

Functorial algebras and coalgebras

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Scheme of talk

structural
problems

representations

old period

new period

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structural
problems

representations

old period

I.

II.

new period

III.

IV.

I. Structural properties (old period)

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Basic definitions: $F : \mathcal{K} \rightarrow \mathcal{K}$

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$$\begin{array}{ccccc} \text{Alg } F : FQ & \xrightarrow{\sim Fh} & FQ' & \xrightarrow{\sim Fh'} & FQ'' \\ \delta \downarrow & & \delta' \downarrow & & \delta'' \downarrow \\ Q & \xrightarrow{\sim h} & Q' & \xrightarrow{\sim h'} & Q'' \end{array}$$

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[Bialgebras $A(F, G)$]

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[Bialgebras $A(F, G)$]

The category $\mathbb{A}(\Sigma)$ of all universal algebras of the type (Σ, ar) [ar being "the arity function" $\Sigma \rightarrow \text{Card}$] is just $\text{Alg}F_\Sigma$ for the polynomial functor $F_\Sigma : \text{Set} \rightarrow \text{Set}$.

$$F_\Sigma X = \prod_{\sigma \in \Sigma} X^{ar\sigma}$$

The oldest bibliography:

- ▶ O. Wyler: Operational categories, Proc. Conference on categorical algebra Springer Verlag Berlin-Haidelberg-New York (1966), 295-316

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- ▶ V. Trnková: Some properties of set functors, Comment. Math. Univ. Carolinae 10 (1969), 323-352

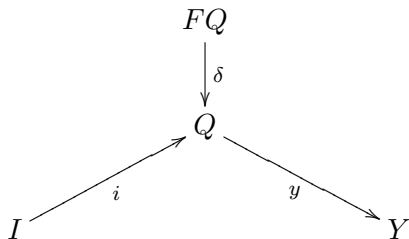
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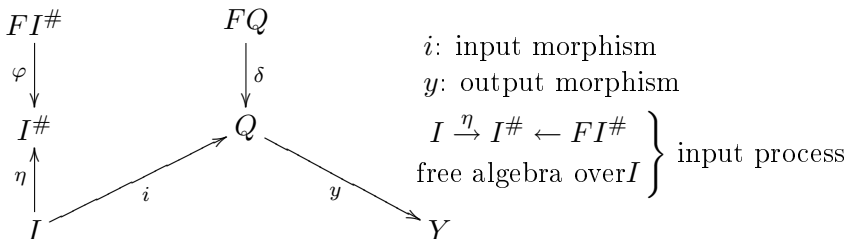
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i : input morphism
 y : output morphism

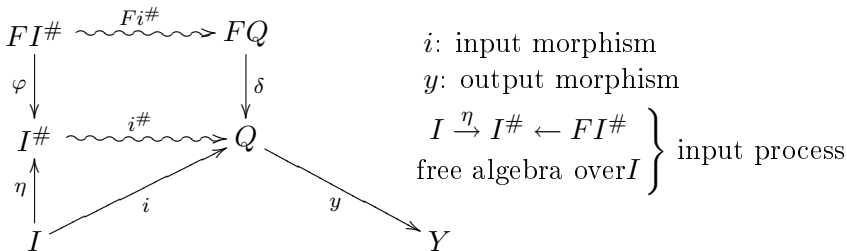
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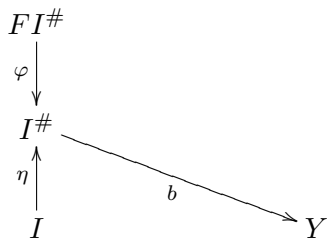
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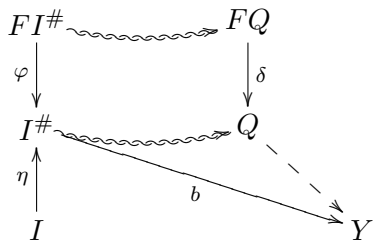
$b = y \circ i\#$ is the behavior of the machine

minimal realization of a given behavior

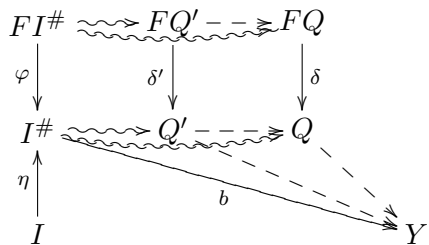
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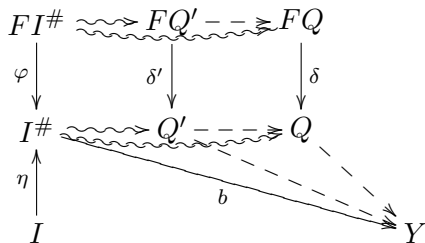
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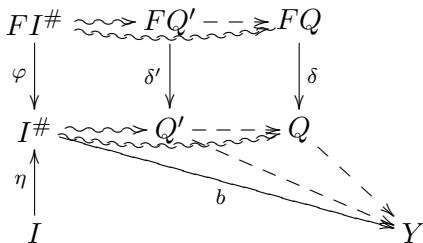


minimal realization of a given behavior



- A) Which functors are input processes?
 B) Which input process admits minimal realization?

minimal realization of a given behavior



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Varietor

Iterative constructions

Iterative constructions

Knaster-Tarski construction of the least fix-point of a monotone map $F : \mathcal{K} \rightarrow \mathcal{K}$, \mathcal{K} complete lattice:

$$\begin{array}{ccccccccccc} 0 & \leq & F0 & \leq & FF0 & \leq & \dots & \sup F^i 0 & \leq & F(\dots) & \leq & \dots \\ \parallel & & \parallel & & \parallel & & & \parallel & & \parallel & & \\ W_0 & & W_1 & & W_2 & & & W_\omega & & W_{\omega+1} & & \end{array}$$

The construction stops if $FW_\alpha = W_\alpha$

Initial-algebra-construction: \mathcal{K} category, $F : \mathcal{K} \rightarrow \mathcal{K}$
endofunctor, \mathcal{K} has an initial object I (ϕ in Set , $\{0\}$ in $Vect$
...) \mathcal{K} has colimits of chains

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$$\begin{array}{ccccccc}
 I & \xrightarrow{w_{0,1}} & FI & \xrightarrow{w_{1,2}} & F^2I & \rightsquigarrow \dots \rightsquigarrow & \text{colim}(F^i I, w_{i,\omega}) \rightsquigarrow F(\dots) \rightsquigarrow \dots \\
 \parallel & & \parallel & & \parallel & & \parallel & & \parallel \\
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 $w_{\alpha,\alpha+1} : W_\alpha \rightarrow FW_\alpha = W_{\alpha+1}$ is an isomorphism.

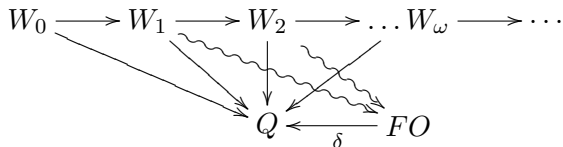
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Adámek: free-algebra-construction

FI

I
 \parallel
 W_0

Adámek: free-algebra-construction

$$\begin{array}{ccc} FI & & \\ & \searrow \varphi_0 & \\ I & \xrightarrow{w_{0,1}} & I + FI \\ \parallel & & \parallel \\ W_0 & & W_1 \end{array}$$

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Then $(W_\alpha, (w_{\alpha,\alpha+1})^{-1} \circ \varphi_{\alpha+1})$ and $\eta : I \rightarrow W_\alpha$ with $\eta = w_{0,\alpha}$ is a free F -algebra over I .

Koubek: push-out construction

$$\begin{array}{ccc} D & \xrightarrow{\alpha} & FW_0 \\ & \searrow \beta & \\ & & W_0 \end{array}$$

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The diagram illustrates a push-out construction. The top row shows a sequence of objects D, FW_0, FW_1, FW_2 connected by arrows $\alpha, F_{w_{0,1}}, F_{w_{2,3}}$. The bottom row shows a sequence of objects W_0, W_1, W_2, W_3 connected by arrows $w_{0,1}, w_{1,2}, w_{2,3}$. A diagonal arrow β maps D to W_0 . Commutative squares are indicated by arrows labeled "P.O." (Push-Out) connecting (FW_0, W_0) , (FW_1, W_1) , and (FW_2, W_2) . The arrow β_1 is also shown between FW_0 and W_0 .

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It depends on \mathcal{K} and F , there are many criteria and many examples.

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YES in *Set*

NO generally

J. Adámek, V. Trnková:

Automata and Algebras in
Categories
Kluwer Academic Publishers 1990

II. Representations (old period)

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J. de Groot 1959: Every group is isomorphic to the autohomeomorphism group of a topological space.

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It was an idea of J. R. Isbell to generalize it to the investigation of full embeddings of categories.

A functor $F : \mathcal{K} \rightarrow \mathcal{H}$ is a *full embedding* if it is isofunctor onto full subcategory (i.e. F maps $\mathcal{K}(a, b)$ onto $\mathcal{H}(Fa, Fb)$).

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A category \mathcal{K} is called *algebraically universal* if every category $\mathbb{A}(\Sigma)$ of universal algebras of an arbitrary signature Σ can be fully embedded in it.

Many categories were proved to be alg-universal (Hedrlín, Pultr, Vopěnka, Kučera, Sichler, ...).

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Theorem(V. Koubek 1983): For a variety $F : Set \rightarrow Set$, $Alg F$ is alg-universal iff F satisfies at least one of the following three conditions:

1. $Ident + Ident$ is subfunctor of F
2. F is faithful and has an unattainable cardinal
3. F is not connected and has an unattainable cardinal.

[A set PX is called an *increase* of a functor $F : Set \rightarrow Set$ on a set X if

$$PX = FX \setminus \bigcup_{\substack{f: Y \rightarrow X \\ |Y| < |X|}} (Ff)[FY]$$

$|X|$ is *unattainable* iff $PX \neq 0$]

Example:

$Alg(2)$ is alg-universal; its variety given by

$$x \cdot x = y \cdot y$$

is not alg-universal; but $Alg(2,0)$ with

$$x \cdot x = y \cdot y$$

is alg-universal.

III. Structural problems (new period)

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J. M. Rutten: Universal coalgebra: a system theory

Theoret. Comp. Sci. 249 (2000), 3-80

H. P. Gumm and T. Schröder: Types and coalgebraic structure

Alg. Universalis 53 (2005), 229-252



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▶ J. Adámek and V. T.: Initial algebras and terminal coalgebras in many-sorted sets, sent for publication

▶ J. Adámek and V. T.: Relatively terminal coalgebras, in preparation

Selected parts:

- ▶ Let $F : \mathcal{K} \rightarrow \mathcal{K}$ admit terminal coalgebra. Must it be obtained by corresponding iterative construction?
YES for $\mathcal{K} = \mathit{Set}$, many sorted sets, vector spaces

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For initial algebras the result is analogous

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[For terminal coalgebras in \mathbf{Set} between λ and 2λ for λ -accessible functors, by J. Worrell, 2005]

Open problem: Which functors $\mathbf{Set} \rightarrow \mathbf{Set}$ admit terminal coalgebras?

IV. Representations

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- ▶ J. Sichler and V. T.: On universal categories of coalgebras, accepted in Alg. Universalis
- ▶ V. Koubek, J. Sichler and V. T.: Universality of categories of coalgebras, accepted in Appl. Cat. Structures

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A category is called *coalgebraically universal* if every category of universal coalgebras can be fully embedded in it.

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A category is called *coalgebraically universal* if every category of universal coalgebras can be fully embedded in it.

Statement: A category is coalgebraically universal if and only if it is algebraically universal.

IV. Representations

- ▶ J. Sichler and V. T.: On universal categories of coalgebras, accepted in Alg. Universalis
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Statements:

- ▶ if F preserves intersections, then $CoalgF$ is universal iff F is not linear (i.e. F is not of the form $FX = X \times A + B$ where A and B are constants).
- ▶ if F preserves non-void preimages then $CoalgF$ is universal iff F does not preserve finite unions of non-empty set.

Examples:

$F_1(X) = X \times X$: both $AlgF_1$ and $CoalgF_1$ are universal.

$F_2(X) = X \times X/\Delta$: $CoalgF_2$ is universal, $AlgF_2$ not.

$F_3(X) = X + X$: $AlgF_3$ is universal $CoalgF_3$ is not.

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Let $\beta : Set \rightarrow Set$ be an ultrafilter functor. Then

- ▶ $Coalg\beta$ contains rigid proper class of coalgebras.
- ▶ But it is not universal, a monoid on 13 elements cannot be represented in it.