

RECAP

$\xi: A \rightarrow B$ clone homo.

- $\xi(\pi_i^n) = \pi_i^n$
- $\xi(f(g_1, \dots, g_m)) = \xi(f)(\xi(g_1), \dots, \xi(g_m))$

- preserves arbitrary composition

$$\xi(f(g(x_1, x_2), x_3)) = \xi(f)(\xi(g)(x_1, x_2), x_3)$$

- preserves identities
associative op. \mapsto associative op.

- $A \leq B$ if $\exists A \rightarrow B$
- $A \sim B$ if $A \leq B \leq A$
 \rightarrow ordering

- $A \leq B \Leftrightarrow B \in \text{EHSP}(A)$

P_{fin} if A, B on finite sets

$\xi: A^{\text{min}} \rightarrow B$ minion homo.

- $\xi(f(\pi_{i_1}^n, \dots, \pi_{i_m}^n)) = \xi(f)(\pi_{i_1}^n, \dots, \pi_{i_m}^n)$

- preserves minors

$$\xi(f(x_1, x_3, x_1, x_2)) = \xi(f)(x_1, x_3, x_1, x_2)$$

- preserves height 1 identities
commutative op. \mapsto commutative op.

- $A \leq^{\text{min}} B$ if $\exists A \xrightarrow{\text{min}} B$
- $A \sim^{\text{min}} B$ if $A \leq^{\text{min}} B \leq^{\text{min}} A$
 \rightarrow ordering

- $A \leq^{\text{min}} B \Leftrightarrow B \in \text{ERP}(A)$

relational counterparts \rightsquigarrow

\mathbb{F} CSP applications

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9.2

Taylor clones

Def. Clone A is Taylor if
• A is idempotent
(∀ f ∈ A f(x₁, ..., xₙ) ≈ x)

• $A \not\leq \text{Proj}$

A is Taylor if Clo(A) is

⇔ A ≠ Proj

on any domain of size ≥ 2

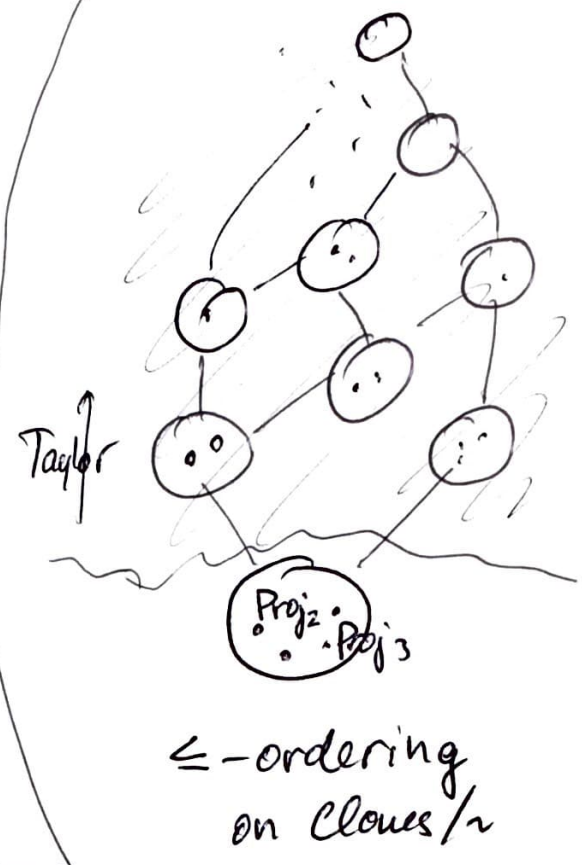
Ex A idempotent & contains commutative operation ⇒ A Taylor
Mal'cev operation

A idempotent & satisfies any nontrivial Mal'cev condition
not satisfiable in every algebra
= not satisfiable by projections

Recall A ≤ Proj_B ⇔ Proj_B ∈ HSP(A)

ie. A is not Taylor ⇔ HSP(A) contains a naked set

• ≥ 2-element algebra whose every operation is a projection



Taylor operation

Def $t: A^n \rightarrow A$ is a Taylor operation if it satisfies identities of the form

$$t(x_1, *, *, \dots) \approx t(y_1, *, *, \dots)$$

$$t(*, x_1, *, \dots) \approx t(*, y_1, *, \dots)$$

$$\dots$$

$$t(*, \dots, *, x) \approx t(*, \dots, *, y)$$

where $*$ are variables (wlog x/y)

Examples • Mal'cev operation is Taylor

$$m(\underline{x}, x, x) \approx m(y, y, x)$$

$$m(x, \underline{x}, x) \approx m(y, y, x)$$

$$m(x, x, \underline{x}) \approx m(x, y, y)$$

• majority operation is Taylor

• semilattice operation (or any commutative) is T.

• constant (non-idempotent) of any arity is Taylor

(1) \circledast A contains a Taylor operation & is idempotent

$\Rightarrow A$ Taylor

even $A \not\stackrel{\text{min}}{=} \text{Proj}$

Theorem [Taylor '70s] \mathcal{A} idempotent. @

- (i) $\mathcal{A} \neq \text{Proj}$ (\mathcal{A} is Taylor) (non-trivial identities)
- (ii) $\mathcal{A} \not\stackrel{\text{min}}{=} \text{Proj}$ (non-trivial w/ identities)
- (iii) \mathcal{A} contains a Taylor operation

Proof: ~~Def~~ (iii) \Rightarrow (ii) \Rightarrow (i) \checkmark (minor of g)

Def. $g: A^m \rightarrow A$, $\alpha: [m] \rightarrow [n]$, $g^\alpha := g(\pi_{\alpha(1)}^n, \dots, \pi_{\alpha(m)}^n)$

$\{1, 2, \dots, m\}$

i.e. $g^\alpha(x_1, \dots, x_n) = g(x_{\alpha(1)}, \dots, x_{\alpha(m)})$

Ex. $g: A^3 \rightarrow A$, $\alpha: [3] \rightarrow [2]$ $g^\alpha(x_1, x_2) = g(x_2, x_1, x_2)$

$1 \mapsto 2$
 $2 \mapsto 1$
 $3 \mapsto 2$ or $\alpha = (2, 1, 2)$

1) $f: \mathcal{A} \rightarrow \mathcal{B}$ minion homomorphism \Leftrightarrow

preserves arity & $\forall g, \alpha \dots \{f(g^\alpha) = [f(g)]^\alpha$

2) g m -ary, $\alpha: [m] \rightarrow [n]$, $\beta: [n] \rightarrow [r]$ $(g^\alpha)^\beta = g^{\beta \circ \alpha}$

3) $t: A^n \rightarrow A$ is a Taylor operation \Leftrightarrow

$\forall i \in [n] \exists \alpha, \beta: [n] \rightarrow \dots \quad t^\alpha = t^\beta$ and $\alpha(i) \neq \beta(i)$

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Star-composition

Def g m -ary, f n -ary. $g * f$ mn -ary defined by
 $g * f(x_1, x_2, \dots, x_{mn}) := g(f(x_1, \dots, x_n), f(x_{n+1}, \dots, x_{2n}), \dots, f(\dots, x_{mn}))$



(4) $g(f^{\alpha_1}, f^{\alpha_2}, \dots, f^{\alpha_m}) = (g * f)^{\alpha_1 \hat{\ } \alpha_2 \hat{\ } \dots \hat{\ } \alpha_m}$
concatenation

Lemma 1 T ... finite set of idempotent operations on A .
Then $\exists F \in \text{Clo}(T) \forall f \in T F^\alpha = f$ for some α

Proof: • Say $T = \{f, g\}$

• $F := g * f$

• $F(x_1, x_1, \dots, x_1, x_2, x_2, \dots, x_2, \dots, x_m, \dots, x_m) = g(x_1, \dots, x_n)$
(ie $F^{(1,1, \dots, 1, 2, 2, \dots, 2, \dots, m, \dots, m)} = g$)

• $F(x_1, x_2, \dots, x_n, x_1, x_2, \dots, x_n, \dots, x_1, x_2, \dots, x_n)$
 $= f(x_1, x_2, \dots, x_n)$
(ie $F^{(1, 2, 3, \dots, m, 1, 2, \dots, m, \dots)} = f$)

in general $T = \{f_1, f_2, \dots, f_n\} \quad F := f_1 * f_2 * \dots * f_n$

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(i) \Rightarrow (ii) $A \neq \text{Proj} \Rightarrow A \not\stackrel{\text{min}}{=} \text{Proj}$

Will prove more: $\forall f: A \xrightarrow{\text{min}} \text{Proj}$ is a ^{clone} homomorphism

• Take $f: A \xrightarrow{\text{min}} \text{Proj}$

• Lemma 2 $t \in a_m, s \in a_n \quad f(\overset{h}{t * s}) = f(t) * f(s)$

• say $f(t) = \pi_{\overset{m}{i}}$, $f(s) = \pi_{\overset{n}{j}}$, $f(h) = \pi_{\overset{mn}{k}}$
index projections from 0 (π_0^m is the 1st proj.)

• RHS = π_{ni+j}^{mn}

• $[f(h)]^{(1,2,\dots,n), (1,2,\dots,n,\dots)} \stackrel{\text{O1}}{=} f(h^{(1,2,\dots,n), (1,2,\dots,n,\dots)})$

= $f(hs) \Rightarrow k \bmod n = j$

• $[f(h)]^{(1,1,\dots,1,2,2,\dots)} \dots \Rightarrow k \text{ div } n = i \quad \left. \begin{matrix} k = \\ ni + j \end{matrix} \right\}$

• want $f \left(\underset{a}{\begin{matrix} g \\ \swarrow \quad \searrow \\ f_1 \quad \dots \quad f_m \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ x_1 \quad \dots \quad x_n \quad \dots \quad x_1 \quad \dots \quad x_n \end{matrix}} \right) = \begin{matrix} f(g) \\ \swarrow \quad \searrow \\ f(f_1) \quad \dots \quad f(f_m) \\ \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ x_1 \quad \dots \quad x_n \quad \dots \quad x_1 \quad \dots \quad x_n \end{matrix}$

• take F from Lemma 1 such that $F^{\alpha_i} = f_i$ ($\alpha_i: [\text{only } F] \rightarrow [n]$)

LHS = $f(g(f_1, \dots, f_m)) = f(g(F^{\alpha_1}, \dots, F^{\alpha_m})) \stackrel{\text{O4}}{=} f(g * F^{\alpha_1 \wedge \alpha_2 \wedge \dots \wedge \alpha_m})$

$\stackrel{\text{O1}}{=} f(g * F)^{\alpha_1 \wedge \alpha_2 \wedge \dots \wedge \alpha_m} \stackrel{\text{Lemma 2}}{=} (f(g * F))^{\alpha_1 \wedge \alpha_2 \wedge \dots \wedge \alpha_m} = (f(g) * f(F))^{\alpha_1 \wedge \alpha_2 \wedge \dots \wedge \alpha_m}$

$\stackrel{\text{O4}}{=} f(g)((f(F))^{\alpha_1}, \dots, (f(F))^{\alpha_m}) \stackrel{\text{O1}}{=} f(g)(f(F^{\alpha_1}), \dots, f(F^{\alpha_m}))$

= $f(g)(f(f_1), \dots, f(f_m)) = \text{RHS}$

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we need to select "important coordinate" for each $t \in A$ so that the selection goes well with minors

(ii) \Rightarrow (iii) no Taylor $\Rightarrow A \xrightarrow{\text{min}} \text{Proj}$

$$\xi(t^\alpha) = [\xi(t)]^\alpha$$

Assume A doesn't contain Taylor operation

Step 1 $\forall t \in A \exists \xi: A \rightarrow \text{Proj}$ $\forall \alpha \dots \xi(t^\alpha) = [\xi(t)]^\alpha$
arity preserving

- t not Taylor $\stackrel{\text{ö3}}{\Rightarrow} \exists i \forall \alpha, \beta \ t^\alpha = t^\beta \Rightarrow \alpha(i) = \beta(i)$
- define $\xi(t^\alpha) := \prod_{\alpha(i)} \xi(t) = \pi_{\alpha(i)} \quad (\Rightarrow \xi(t) = \pi_i)$
 - makes sense
 - $\xi(t^\alpha) = \pi_{\alpha(i)} = [\xi(t)]^\alpha$

Step 2 $\forall T \subseteq_{\text{fin}} A \exists \xi \forall t \in T \forall \alpha \dots \xi(t^\alpha) = [\xi(t)]^\alpha$

- Lemma 1 - $\exists F \forall t \in T \ t = F^\beta$ for some β
- $\xi :=$ from Step 1 for F
- $\xi(t^\alpha) = \xi((F^\beta)^\alpha) \stackrel{\text{ö2}}{=} \xi(F^{\alpha\beta}) \stackrel{\text{step 1}}{=} [\xi(F)]^{\alpha\beta} = \stackrel{\text{ö2}}{=} ([\xi(F)]^\beta)^\alpha \stackrel{\text{step 1}}{=} (\xi(F^\beta))^\alpha = (\xi(t))^\alpha$

Step 3 \exists minion homo $\xi: A \rightarrow \text{Proj}$ - compactness argument

any symbols: $\pi_i^n, \xi(t)$

sentences: $n \in \mathbb{N} \bullet \pi_1^n \neq \pi_2^n \neq \dots$

$\forall t \in A \text{ arity } n \bullet \xi(t) = \pi_1^n \vee \xi(t) = \pi_2^n \vee \dots$

$\forall t \forall \alpha \ s := t^\alpha \ \xi(s) = (\xi(t))^\alpha \bullet \xi(t) = \pi_1^n \Rightarrow \xi(s) = \pi_{\alpha(i)}^n, \dots$

\forall finite subset has a model (Step 2)

$\Rightarrow \exists$ model

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9.8

We proved: A idempotent: \mathbb{P}

- $A \not\leq \text{Proj}$ } in general
- $A \not\leq^{\text{min}} \text{Proj}$ }
- A contains a Taylor op.
- ~~$A \leq \text{Proj}$~~ $\text{Proj}_2 \in \text{HSP}(A)$

Corollary A finite (not necessarily idempotent)

- (i) $A \not\leq^{\text{min}} \text{Proj}$ (this is what interests us for CSP)
- (ii) A contains a Taylor operation

Proof: (ii) \rightarrow (i) \checkmark

(i) \rightarrow (ii) • $\exists B$ idempotent $A \sim B$

• $B \not\leq^{\text{min}} \text{Proj} \Rightarrow B$ contains a Taylor op.

• $B \leq A \Rightarrow A$ contains a Taylor op.
(since Taylor identities are height 1)

But • for finite (non-idempotent) $A \leq^{\text{min}} \text{Proj} \not\Rightarrow A \leq \text{Proj}$

• for infinite (idempotent) $A \not\leq^{\text{min}} \text{Proj} \not\Rightarrow A$ contains a Taylor

Viewpoint via important ^{variables /} coordinates

$\{ : A \leq^{\text{min}} \text{Proj} \Leftrightarrow \exists I: A \rightarrow N$ such that

• $I(f) \in \{1, 2, \dots, \text{arity } f\}$ ($\forall f \in A$)

• $I(f^x) = x(I(f))$

($\forall f \in A \forall x: [\text{arity } f] \rightarrow \mathbb{N}$)

Theorem

A idempotent finite \mathcal{A}

~~$\mathcal{A} \in \text{Proj}_2$~~ (Taylor)

• $\mathcal{A} \notin \text{Proj}_2$

practical

• $\text{Proj}_2 \in \text{HSP}(\mathcal{A})$

• $\text{Proj}_2 \in \text{HS}(\mathcal{A})$ [Bulatov '00s]

• \mathcal{A} contains a Taylor operation [Taylor '70s]

• \mathcal{A} contains a weak near-unanimity operation
 $w(y, x, x, \dots, x) \approx w(x, y, x, \dots, x) \approx \dots \approx w(x, x, \dots, x, y)$
(arity ≥ 2) [McKenzie, Maroti '00s]

• \mathcal{A} contains a p -ary cyclic operation for every prime $p > |\mathcal{A}|$ [Barto, Kozik '10s]
 $c(x_1, \dots, x_p) \approx c(x_2, \dots, x_p, x_1)$

• \mathcal{A} contains a 6-ary Siggers operation [Siggers '10s]
 $s(x, y, x, z, y, z) \approx s(y, x, z, x, z, y)$

we will prove this

• \mathcal{A} contains a 4-ary Siggers operation

$s(r, a, r, e) \approx s(a, r, e, a)$ [Keames, Markovic, McKenzie]

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9.10

Theorem A idempotent $\Downarrow @$

- $A \notin \text{Proj}$
- $A \notin^{\text{min}} \text{Proj}$
- A has a Taylor operation

• A has an Olšák operation [Olšák'10s]

$$\begin{aligned} & o(x_1 x_1 y_1 \quad y_1 y_1 x_1) \\ \approx & o(x_1 y_1 x_1 \quad y_1 x_1 y_1) \\ \approx & o(y_1 x_1 x_1 \quad x_1 y_1 y_1) \end{aligned}$$

This theorem & some items in the previous one
uses the concept of absorbing subuniverses

\swarrow
the rest of the course