

RECAP

arity ≥ 1

Clone on A = set of operations on A closed under forming term operations

- = set of operations on A that
 - contains all the projections π_i^n
 - is closed under composition

$$f(g_1, \dots, g_m)(\bar{a}) = f(g_1(\bar{a}), \dots, g_m(\bar{a}))$$

$\text{Clo}(A) =$ the clone of term operations of A

$\text{Clo}_n(A) =$ n-ary operations in $\text{Clo}(A)$

! $\text{Clo}_n(A) \leq A^{A^n}$ generated by π_1^n, \dots, π_n^n

because $f^{A^{A^n}}(g_1, \dots, g_m) = \underbrace{f(g_1, \dots, g_m)}_{\text{composition}}$

↑
basic operation
in A^{A^n}

THM: $\text{Clo}_n(A) \cong F_{\{A\}}(\{x_1, \dots, x_n\})$
 (= $F_{\text{HSP}(A)}(\{x_1, \dots, x_n\})$)

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~~10.8~~

11.2

COMPATIBILITY

Def. $f: A^n \rightarrow A, R \subseteq A^m$

f is compatible with R (or R invariant (compatible) under f)

if $R \subseteq (A_i f)$

Recall

$$\begin{pmatrix} a_{11} & a_{12} & & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & & a_{mn} \end{pmatrix} \begin{matrix} \xrightarrow{f} \\ \xrightarrow{f} \\ \\ \Rightarrow \end{matrix} \begin{pmatrix} \\ \\ \\ \cap \\ R \end{pmatrix}$$

$\cap R \quad \cap R \quad \cap R \quad \cap R$

1. \emptyset the set of all operations compatible with $R \subseteq A^m$ is a clone on A (see 10.2)

- some clones from the examples are of this sort

- 2. f is compatible with $\{0\} \subseteq A'$ iff ...
- f is compatible with $\{a\} \subseteq A' \forall a \in A$ iff ...
- f is \equiv $B \subseteq A' \forall B \subseteq A$ iff ...
- f is \equiv $R \subseteq \{0,1\}^2 \dots R = \equiv$
($R = \{(0,0), (0,1), (1,1)\}$)

iff ...

2. Is \wedge/\vee compatible with " $x \wedge y \rightarrow z$ " ie. $R = \{0,1\}^3 \setminus \{(1,1,0)\}$?

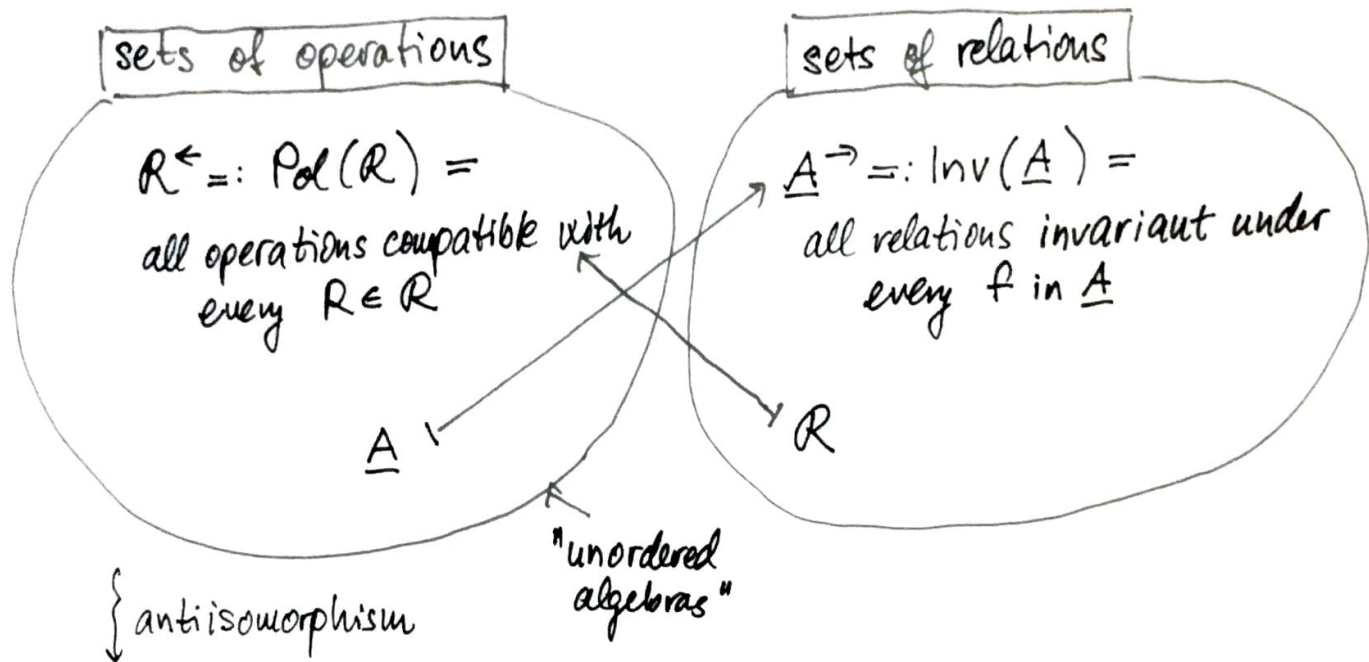
Pol-Inv: #2 Galois correspondence in UA

A - fixed

$Op :=$ all operations on A (arity ≥ 1)

$Rel :=$ all relations on A (nonempty)

relation $\mathcal{R} \subseteq Op \times Rel$ $f \in \mathcal{R}$ iff f compatible with R



(closed sets of operations, \subseteq) \leftrightarrow (closed sets of relations, \subseteq)

$= \{ Pol(\mathcal{R}); \mathcal{R} \subseteq Rel \}$ $= \{ Inv(\underline{A}); \underline{A} \subseteq Op \}$

For finite A

$\nabla \nabla$ closed sets of operations = clones

$\circ \circ$ closed sets of relations = coclones

Remark

Notation $Inv(\underline{A})$ also used for algebras

$$Inv(A; f_1, f_2, \dots) := Inv(\{f_1, f_2, \dots\})$$

Note $R \in Inv(\underline{A})$ iff R is a subpower of \underline{A}

o $\forall R \subseteq Rel \quad Pol(R)$ is a clone

Proof: $R = \{R_1, R_2, \dots\}$

$$\begin{aligned}
 Pol(R) = & \{f_i \mid f \text{ compatible with } R_1\} \\
 & \cap \{f_i \mid \text{---} \text{---} \text{---} R_2\} \\
 & \cap \dots
 \end{aligned}
 \begin{array}{l}
 \swarrow \\
 \searrow \\
 \rightarrow \text{ clones}
 \end{array}$$

→ every closed set of operations is a clone

Examples

- $Pol(\emptyset) =$ all operations
- $Pol(\text{all relations}) =$ projections Exercise
- $A = \{0, 1\}$ $Pol(\{\leq\}) =$ monotone operations
(= $Clo(\{0, 1\}; 0, 1, \wedge, \vee)$)
- $Pol(\{B_i \mid B \subseteq A\}) =$ conservative operations
- $Pol(\{\{a\} \mid a \in A\}) =$ idempotent operations

? $Pol(?) = Clo(\{0, 1\}; \rightarrow)$

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closed ^{sets} classes of operations = clones

THEOREM (Geiger '68, Bodnarcuk, Kalužnin, Kotov, Lauou '69)

For every finite algebra \underline{A} , $\text{Clo}(\underline{A}) = \text{Pol}(\text{Inv}(\underline{A}))$

Proof: \subseteq ✓ (since $\text{Pol}(\)$ is always a clone)

\supseteq • take $f \in \text{Pol}(\text{Inv}(\underline{A}))$, say n -ary

• fix a linear ordering of A^n and recall

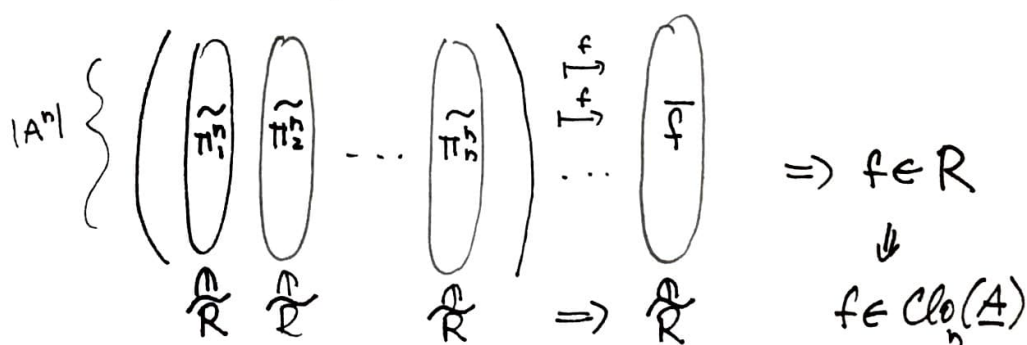
$g \in A^{A^n} \leftrightarrow \text{tuple } \tilde{g} \in A^{A^n}$

$R \subseteq A^{A^n} \leftrightarrow \text{relation } \tilde{R} = \{\tilde{g} : g \in R\} \subseteq A^{A^n}$

• ~~take~~ define $R := \text{Clo}_n(\underline{A})$

• recall $R \subseteq \underline{A}^{A^n} \Rightarrow \tilde{R} \subseteq \underline{A}^{A^n} \Rightarrow \tilde{R} \in \text{Inv}(\underline{A})$

• so f is compatible with \tilde{R} , in particular



so we have antiisomorphism

$(\text{clones}, \subseteq) \leftrightarrow (\text{closed sets of relations}, \subseteq)$

(?) Is $p(x,y) = x+y$ in $\text{Clo}(A; +, \nu)$?

THE OTHER SIDE

Def. R, R_1, \dots, R_n relations on A .

R is primitively positively definable from R_1, \dots, R_n if R can be defined by a first order formula using

- $R_1, R_2, \dots, R_n, =$
- conjunction
- existential quantification

or pp-definable

Example $R_1 \subseteq A, R_2 \subseteq A^2, R_3 \subseteq A^3$

$$R(x, y) \stackrel{\text{def}}{=} \exists u, v, w \underbrace{R_1(u) \wedge (x=v) \wedge R_2(v, y) \wedge R_3(y, x, x)}_{\text{pp-definition}}$$

What can we do

- intersection R_1, \dots, R_n m -ary $\Rightarrow R_1 \cap R_2 \cap \dots \cap R_n$ pp-definable from R_1, \dots, R_n
- permutation of coordinates R 3-ary $\Rightarrow S = \{(a, b, c); (c, a, b) \in R\}$ pp-definable from R
- introduction of dummy coordinates R 2-ary $\Rightarrow S = \{(a, b, c); (a, b) \in R\}$
- projection R ternary $\Rightarrow S = \{(b, c); \exists a (a, b, c) \in R\}$ pp-def. from R
- composition of binary relations

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11.7

Def A clone (or relational clone) on A is a set of relations on A that is closed under pp-definitions & contains \emptyset

For a set of relations \mathcal{R}

$\text{Coclo}(\mathcal{R}) :=$ the smallest clone containing \mathcal{R}
 $=$ all relations pp-definable from $\mathcal{R} + \emptyset$
 (since a relation pp-def. from relations pp-def. from \mathcal{R} is pp-def. from \mathcal{R})

$\text{Inv}(\underline{A})$ is a clone (for each set of operations \underline{A})

THEOREM For every set of relations \mathcal{R} on a finite set A
 $\text{Inv}(\text{Pol}(\mathcal{R})) = \text{Coclo}(\mathcal{R})$

so we have mutually inverse anti isomorphisms

