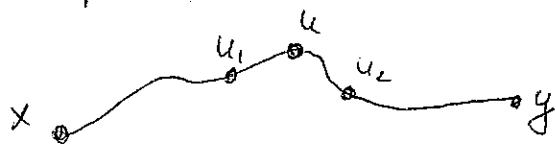


## Recommended Problems 8 -- Solutions

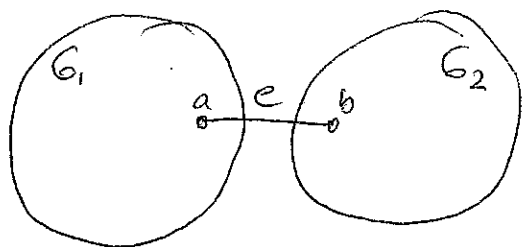
8.1 Let  $u \in V(G')$  be arbitrary. We need to show that  $G' - u$  is connected. So, let  $x, y \in V(G' - u)$  be arbitrary vertices and we want to find a  $x, y$ -path in  $G' - u$ . Since  $G$  is connected, there exists a  $x, y$ -path  $P$  in  $G$ . If  $P$  doesn't contain  $u$ , then  $P$  is a  $x, y$ -path in  $G' - u$  and we are done. Otherwise, let  $u_1$  be the last vertex before  $u$  on the path  $P$  and let  $u_2$  be the first vertex after  $u$ :



If  $u_1, u_2 \in E(G)$  then  $x \rightsquigarrow u_1 - u_2 \rightsquigarrow y$  is an  $x, y$ -path in  $G' - u$ .

If  $u_1, u_2 \notin E(G)$  then  $d_G(u_1, u_2) = 2$  therefore  $u_1, u_2$  are adjacent in  $G'$  and, again,  $x \rightsquigarrow u_1 - u_2 \rightsquigarrow y$  is an  $x, y$ -path in  $G' - u$ .

8.2. Since  $\kappa(G) = \kappa'(G)$  for 3-regular graphs, such a graph  $G$  has edge-connectivity 1, therefore it has a cut edge  $e$ . Let  $G_1, G_2$  be the components of  $G - e$



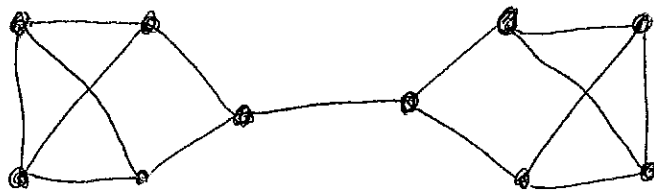
Each component has one vertex of degree 2 ( $a$  in  $G_1$ ,  $b$  in  $G_2$ ) and all the other vertices have degree 3. Then  $|V(G_1)| \geq 4$  (and  $|V(G_2)| \geq 4$ ).

But  $|V(G_1)|$  (or  $|V(G_2)|$ ) cannot have 4 vertices, because the number of vertices of odd degree must be even. Therefore  $|V(G_1)|, |V(G_2)| \geq 5$ , thus  $|V(G)| \geq 10$ .

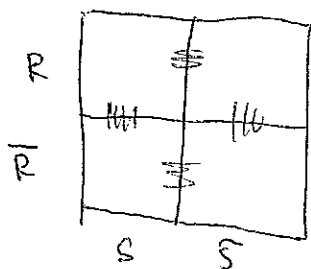
On the other hand, there exists a <sup>3-regular</sup> graph with 10 vertices and connectivity 1:

# Recommended Problems 8 - Solutions

8.2 cutd



8.3 Let  $[S, \bar{S}], [R, \bar{R}]$  be different edge-cuts.



$[S, \bar{S}]$  contains all the edges between  $S, \bar{S}$   
 $[R, \bar{R}]$  contains all the edges between  $R, \bar{R}$

$[S, \bar{S}] \Delta [R, \bar{R}]$  contains all the edges between  $S \cap R, \bar{S} \cap R$   
 and  $S \cap \bar{R}, \bar{S} \cap \bar{R}$   
 and  $S \cap \bar{R}, S \cap R$   
 and  $\bar{S} \cap R, \bar{S} \cap \bar{R}$

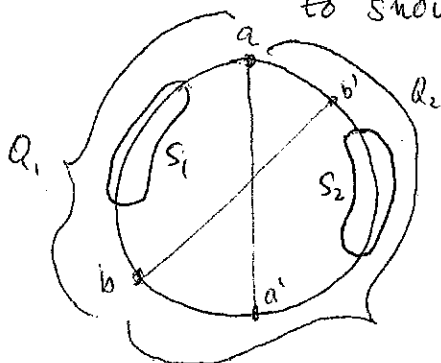
which is equal to the set of edges between  $S \cap R \cup \bar{S} \cap \bar{R}, S \cap \bar{R} \cup \bar{S} \cap R$

$$= [S \cap R \cup \bar{S} \cap \bar{R}, S \cap \bar{R} \cup \bar{S} \cap R]$$

↓ this is an edge cut (both sets are nonempty since  $[S, \bar{S}], [R, \bar{R}]$  are different)

8.4  $\kappa(G) \leq k$  because if we delete all neighbors of a vertex, the graph becomes disconnected

$\kappa(G) \geq k$  as let  $S \subseteq V(G), |S| = k-1$  and let  $a, b \in V(G) - S$ . We need to show that  $a$  is connected to  $b$  in  $V(G) - S$ .



Denote by  $Q_1$  ( $Q_2$ , resp.) the vertices between  $a, b$  on the shorter (longer, resp.) part of the circle and let  $S_1 = S \cap Q_1, S_2 = S \cap Q_2$ .

If  $S_1$  or  $S_2$  doesn't contain  $\frac{k-1}{2}$  consecutive vertices, then "we can jump across all the gaps" to get from  $a$  to  $b$

## Recommended Problems 8 - Solutions

8.4 cubed.

Therefore we can assume that  $|S_1| = |S_2| = \frac{k-1}{2}$  and both  $S_1$  and  $S_2$  are formed by consecutive vertices.

Let  $a'$  be the vertex opposite to  $a$  and  $b'$  the vertex opposite to  $b$ .

If  $S_2$  contains no vertices between  $a$  and  $b'$  we can get from  $a$  to  $b'$  and then to  $b$ .

Similarly, if  $S_2$  contains no vertices between  $b$  and  $a'$ ,  $a$  and  $b$  are connected.

Otherwise  $S_2$  contains both  $a'$  and  $b'$  which is impossible as  $d(a', b') = d(a, b) > \frac{k-1}{2}$  (otherwise  $a$  and  $b$  are adjacent).