

On the equational theory of projection lattices of finite von Neumann factors

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on joint work with

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Jardafest, Praha July 22, 2010

0. MODULAR (ORTHO)LATTICES

We consider *modular lattices* (shortly MLs) - we write $a + b$ for joins and ab for meets - and *modular ortholattices* (shortly MOLs), a subject started by Birkhoff and von Neumann [1]. These have constants 0 and 1 and a unary fundamental operation $x \mapsto x'$ which is an *involution* and, moreover, an *orthocomplementation*

$$x'' = x, \quad x \leq y \Rightarrow y' \leq x'; \quad x \oplus x' = 1,$$

The principal examples are the lattices $L(M)$ of all submodules of some R -module M and the lattices

$$L(V, \Phi) = \{X, X^\perp \mid X \in L(V), \dim X < \infty\}, \quad X \mapsto X^\perp$$

where (V, Φ) is an inner product space. Considering lattice identities $\forall \bar{x}. f(\bar{x}) = g(\bar{x})$, we may assume that $f(\bar{x}) \leq g(\bar{x})$ holds in all lattices. In MOLs it suffices to consider identities of the form $\forall \bar{x}. t(\bar{x}) = 0$ - put $t = gf'$.

1. IDENTITIES IN THE ATOMISTIC CASE

Proposition 1.1. *If L is an atomistic ML or MOL, then $L \models \forall \bar{x}. f(\bar{x}) = g(\bar{x})$ if and only if $[0, u] \models \forall \bar{x}. f(\bar{x}) = g(\bar{x})$ for all interval subalgebras $[0, u]$ where $\dim u$ is at most the number of occurrences of variables in f and g , together.*

Here, in the case of MOLs, $[0, u]$ is endowed with the orthocomplement $x^u = x'u$.

Lemma 1.2. *In an atomistic ML, if $h(\bar{x})$ is a lattice term with unique occurrence of variables, and if $p \leq h(\bar{a})$ for some atom p , then there is a substitution by atoms such that $p \leq h(\bar{q})$.*

Proof. By induction. If $p \leq h_1(\bar{a}_1) + h_2(\bar{a}_2)$ then there are $p_i \leq h_i(\bar{a}_i)$ such that $p \leq p_1 + p_2$ - this is, in essence, the Theorem on Joins in projective spaces. By inductive hypothesis we have $p_i \leq h_i(\bar{q}_i)$ for some substitutions by atoms. \square

Proof. Prop.1. In the lattice case, assume $f(\bar{a}) < g(\bar{a})$. Then there is $p \leq g(\bar{a})$, $p \not\leq f(\bar{a})$. Consider $h(\bar{z})$ with unique occurrence and substitution $h(\bar{a}) = g(\bar{a})$. Then by the Lemma there is a substitution by atoms such that $p \leq h(\bar{q})$. Collect all q associated with the same a into the join $c \leq a$ and let u be the join of all c . Then $\dim u$ is small as indicated and, by monotonicity of lattice terms, we have $p \leq g(\bar{c})$ but $p \not\leq f(\bar{c})$ in $[0, u]$.

In the MOL case, we replace $t(\bar{x})$ by its negation normal form and associate with that a lattice term $h(\bar{y}, \bar{z})$ with unique occurrence of variables such that any y stands for a positive occurrence of some x , any z for an occurrence of x' . Now, if $0 < t(\bar{a})$, then $p \leq t(\bar{a}) = h(\bar{b}, \bar{c})$ where all b 's are a 's and all c 's are a' s. By the Lemma, $p \leq h(\bar{q}, \bar{r})$ for some substitution by atoms whence $p \leq t(\bar{d})$ for some $\bar{d} \in [0, u]$, u the joins of all q 's and r 's. Again, $\dim u$ is bounded as stated. \square

Theorem 1.3. Huhn, H., Czédli, Hutchinson. *For any division ring D with prime subfield F_p ,*

$$\text{Th}_{eq}\{L(V) \mid V \text{ a } D\text{-vector space}\} = \text{Th}_{eq}\{L(F_p^n) \mid n < \infty\}$$

and this equational theory is decidable.

Proof. [7]. $L(F_p^n)$ is a sublattice of $L(D^n)$. Any $L(V_D)$ is a sublattice of $L(V_{F_p})$. The latter is atomistic, whence by Prop.1 in the variety generated by the $L(F_p^n)$. This proves that the varieties coincide. Thus, the set of non-valid identities is recursively enumerable. On the other hand, the quasi-variety generated by the $L(V_D)$ is recursively axiomatizable - using Mal'cev's method of axiomatic correspondence. Thus, the equational theory is recursively enumerable, too. A reasonable decision procedure has been provided by Czédli and Hutchinson [3]. \square

For MOLs we define the *von Neumann variety*

$$\mathcal{N} = \text{HSP}\{L(\mathbb{R}^n) \mid n < \infty\}$$

Theorem 1.4. $\mathcal{N} = \text{HSP}\{L(\mathbb{C}^n) \mid n < \infty\}$ and $\text{Th}_{eq}\mathcal{N}$ is decidable.

Here, we consider the canonical real resp. complex scalar products.

Proof. We have the following embeddings $L(\mathbb{R}^n) \subseteq L(\mathbb{C}^n) \subseteq L(\mathbb{R}^{2n})$. Also, due to Tarksi [22], $\text{Th}L(\mathbb{R}^n)$ is decidable. Now, apply Prop.1. \square

2. INTERPRETATION OF RINGS VIA FRAMES

A (von Neumann) *n-frame* is a system a_{ij} ($1 \leq i, j \leq n$) of constants and relations such that in a lattice $L(M)$, M a free R -module on generators e_1, \dots, e_n and R a ring with unit, these relations are satisfied

for

$$a_{ii} = e_i R, \quad a_{ij} = a_{ji} = (e_i - e_j)R$$

and, conversely, any system satisfying the relations is, up to isomorphism, this canonical one. In particular, the ring R can be interpreted into $L(M)$ via $r \mapsto (e_1 - e_2r)R$.

- (1) For any modular L with an n -frame, $n \geq 4$, one obtains a ring on $\{x \in L \mid x \oplus a_2 = a_1 a_2\}$. Von Neumann [19], mimicking the above.
- (2) There are terms $t_{ij}(\bar{x})$ such that for an \bar{a} in a modular lattice the $t_{ij}(\bar{a})$ form an n -frame in the interval $[\prod_{ij} t_{ij}(\bar{a}), \sum_{ij} t_{ij}(\bar{a})]$. Moreover, $t_{ij}(\bar{a}) = a_{ij}$ if \bar{a} is an n -frame, already. G. Bergman and A. Huhn [13].
- (3) For MOLs this extends to *orthogonal n -frames*: $a_{ii} \leq a'_{jj}$ for $i \neq j$ and $\prod_{ij} a_{ij} = 0$. R. Mayet and M. Roddy [16].

3. UNIFORM WORD PROBLEM

A *quasi-identity* is a first order sentence of the form

$$\forall \bar{x}. \left(\bigwedge_{i=1}^n s_i(\bar{x}) = t_i(\bar{x}) \right) \Rightarrow s(\bar{x}) = t(\bar{x})$$

Solvability of the *uniform word problem* for a class \mathcal{C} of algebraic structures means decidability of the set of quasi-identities valid in \mathcal{C} . Let \mathcal{Q} denote the quasi-variety generated by \mathcal{C} , and \mathcal{Q}_L the class of lattices embedded into reducts of members of \mathcal{Q} .

Theorem 3.1. *If, for some field F , $L(F^n) \in \mathcal{Q}_L$ for all $n < \infty$, then the uniform word problem for \mathcal{C} is unsolvable.*

Proof. Let \mathcal{S} denote the class of all semigroups, $F^{n \times n}$ the ring of all $n \times n$ -matrices over F , and F_p the prime subfield.

- (1) $L(F^n) \cong \overline{L}(F^{n \times n})$, the lattice of principal right ideals
- (2) $\text{Th}_{\forall}\{F^{n \times n} \mid n < \infty\} = \text{Th}_{\forall}\{F_p^{n \times n} \mid n < \infty\} = \text{Th}_{\forall}\mathcal{S}_{fin}$ considering multiplicative semigroups Lipshitz [15]
- (3) $\text{Th}_{qid}\mathcal{S} \subseteq \Gamma \subseteq \text{Th}_{qid}\mathcal{S}_{fin}$ for no recursive Γ . Gurevich, Lewis [6].
- (4) Interpret $\mathcal{S} \rightarrow \text{Rings} \rightarrow \text{ML}$ via $F \mapsto F[S]$ and $R \mapsto L$ via frames

□

4. RESTRICTED WORD PROBLEM

The restricted word problem for \mathcal{C} considers in each instance a fixed premise $\bigwedge_{i=1}^n s_i(\bar{x}) = t_i(\bar{x})$. For a quasi-variety, this amounts to considering a finite presentation. Unsolvability means the existence of some instance with undecidable decision problem.

Theorem 4.1. Lipshitz, Hutchinson [15, 14]. *If $L(M) \in \mathcal{Q}_L$ for some free module M on an infinite basis, then the restricted word problem for \mathcal{C} is unsolvable - there is a presentation on 5 lattice generators.*

Indeed, any finitely presented semigroup can be interpreted into some $L(M^n)$.

Theorem 4.2. Cohn, McIntyre. *There is a finitely presented division ring D with unsolvable word problem*

Corollary 4.3. *If $L(D^n) \in \mathcal{Q}_L$ for such D and some $n \geq 4$ then the restricted word problem for \mathcal{Q} is unsolvable.*

Theorem 4.4. Roddy [21]. *The restricted word problem for MOLs is unsolvable - there is a presentation on 3 generators.*

This is based on an intricate construction of a division ring as above admitting a scalar product on some D^n - $n = 14$.

5. UNDECIDABLE EQUATIONAL THEORIES

Theorem 5.1. Freese [4]. *The equational theory of all modular lattices is undecidable - 5 and even 4 variables suffice.*

The proof is based on the above division rings, frames, and an ingenious device allowing to force relations via terms in free modular lattices.

Proposition 5.2. *For D as above and $n \geq 3$, $\text{Th}_{eq}L(D^n)$ is undecidable.*

Proof. The terms for an n -frame will either yield an n -frame of $L(D^n)$ or just a single element. In the first case, one has terms giving elements of the ring associated with the frame or else a collapse of the frame. Again, considering relations on those ring elements one has terms enforcing these relations simultaneously - or else a collapse. Also, when applied to elements satisfying the relations, these remain unchanged. \square

6. SATISFIABILITY PROBLEMS

Dealing with modular lattices of finite height we consider 0 and 1 as constants.

Lemma 6.1. *Let L be an ML of height n with an n -frame a_{ij} . For any pair $f(\bar{x}), g(\bar{x})$ of lattice terms one can construct lattice polynomials $f^-(\bar{x}), g^+(\bar{x})$ with constants a_{ij} such that the following are equivalent*

- (1) $L \models \exists \bar{x}. f(\bar{x}) < g(\bar{x})$
- (2) $L \models \exists \bar{x}. f^-(\bar{x}) = 0 \ \& \ g^+(\bar{x}) = 1$

If L admits an involution such that $a_{jj} \leq a'_{ii}$ for $j \neq i$, then one can construct $h(\bar{x}, \bar{y})$ and add

- (3) $L \models \exists \bar{x}. h(\bar{x}, \bar{x}') = 1$

Construction and identification of the output polynomials, as well as reconstruction of the input terms can all be done in PTIME.

Proof. Define $b_k = \sum_{i \leq k} a_{ii}$, $f_k^- = a_{kk}(b_{k-1} + f)$, and $f^- = \prod_k f_k^-$. Similarly, for g^+ . Put $h = g(\bar{x})\tilde{f}(\bar{y})$ where \tilde{f} arises from f^- by interchanging $+$ and \cdot and replacing the constants by the corresponding elements of the dual n -frame canonically associated with the given n -frame. In the case of MOLs this gives rise to a ternary discriminator polynomial on L [9]. \square

Theorem 6.2. *Let F be a field and $n \geq 3$.*

- (i) *With each polynomial $p(\bar{x})$ over F one can associate lattice terms $p^-(\bar{y})$ and $p^+(\bar{y})$ such that*

$$F \models \exists \bar{x}. p(\bar{x}) = 0 \Leftrightarrow L(F^n) \models \exists \bar{y}. p^-(\bar{y}) = 0 \ \& \ p^+(\bar{y}) = 1$$

- (ii) *With any pair $s(\bar{y}), t(\bar{y})$ of lattice terms one can associate polynomials $p_1(\bar{x}), \dots, p_n(\bar{x})$ with integer coefficients such that*

$$L(F^n) \models \exists \bar{y}. s(\bar{y}) = 0 \ \& \ t(\bar{y}) = 1 \Leftrightarrow F \models \exists \bar{x}. p_1(\bar{x}) = \dots = p_n(\bar{x}) = 0$$

- (iii) *If F^n admits an inner product Φ then with any ortholattice term $t(\bar{y})$ one can associate integer $p_i(\bar{x})$ such that*

$$L(F^n, \Phi) \models \exists \bar{y}. t(\bar{y}) = 1 \Leftrightarrow F \models \exists \bar{x}. p_1(\bar{x}) = \dots = p_n(\bar{x}) = 0$$

All this can be done in PTIME and does not depend on F for polynomials with integer coefficients.

Here, we conceive the $p(\bar{x})$ primarily as terms. But, as far as solvability is concerned, transition to a linear combination of monomials can be done in PTIME - adding variables.

Proof. In (i) use the lattice terms providing an n -frame and the interpretation of F into $L(F^n)$. In (ii) and (iii) replace the (ortho)lattice variables by matrices with variables for elements of F and recall the descriptions of joins, meets, and orthocomplements in $\overline{L}(F^{n \times n})$. Solving $t(\bar{y}) = 1$ amounts to capturing the identity matrix. \square

Corollary 6.3. *Let F be a subfield of \mathbb{R} and Φ the canonical scalar product on F^n where $n \geq 3$. Then the following satisfiability problems are polynomially equivalent*

$$\begin{array}{ll} L(F^n) \models \exists \bar{x}. f(\bar{x}) < g(\bar{x}) & f \leq g \text{ lattice terms} \\ L(F^n) \models \exists \bar{x}. f(\bar{x}) = 0 \& g(\bar{x}) = 1 & f, g \text{ lattice terms} \\ L(F^n, \Phi) \models \exists \bar{x}. t(\bar{x}) = 1 & t \text{ ortholattice term} \\ F \models \exists \bar{x}. p(\bar{x}) = 0 & p \text{ integer polynomial} \end{array}$$

As remarked by George McNulty, decidability of the latter is an open and controversial question for $F = \mathbb{Q}$ [20]. To get p from the p_i put $p = \sum_i p_i^2$.

7. REAL COMPLEXITY

Henceforth, we consider \mathbb{R}^n and \mathbb{C}^n always with the canonical scalar product Φ .

Corollary 7.1. *The decision problems for each single $\text{Th}_{eq}L(\mathbb{R}^n, \Phi)$, $n \geq 3$, als well as for $\text{Th}_{eq}\mathcal{N}$ are polynomially equivalent and $coBP(NP_{\mathbb{R}}^0)$ -complete. In particular, they are $coNP$ -hard and in $PSPACE$.*

Here, $BP(NP_{\mathbb{R}}^0)$ refers to non-deterministic polynomial time in the Blum-Shub-Smale model of real computation with constants 0, 1, only, and binary input.

Proof. The equational theory of a class \mathcal{C} is just the complement of the set of sentences $\exists \bar{x}. t(\bar{x}) = 1$ satisfiable in some member of \mathcal{C} . Thus, the claim about the $L(\mathbb{R}^n, \Phi)$ follows from Cor.6.3 and the fact that feasability of integer polynomials over \mathbb{R} is known to be $BP(NP_{\mathbb{R}}^0)$ -complete cf. [17]. Also, with Prop.1 it follows that $\text{Th}_{eq}\mathcal{N}$ is in $coBP(NP_{\mathbb{R}}^0)$. To prove completeness, we interpret feasability of integer polynomials via 3-frames into $L(\mathbb{R}^{3n}, \Phi)$ for all $n \geq 1$ simultaneously: according to $L(\mathbb{R}^{3n}) \cong L((\mathbb{R}^{n \times n})^3)$ we see the x_i as variables for matrices $A_i \in \mathbb{R}^{n \times n}$. Imposing the relations $A_i = A_i^t$ and $A_i A_j = A_j A_i$, which we can enforce via ortholattice terms to be built into the identity $t(\bar{y}) = 1$, we achieve that the A_i are simultaneously diagonalizable, whence from $p(\bar{A}) = 0$ we obtain a solution in \mathbb{R} . The cases $3n+1$ and $3n+2$ are dealt with considering the $3n$ -part of the frame. \square

8. MOL-REPRESENTATIONS

Given an inner product space (V, Φ) which is an elementary extension of an unitary space, an *e-unitary representation* of an MOL L is a 0-lattice embedding $\varepsilon : L \rightarrow L(V)$ such that

$$\varepsilon(a') = \varepsilon(a)^\perp \text{ for all } a \in L.$$

Theorem 8.1. Bruns, Roddy, H. [2, 8, 11]. *For any e-unitary representation of an MOL, there is an atomic MOL \tilde{L} which is a sublattice of $L(V)$ and contains both $\varepsilon(L)$ and $L(V, \Phi)$ as sub-OLs.*

Proposition 8.2. H., Roddy [8, 11]. *$L \in \mathcal{N}$ if L admits an e-unitary representation. For subdirectly irreducibles, the converse holds, too.*

Proof. Prop.1 and the fact, that all sections of fixed finite height have the same first order theory. Conversely, by the Jónsson Lemma we have $L \in \text{HSP}_u\{L(\mathbb{C}^n) \mid n < \infty\}$ and show that representability is preserved. \square

9. *-REGULAR RINGS AND REPRESENTATIONS

An associative ring (with or without unit) R is (von Neumann) *regular* if for any $a \in R$ there is a *quasi-inverse* $x \in R$ such that $axa = a$. A $*$ -ring is a ring with an involution $*$ as additional operation:

$$(x + y)^* = x^* + y^*, \quad (xy)^* = y^*x^*, \quad x^{**} = x.$$

e is a *projection* if $e = e^* = e^2$. A $*$ -ring is *$*$ -regular* if it is regular and, moreover, *positive*: $xx^* = 0$ only for $x = 0$. Equivalently, for any $a \in R$ there is a (unique) projection e such that $aR = eR$. Examples are the $\mathbb{C}^{n \times n}$ with r^* the adjoint matrix. The projections of a $*$ -regular ring with unit form an MOL $\overline{L}(R)$ where $e \leq f \Leftrightarrow e = ef$ and $e' = 1 - e$. Now, $e \mapsto eR$ is an isomorphism of $\overline{L}(R)$ onto the ortholattice of principal right ideals of R and we may use the same notation for both.

Let (V, Φ) be an elementary extension of a unitary space. Denote by ϕ^* the adjoint of ϕ - if it exists. An *e-unitary representation* of a $*$ -ring R is a ring embedding $\iota : R \rightarrow \text{End}(V)$ such that $\iota(r^*) = \iota(r)^*$ for any $r \in R$.

Proposition 9.1. Giudici [5]. *If $\iota : R \rightarrow \text{End}(V)$ is an e-unitary representation of the $*$ -regular ring R , then*

$$\varepsilon(eR) = \text{Im } \iota(e)$$

is an e-unitary representation of the MOL $\overline{L}(R)$ in (V, Φ) .

10. VON NEUMANN ALGEBRAS

A *von-Neumann algebra* \mathbf{M} is an unital involutive \mathbb{C} -subalgebra of the algebra $\mathcal{B}(H)$ of all bounded operators of a separable Hilbert space H with $\mathbf{M} = \mathbf{M}''$ where $\mathbf{A}' = \{\phi \in \mathcal{B}(H) \mid \phi\psi = \psi\phi \quad \forall \phi \in \mathbf{A}\}$ is the *commutant* of \mathbf{A} . \mathbf{M} is *finite* if $rr^* = 1$ implies $r^*r = 1$. For such, the projections of \mathbf{M} form a (continuous) MOL $L(\mathbf{M})$. A finite von-Neumann algebra is a *factor* if its center is $\mathbb{C} \cdot 1$. Particular examples of finite factors are the algebras $\mathbb{C}^{n \times n}$ of all complex n -by- n -matrices.

Theorem 10.1. Murray-von-Neumann [18]. *Any finite von-Neumann algebra factor is either isomorphic to $\mathbb{C}^{n \times n}$ for some $n < \infty$ (type I_n) or contains for any $n < \infty$ a subalgebra isomorphic to $\mathbb{C}^{n \times n}$ (type II_1).*

Theorem 10.2. Murray-von-Neumann [18]. *For every finite factor \mathbf{M} , there is a $*$ -regular ring $U(\mathbf{M})$ of unbounded operators on H having \mathbf{M} as $*$ -subring and such that ϕ^* is adjoint to ϕ . Moreover, \mathbf{M} and $U(\mathbf{M})$ have the same projections.*

Theorem 10.3. *$U(\mathbf{M})$ admits an e -unitary representation.*

Proof. By the Compactness Theorem, it suffices to consider countable $*$ -subrings R of $U(\mathbf{M})$. A representation of R is constructed from the given Hilbert space H . Let H_0 be the intersection of all domains of operators $\phi \in R$. Define, recursively, H_{n+1} as the intersection of H_n and all preimages $\phi^{-1}(H_n)$ where $\phi \in R$. $H_\omega = \bigcap_{n < \omega} H_n$. Due to Murray and von Neumann, all H_n and H_ω are dense in H . It easily follows, that $\varepsilon(\phi) = \phi|_{H_\omega}$ defines a representation. \square

Corollary 10.4. $\text{Th}_{eq}\mathcal{N} = \text{Th}_{eq}L(\mathbf{M})$ for any finite von Neumann algebra factor \mathbf{M} of infinite dimension.

Proof. Observe $L(\mathbf{M}) \cong \overline{L}(U(\mathbf{M}))$ and apply Prop.8.2 and 9.1. \square

Corollary 10.5. *For any finite von Neumann algebra factors \mathbf{M} and \mathbf{N}*

$$U(\mathbf{N}) \in \text{HSP}_u U(\mathbf{M}), \quad U(\mathbf{N}) \in \text{HSP}_u \{\mathbb{C}^{n \times n} \mid n < \infty\}$$

and, analogously, for the projection lattices.

Proof. With suitable choice of quasi-inverse, $*$ -regular rings form a congruence distributive variety - the congruence lattice of R is isomorphic to that of $\overline{L}(R)$. The $L(\mathbf{M})$ are simple. Thus, the Jónsson Lemma can be applied. \square

A question, raised by Connes and still unanswered, asks whether the Banach-space version of this result is true.

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