

UA2 6.1

RECAP

$\Sigma \dots$  finite signature

$E$  is a base of identities in  $Id(V)$

$\mathcal{V}$  is finitely based if

$\exists$  finite  $E$  such that  $\mathcal{V} \models s \approx t$  iff  $E \models s \approx t$

$\Leftrightarrow$  finitely axiomatizable

A finite @

- $HSP(A)$  finitely based
- $\exists_n$   $\{n\text{-variable identities}\}$  is a base

$\exists$  finite algebras A that are not finitely based

TODAY

A finite basis result

- does not cover a "large" class of varieties
- its successors do

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**THEOREM**

[McKenzie '78]

Suppose  $\mathcal{V}$  has

- (1) definable principal congruences (DPC) and
- (2) only finitely many SIs (up to isomorphism), all of them finite.

Then  $\mathcal{V}$  is finitely based

(2) satisfied if eg.  $\mathcal{V} = \text{HSP}(\underline{A})$  and  $\mathcal{V}$  is CD  
(Jónsson's lemma  $\rightarrow$  all SIs are in  $\text{HS}(\underline{A})$ )  
e.g.  $\underline{A}$  finite lattice

(1) "almost never" satisfied but weaker conditions  
"quite often" (proofs more difficult)

recall • lattices have majority term  
 $(x_1 y) \vee (x_1 z) \vee (y_1 z)$

•  $\mathcal{V}$  has a majority term  
 $\Rightarrow \mathcal{V}$  is CD (practicals last semester)

also recall: varieties determined by SIs

VA2 [6.3]

**Def.**  $\mathcal{V}$  has DPC (definable principal congruences) if  
 $\exists$  formula  $\varphi(x, y, a, b)$  equivalent in  $\mathcal{V}$  to  
 $\underbrace{''(x, y) \in Cg(a, b)''}_{\text{free variables}}$ , i.e.

$$\forall \underline{A} \in \mathcal{V} \quad \forall x, y, a, b \in A \quad (x, y) \in Cg_{\underline{A}}(a, b) \Leftrightarrow \underline{A} \models \varphi(x, y, a, b)$$

$$\Leftrightarrow \forall \underline{A} \in \mathcal{V} \quad \forall a, b \in A \quad \varphi - Cg^{\underline{A}}(a, b) = Cg^{\underline{A}}(a, b)$$

where  $\varphi - Cg^{\underline{A}}(a, b) := \{ (x, y) \in A^2; \underline{A} \models \varphi(x, y, a, b) \}$

$\nearrow$  "what  $\varphi$  thinks is  $Cg(a, b)$ "

**Examples**

- $\mathcal{V}$  = commutative rings has DPC  $\varphi \equiv \exists z \ x - y = z \cdot (a - b)$
- $\mathcal{V}$  = abelian groups of exponent 1000 has DPC  
 $\varphi: (x - y = 0) \vee (x - y = a - b) \vee (x - y = b - a) \vee (x - y = 2(a - b)) \vee \dots$   
 $\underbrace{\hspace{10em}}_{\text{"the } x - y \text{ is in the cyclic subgroup gen. by. } a - b \text{"}}$
- $\mathcal{V}$  = distributive lattices } have DPC  
 $\mathcal{V}$  = semilattices }
- $\mathcal{V}$  = abelian groups does not have DPC

$\mathcal{U}$  = abelian groups does not have DPC

$\varphi_n: (x-y=0) \vee (x-y=a-b) \vee \dots \vee (x-y=n(a-b)) \vee (x-y=n(b-a))$

clearly  $\forall A \in \mathcal{U} \forall x, y, a, b (x=y) \in C_{\mathcal{U}}(a, b)$  iff  $\varphi_1(x, y, a, b) \vee \varphi_2(x, y, a, b) \vee \dots$

assume  $\varphi(x, y, a, b)$  defines principal congruences

so  $\varphi(x, y, a, b) \leftrightarrow \underbrace{\varphi_1(x, y, a, b) \vee \varphi_2(x, y, a, b) \vee \dots}_{\text{infinite disjunction}}$   
within  $\mathcal{U}$

Then  $\varphi(x, y, a, b) \leftrightarrow$  finite subdisjunction which is false (see  $\mathbb{Z}$ )

Why?  $\rightarrow$

$\Sigma = \{+, -, 0, \underbrace{c, d, e, f}_{\text{new 0-ary symbols}}\}$

sentences  $\left. \begin{array}{l} \text{ab. group axioms} \\ \varphi(c, d, e, f) \\ \neg \varphi_1(c, d, e, f) \\ \neg \varphi_2(c, d, e, f) \\ \vdots \end{array} \right\} \text{have no model}$

compactness for logic  $\uparrow$

$\Rightarrow$  some finite set of sentences has no model

$\Rightarrow \left. \begin{array}{l} \text{ab. group axioms} \\ \varphi(c, d, e, f) \\ \neg \varphi_{n_1}(c, d, e, f) \\ \vdots \\ \neg \varphi_{n_k}(c, d, e, f) \end{array} \right\} \text{no model} \Leftrightarrow \begin{array}{l} \varphi(x, y, a, b) \rightarrow \\ \varphi_{n_1}(x, y, a, b) \vee \\ \vdots \\ \vee \varphi_{n_k}(x, y, a, b) \end{array}$

We will need "nice" formulas defining principal congruences

**Def**  $\varphi(x, y, a, b)$  is conservative if  $\forall B$  in the signature

$$\forall a, b \in B \quad \varphi\text{-}Cg_B(a, b) \subseteq Cg_B(a, b)$$

i.e. " $\varphi(x, y, a, b) \Rightarrow (x, y) \in Cg(a, b)$ "

Examples

- $\mathcal{V}$  = commutative rings

- $\varphi = \exists z \quad x - y = z \cdot (a - b)$  is NOT conservative  
(take  $B$  such that  $\forall u, v \in B \quad u - v = u, u \cdot v = u$ )

- $\varphi = \exists z, r \quad (x = r \cdot a + z) \wedge (y = r \cdot b + z)$  is conservative  
(and defines principal congruences in  $\mathcal{V}$  like  $\sim$ )

- more complex conservative formula ( $\Sigma = \{+, *, \cdot\}$ )

$$\begin{aligned} \exists z_1, z_2, z_3 \quad & (x = z_1 \cdot a + z_2) \\ & \wedge (z_1 \cdot b + z_2 = (z_3 * b) * z_4) \\ & \wedge_{z_4} ((z_3 * a) * z_4 = y) \end{aligned}$$

i) disjunction of conservative formulas is conservative

ii) For every signature " $(x, y) \in Cg(a, b)$ " is equivalent to infinite disjunction of conservative formulas

why?  $Cg^+(a, b)$  is the transitive closure of the subalgebra of  $A^2$  generated by the equivalence  $(a, b) \cdot 1 \cdot 1 \cdot 1 \cdot 1 \dots$

- this can be written (Exercise!) as infinite disjunction of conservative formulas (like see  $\checkmark$ )

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$\rightsquigarrow \mathcal{V}$  has DPC, then  $(x,y) \in C_g(a,b)$   
is equivalent to  $\varphi(x,y,a,b)$  (within  $\mathcal{V}$ )  
and this is equivalent to infinite  $\vee$  of conservative  
formulas

$\Rightarrow (x,y) \in C_g(a,b)$  is equivalent to a conservative formula

More generally:

Theorem Assume  $\mathcal{V}$  is axiomatizable and  
 $\varphi(x_1, \dots, x_n) \leftrightarrow \underbrace{\varphi_1(x_1, \dots, x_n) \vee \varphi_2(x_1, \dots, x_n) \vee \dots}_{\text{infinite disjunction}} \quad (\text{in } \mathcal{V})$

Then  $\varphi(x_1, \dots, x_n) \leftrightarrow$  finite subdisjunction (in  $\mathcal{V}$ )

Proof:  $\leftarrow$  for any subdisjunction  
 $\rightarrow$  similar idea as  $\mathcal{V}$ -ab. gr. doesn't have DPC  
add constants  $a_1, \dots, a_n$  to the signature

axioms for  $\mathcal{V}$   
 $\varphi(a_1, \dots, a_n)$   
 $\neg \varphi_1(a_1, \dots, a_n)$   
 $\vdots$  } no model  $\Rightarrow$  some finite subset has no model  $\Rightarrow$

$\Rightarrow$  axioms for  $\mathcal{V}$   
 $\varphi(a_1, \dots, a_n)$   
 $\neg \varphi_{m_1}(a_1, \dots, a_n)$   
 $\vdots$   
 $\neg \varphi_{m_k}(a_1, \dots, a_n)$  } no model  $\Rightarrow$   $\rightarrow$

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Theorem  $\mathcal{V}$  has DPC & finitely many SIs, all finite  
 $\Rightarrow \mathcal{V}$  is finitely based

Prf:  $\mathcal{V}$  has DPC witnessed by conservative  $\varphi(x, y, a, b)$

$$SI(\mathcal{V}) = \{ \underline{A}_1, \dots, \underline{A}_k \} \text{ (up to iso)}$$

consider  ~~$\varphi_1 \wedge (\varphi_2 \rightarrow \varphi_3)$~~   $\alpha_1 \wedge (\alpha_2 \rightarrow \alpha_3)$

where  $\alpha_1 \equiv$  " $\forall a, b \varphi\text{-}Cg^A(a, b)$  is a congruence of  $\underline{A}$  containing  $(a, b)$ "

$\alpha_2 \equiv$  " $\underline{A}$  is SI according to  $\varphi$ "

$$\exists a \neq b \forall x \neq y \varphi(a, b, x, y)$$

$\alpha_3 \equiv$  " $\underline{A} \in \{ \underline{A}_1, \dots, \underline{A}_k \}$ "

should be clear

$$\forall a, b, x \varphi(x, x, a, b)$$

" $\varphi\text{-}Cg(a, b)$  is reflexive"

$$\wedge \forall a, b, x, y \varphi(x, y, a, b) \rightarrow \varphi(y, x, a, b)$$

" $\varphi\text{-}Cg(a, b)$  is symmetric"

$\wedge$  " $\varphi\text{-}Cg(a, b)$  is transitive"

$$\wedge \forall a, b Cg \varphi(a, b, a, b)$$

$\wedge$  " $\varphi\text{-}Cg(a, b)$  compatible with operations"

$\mathcal{V} \models \alpha_1 \wedge (\alpha_2 \rightarrow \alpha_3) \Rightarrow Id(\mathcal{V}) \vdash \alpha_1 \wedge (\alpha_2 \rightarrow \alpha_3) \Rightarrow$   
 $\Rightarrow \exists \mathcal{E} \subseteq_{fin} Id(\mathcal{V})$  such that  $\mathcal{E} \vdash \alpha_1 \wedge (\alpha_2 \rightarrow \alpha_3)$ .

$\mathcal{W} = Mod(\mathcal{E})$ , clearly  $\mathcal{V} \subseteq \mathcal{W}$

$\mathcal{W} \stackrel{?}{\subseteq} \mathcal{V}$ . Enough  $\underline{A} \in \mathcal{W}$  is SI  $\Rightarrow \underline{A} \in \{ \underline{A}_1, \dots, \underline{A}_k \}$

$a, b \in A$   $\underline{A} \models \alpha_1 \rightarrow \varphi\text{-}Cg^A(a, b)$  is congruence containing  $(a, b)$  }  $Cg^A(a, b) =$   
 $\varphi$  is conservative  $\rightarrow \varphi\text{-}Cg^A(a, b) \subseteq Cg^A(a, b)$  }  $= \varphi\text{-}Cg^A(a, b)$