_N\|_2 \le \min_{2 \le p \le k-2} \left( \left(1 + \frac{2(k-p)}{p-1}\right) \lambda_{k-p+1} + \frac{2e^2k}{p^2 - 1} \sum_{j=k-p+1}^n \lambda_j \right)$$

where  $\lambda_i \geq \lambda_{i+1}$  are the eigenvalues of A.

Finite precision error:  $A_N - \hat{A}_N$ 

Assumptions:

- A is stored in precision  $u_p$  and matrix-matrix product AQ is computed in precision  $u_p$
- All other quantities stored and computed in precision  $u \ll u_p$

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The more approximate the low-rank representation, the lower the precision we can use!

## Condition Number Bounds

Let  $E = A - A_N$ ,  $\mathcal{E} = A_N - \hat{A}_N$ , and assume  $(A + \mu I)$  is SPD.

Let

$$\widehat{P}^{-1} = I - \widehat{U}\widehat{U}^T + (\widehat{\lambda}_k + \mu)\widehat{U}(\widehat{\Theta} + \mu I)^{-1}\widehat{U}^T$$

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Then

$$\max\left\{1, \frac{\hat{\lambda}_k + \mu - \|\mathcal{E}\|_2}{\mu + \lambda_{min}(A)}\right\} \le \kappa \left(\hat{P}^{-1/2}(A + \mu I)\hat{P}^{-1/2}\right) \le 1 + \frac{\hat{\lambda}_k + \|E\|_2 + 2\|\mathcal{E}\|_2}{\mu - \|\mathcal{E}\|_2}$$

where the upper bound holds if  $\mu > \|\mathcal{E}\|_2$ .

Regardless of this constraint, if A is positive definite, then

$$\kappa \left( \hat{P}^{-1/2} (A + \mu I) \hat{P}^{-1/2} \right) \leq \left( \hat{\lambda}_k + \mu + \|E\|_2 + \|\mathcal{E}\|_2 \right) \left( \frac{1}{\hat{\lambda}_k + \mu} + \frac{\|\mathcal{E}\|_2 + 1}{\lambda_{min}(A) + \mu} \right).$$

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Tropp, Udell, 2021] for exact case.

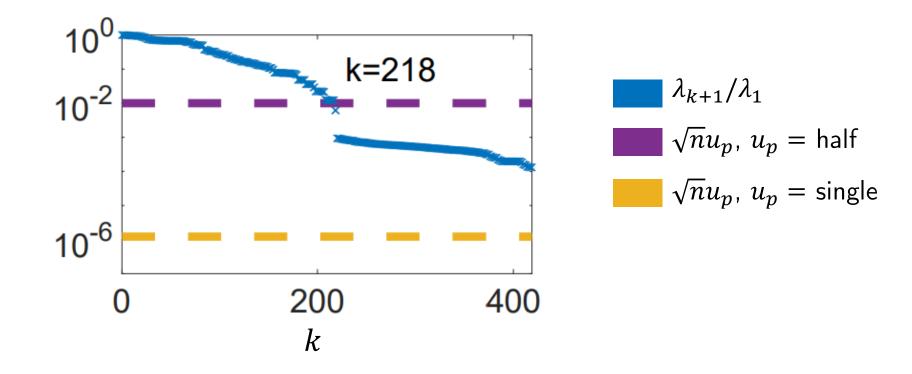
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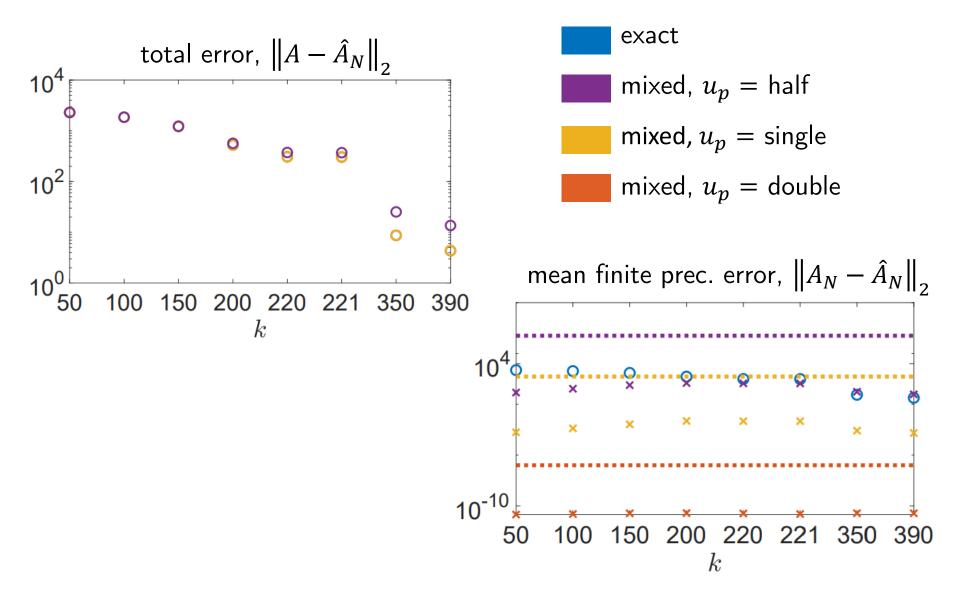
### Numerical Experiment

Matrix: bcsstm07, n = 420

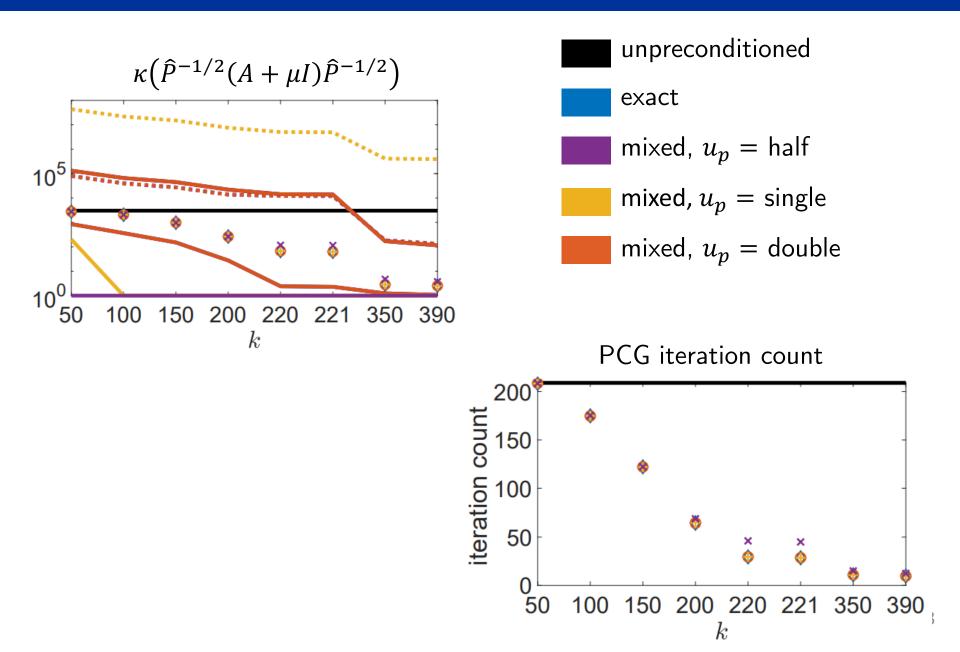


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### Numerical Experiment



- Similar to the linear system case, we can use a lower precision factorization for even more ill-conditioned problems if we **improve the effective precision of the solver**
- Again, don't want to compute an LU factorization of the augmented system
- How can we use existing QR factors to construct an effective and inexpensive preconditioner for the augmented system?
- Note that augmented system is a saddle-point system; lots of existing work (block diagonal, triangular, constraint-based, ... )

• Ex: block diagonal preconditioner ([Murphy, Golub, Wathen, 2000], [Ipsen, 2001])

$$\begin{bmatrix} \alpha I & 0 \\ 0 & \frac{1}{\alpha} \hat{R}^T \hat{R} \end{bmatrix} = \begin{bmatrix} \sqrt{\alpha} I & 0 \\ 0 & \frac{1}{\sqrt{\alpha}} \hat{R}^T \end{bmatrix} \begin{bmatrix} \sqrt{\alpha} I & 0 \\ 0 & \frac{1}{\sqrt{\alpha}} \hat{R} \end{bmatrix} \equiv M_1 M_2$$

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• Assuming QR factorization is exact,

$$M_{2}^{-1}M_{1}^{-1}\tilde{A} = \begin{bmatrix} I & \frac{1}{\alpha}A \\ \alpha \,\hat{R}^{-1}\hat{R}^{-T}A^{T} & 0 \end{bmatrix}$$

is nonsymmetric, diagonalizable, with eigenvalues  $\left\{1, \frac{1}{2}\left(1 \pm \sqrt{5}\right)\right\}$ .

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- However, condition number can still be quite large; unsuitable for proving backward stability of GMRES
- If we take split preconditioner

$$M_1^{-1}\tilde{A}M_2^{-1} = \begin{bmatrix} I & A\hat{R} \\ \hat{R}^{-T}A^T & 0 \end{bmatrix}$$

we will have a well-conditioned system

- However, split-preconditioned GMRES is not backward stable
- Potentially useful in practice, not but in theory

• One option:

$$M = \begin{bmatrix} \alpha I & \hat{Q}_1 \hat{R} \\ \hat{R}^T \hat{Q}_1^T & 0 \end{bmatrix}$$

• Then we can prove that for the left-preconditioned system,  $\kappa \big( M^{-1} \tilde{A} \big) \leq \Big( 1 + {\pmb u_f} c \; \kappa(A) \Big)^2$ 

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• So for GMRES-based LSIR,  $u_s \equiv u$ ; expect convergence of forward error when  $\kappa_{\infty}(A) < u^{-1/2} u_f^{-1}$  [C., Higham, Pranesh, SISC 2020]

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