Domain Decomposition Strategies for Adaptively Refined Meshes

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Many ways how to improve performance of finite element codes, two very different strategies are:

- parallel, domain decomposition, the use of supercomputers, very large linear systems
- adaptivity, higher order, the goal is to have smaller linear systems solvable on PC, still with good accuracy
- It would be nice to combine both strategies
 - Parallel calculations
 - take advantage from similarity of structure over the domain
 - most of the research done in linear solver part
 - Exascale brings many challenges. It could be good alternative to get more from petascale instead.
 - Adaptivity
 - different treatment of different areas, based on solution behavior
 - a lot of effort in assembly part, trying to minimize number of DOFs
 - There are certain limits for single PC, no matter how smart the algorithm is.



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Ingredients

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I Experience with adaptivity and higher order finite elements

[P. K., P. Šolín, D. Andrš, *Arbitrary-level hanging nodes for adaptive hp-FEM approximations in 3D*, JCAM, 270, pp. 121–133, 2014]

Experience with domain decomposition, BDDCML library

[B. Sousedík, J. Šístek, and J. Mandel, *Adaptive-Multilevel BDDC and its parallel implementation*, Computing, 95 (12), pp. 1087–1119, 2013.]



3 Parallel mesh handler p4est



[C. Burstedde, L. Wilcox, and O. Ghattas, *p4est: Scalable Algorithms for Parallel Adaptive Mesh Refinement on Forests of Octrees*, SIAM J. Sci. Comput., 3 (33), pp. 1103–1133, 2011.]





- Without balancing has no sense in parallel
- Most of the refinements would concentrate in those domains, where singularities, boundary or internal layers are present
- It might be just few subdomains





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- Element incidence graph can be created and partitioned
- Advantage: usually nice shape of subdomains
- Disadvantage: it is not scalable to large number of processors





- Z-order (space filling) curve can be used
- Each element in the refinement hierarchy might be identified by number, coding bitwise its position and level of refinement
- Can be used both in 2D (quadrilaterals) or 3D (hexahedra)





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- Can be used for partitioning
- Initial mesh can be small, just to express the geometry. It has to be made of quadrilaterals or hexahedra, which might be limiting.
- Disadvantage: shape of the subdomains far from optimal





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Hanging nodes

- Hanging nodes have to be eliminated
- They can also appear at the subdomain interface

2 Shape of the subdomains

- The shape is far from perfect
- Subdomains might be disconnected or only loosely coupled (e.g. by one node in elasticity)

The BDDC preconditioner



- Balancing Domain Decomposition based on Constraints [Dohrmann (2003)], [Cros (2003)], [Fragakis, Papadrakakis (2003)]
- continuity at corners, and of averages (arithmetic or weighted) over edges or faces considered
- enough constraints to fix floating subdomains $a(\cdot, \cdot)$ symmetric positive definite on \widetilde{W}
- corresponding matrix \widetilde{A} symmetric positive definite, almost block diagonal structure, larger dimension than A
- used to construct (an action of) preconditioner M^{-1} to solve

$$M^{-1}Su_{\Gamma} = M^{-1}g$$



- nullspaces of subdomain matrices unknown a priori
- detect graph components of subdomain mesh
- components independent during classification of interface into faces, edges and vertices
- corners selected by the face-based algorithm [Šístek et al. (2011)]
- size of local problems unchanged, but larger nullspaces lead locally to more constraints — still potential load imbalance







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Parallel implementation

Parallel FEM solver with AMR

- experimental in-house code
- high order finite elements
- Poisson equation and linear elasticity
- C++ (object oriented) + MPI

p4est mesh manager for AMR

- rebalancing based on Z-curves
- ANSI C + MPI
- open-source (GPL)
- scalability reported for 1e5–1e6 cores

http://www.p4est.org

BDDCML equation solver

- Adaptive-Multilevel BDDC
- Fortran 95 + MPI
- open-source (LGPL)
- current version 2.5 (8/6/'15)
- tested on up to 65e3 cores and 2e9 unknowns

http://www.math.cas.cz/~sistek/

software/bddcml.html

P. Kůs









$$\begin{split} - \triangle u &= f \quad \text{on} \quad (0,1)^d \\ u &= \arctan \Bigl(s \cdot \Bigl(r - \frac{\pi}{3} \Bigr) \Bigr) \end{split}$$

- solution exhibits sharp internal layer
- $\blacksquare r$ is a distance from a given point
- $\blacksquare~s$ controls "steepness" of the layer

- Adaptivity tested for element orders 1-4 (showed order 1)
- \blacksquare Guided by exact solution, using H^1 semi-norm for error calculation





Iteration 3, mesh and solution

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Iteration 5, mesh and solution

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Iteration 8, mesh and solution

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Iteration 13, mesh and solution

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Iteration 18, mesh and solution.









Convergence of adaptivity in 2D, 8 subdomains

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Conclusions



Development of parallel FEM with AMR

level-1 hanging nodes simple to handle and interface with DD solver
disconnected and loosely coupled subdomains handled by detecting components of subdomain mesh

[P. K., J. Šístek, Coupling parallel adaptive mesh refinement with a nonoverlapping domain decomposition solver, Advances in Engineering Software 110, 34–54, 2017]

Future work

- improve the performance of the preconditioner by better subdomain shapes; Hilbert curves
- currently each refinement changes all subdomains not optimal
- multiple subdomains per compute node
- try to keep most of the subdomains intact and re-use part of the preconditioner work
- distribute created or changed subdomains to ensure load balancing
- promising starting point for embedded domain FEM
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Thank you for your attention.



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Mesh refinements - extra

- We want to focus the computational effort to troubled areas (e.g. singularities of electrostatic field, boundary layers in flow simulations, etc.)
- Algorithmic complexity of the software grows
- Different approaches
 - complete re-meshing
 - change of vertices positions(*r*-adaptivity)
 - element refinements (*h*-adaptivity)
 - different polynomial orders (*p*-adaptivity)
 - combination of both (*hp*-adaptivity)





benchmark



uniform mesh







- No global degrees of freedom assigned to hanging nodes
- Contributions to the local local stiffness matrix and RHS have to be adjusted to ensure continuity
- Works for higher-order basis functions as well
- Works for 2D and 3D for 1-irregular mesh and elements of same order
- Not numbering hanging nodes natural for p4est and BDDCML





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An abstract problem

$$u \in U : a(u, v) = \langle f, v \rangle \quad \forall v \in U$$

- $\blacksquare \ a \left(\cdot, \cdot \right)$ symmetric positive definite form on U
- $\blacksquare~\langle\cdot,\cdot\rangle$ is inner product on Hilbert space U
- U is finite dimensional space (typically finite element space)

Matrix form $u \in \mathbb{R}^n : Au = f$ A symmetric positive definite matrix

• A large, sparse, condition number $\kappa(A) = \frac{\lambda_{\max}}{\lambda_{\min}} = \mathcal{O}(1/h^2)$



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Iterative substructuring - extra





- $\Omega_1, \Omega_2 \dots$ subdomains (substructures), do not overlap • $\Gamma \dots$ interface
- The goal is to do as much work as possible locally

Reduced (Schur complement) problem on interface Γ

 $Su_{\Gamma} = g$

- \blacksquare *S* . . . Schur complement matrix
- $\blacksquare\ S$ much smaller than A
- solved by PCG

 $\blacksquare\ S$ never constructed, only its action on vector used

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x axis: Estimated element error y axis: Number of elements

Goal: refine prescribed fraction of elements in each step It might be too expensive to communicate error estimate of all elements to all processors

■ The first step is to communicate globally maximal element error





- On each processor, count elements in moderate number of error intervals (the same "bins" on all processors)
- 3 Communicate to the remaining processors this can be done









- 5 Refine elements of estimated error larger than threshold
- It can be none or all on some processors, but globally as required

Global sum

Continuity of basis functions - extra







refinement 1-irregularity rule



refinements forced by the previous step

- Standard H¹-conforming elements
- continuity of basis functions has to be enforced
- "gluing" of basis functions
- for hanging nodes, combinations of shape functions



 For higher-order hanging nodes, many shape functions may contribute









$(L_1,,, L_n)$	$L_{2},$	L_3 ,	$L_4)$	\leftrightarrow ($g_2,$	g_3, g	$_{5}, g_{6})$	
G_K	=	$\begin{bmatrix} 0\\0\\0\\0\\0 \end{bmatrix}$	1 0 0 0	0 1 0 0	0 0 0 0	0 0 1 0	0 0 0 1	
$A_K =$	$\begin{bmatrix} a \\ a \\ a \end{bmatrix}$	$L_1,$ \vdots $L_4,$	$L_1)$ $L_1)$	· · · .		a(L) $a(L)$	$\left[\begin{array}{c} 1,L_4 \end{array} ight] \\ \left[\begin{array}{c} 2\\ 2\\ 4\\ 4\\ \end{array} ight] L_4 ight)$	
$A = \sum_{K \in \mathcal{T}_k} G_K^T A_K G_K$								

Regular element

- \blacksquare Matrix G_K represents the relationship between local and global DOFs
- Local stiffness matrix distributed to the Global one
- Similarly for right-hand side





First option – constraints added to global matrix

- There are global degrees of freedom assigned to hanging nodes
- Matrix A assembled as in the regular case
- Solution would be discontinuous \rightarrow the global system has to be extended by constraints C
- The matrix \tilde{A} is rectangular, has to be modified before actual solution – various techniques







$$(L_1, L_2, L_3, L_4)^T = T_K(g_2, g_7, \mathbf{g_5}, g_8)^T$$
$$T_K = \begin{bmatrix} \mathbf{1} & 0 & 0 & 0\\ 0 & \mathbf{1} & 0 & 0\\ \mathbf{1/2} & 0 & \mathbf{1/2} & 0\\ 0 & 0 & 0 & \mathbf{1} \end{bmatrix}$$
$$\bar{A}_K = T_K^T A_K T_K$$
$$A = \sum_{K \in \mathcal{T}_k} G_K^T \bar{A}_K G_K$$

Second option – local change of basis

- No global degrees of freedom assigned to hanging nodes
- Matrix T_K used to modify the local stiffness matrix and RHS
- Corresponds to the construction of globally continuous basis function (depicted the one corresponding to *g*₅)





$$L_{1}, L_{2}, L_{3}, L_{4})^{T} = T_{K}(\mathbf{g}_{2}, g_{8}, g_{5}, g_{9})^{T}$$
$$T_{K} = \begin{bmatrix} \mathbf{1/2} & 0 & \mathbf{1/2} & 0\\ 0 & \mathbf{1} & 0 & 0\\ 0 & 0 & \mathbf{1} & 0\\ 0 & 0 & 0 & \mathbf{1} \end{bmatrix}$$
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Second option – local change of basis

- Works for higher-order basis functions as well
- Works for 2D and 3D for 1-irregular mesh and elements of same order
- \blacksquare Not numbering hanging nodes natural for p4est \rightarrow we use this approach

We use **non-overlapping** domain decomposition method, namely **BDDC**, a Balancing Domain Decomposition by Constraints. Experiments with two different libraries:

- Fempar library
 - A FORTRAN library for the development of Finite Element Multiphysics PARallel solvers
 - Developed at CIMNE, Barcelona, Spain by the group of Santiago Badia
 - Scales to hundreds of thousands of processor cores
 - Large project, includes all FEM machinery (space discretization, integration, assembling, physic-based preconditioners, ...)
 - A lot of the code has to be aware of the hanging nodes

BDDCML library

- A FORTRAN library, BDDC Multi Level
- Developed at IM AS CR by Jakub Šístek
- Only linear algebra solver
- Receives discrete system and some geometry information
- It is much easier to deal with hanging nodes on the interface

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Algorithm for generating constraints

1 classify interface to faces, edges and vertices — components independent

2 for each subdomain for each component

for each face

 select nodes on interface shared with neighbouring subdomain (generally larger set than the face under consideration) and detect components

■ select (in 3D) three nodes as corners from each such set as

- pick arbitrary node of the set
- \blacksquare find the first corner as the most remote node from the arbitrary node
- find the second corner as the most remote node from the first corner
- find the third corner as the node maximizing the are of the triangle
- 3 select corners as union of vertices and face-based selection
- 4 remove corners from edges and faces
- 5 use arithmetic averages on edges and faces

■ extension on face-based selection of corners [Šístek et al. (2011)]

amenable for parallelization

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- amenable for parallelization

- geometric information useful for the solver
- component detection crucial for robustness of corner selection
- amenable for parallelization

How to set up the scaling test for adaptivity?

- different trajectories of adaptive computations for changing number of cores
- what time to measure total solution time, time of the last problem from adaptive loop?
- setup of weak scaling tests unclear

Strong scaling tests on the final problem

- \blacksquare refinements are prescribed and always the same, not an adaptive run
- investigate the behaviour of DD on these nonstandard meshes

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Strong scaling in 2D (preliminary results)

Same 2D mesh with 28M DOFs, increasing number of subdomains

Illustrative mesh, 5000 DOFs, 5 subdomains

Strong scaling in 2D (preliminary results)

Same 2D mesh with 28M DOFs, increasing number of subdomains

Illustrative mesh, 5000 DOFs, 10 subdomains

Strong scaling in 2D (preliminary results)

Same 2D mesh with 28M DOFs, increasing number of subdomains

Illustrative mesh, 5000 DOFs, 20 subdomains.





























p4est vs. Metis



The same mesh partitioned from Z-curve (p4est) and graph (Metis)



subdomains by p4est

subdomains by Metis

- no significant differences in numbers of iterations
- more tests required in this direction