CSP lecture 17/18 winter semester – Problem Set 3

An *n*-ary operation on a set A is a mapping $A^n \to A$. The *n*-ary projection onto the *i*-th coordinate (on a set A) is the operation π_i^n defined by $\pi_i^n(a_1, \ldots, a_n) = a_i$ for any $a_1, \ldots, a_n \in A$.

An *n*-ary operation $f : A^n \to A$ is *compatible* with an *m*-ary relation $R \subseteq A^m$ if $f(\mathbf{r}_1, \ldots, \mathbf{r}_n) \in R$ (operation is applied coordinate-wise) whenever $\mathbf{r}_1, \ldots, \mathbf{r}_n \in R$. In other words, for any $m \times n$ matrix whose columns are in R, f applied to the rows of this matrix gives a tuple in R. In such a situation, we also say that R is compatible with f, or R is invariant under f.

An operation $A^n \to A$ is a *polymorphism* of a relational structure $\mathbb{A} = (A; ...)$ if it is compatible with all the relations in \mathbb{A} . The set of all polymorphisms of \mathbb{A} is denoted Pol(\mathbb{A}).

Problem 1. Observe that

- f: Aⁿ → A is compatible with every singleton unary relation {a}, a ∈ A, iff f(a,...,a) = a for all a ∈ A;
- the constant unary operation $c_a : A \to A$ (defined by $c_a(b) = a$ for any $b \in A$) is compatible with $R \subseteq A^n$ iff R contains the tuple (a, a, \ldots, a) .

Problem 2. Let A be a set. Prove that f is compatible with every relation on A if and only if f is a projection.

Problem 3. Let $\mathbb{A} = (A; ...)$ be a relational structure, $f \in Pol(\mathbb{A})$ a binary polymorphism and $g \in Pol(\mathbb{B})$ a ternary polymorphism. Then the 4-ary operation h defined by

$$h(a, b, c, d) = g(a, f(c, g(b, b, d)), c), \quad a, b, c, d \in A$$

is a polymorphism of \mathbb{A} as well. Try to formulate a general statement.

Problem 4. Find all unary and binary polymorphisms of the structure $\mathbb{A} = (\{0, 1\}; H, C_0, C_1)$ from Problem Set 1 (Problem 2 – HORN-SAT).

Problem 5. Find all unary and binary polymorphisms of the structure

 $\mathbb{A} = (\{0, 1\}; \text{all unary and binary relations})$

from Problem Set 1 (Problem 1 – 2-SAT). Find some nice nontrivial (= not a projection) polymorphism of \mathbb{A} .

Problem 6. Find all unary, binary, and ternary polymorphisms of the structure $\mathbb{A} = (\{0, 1\}; C_0, C_1, G_1, G_2)$ from Problem Set 1 (Problem 3 – LIN-EQ(\mathbb{Z}_2)).

A relation $R \subseteq A^m$ is *pp-definable* from $\mathbb{A} = (A; ...)$ if it can be defined from relations in \mathbb{A} by a pp-formula, that is, a formula which only uses conjunction, equality, and existential quantification. A relational structure $\mathbb{B} = (B; ...)$ is pp-definable from \mathbb{A} if A = B and each relation in \mathbb{B} is pp-definable from \mathbb{A} . We also say that \mathbb{A} pp-defines \mathbb{B} .

Problem 7. Prove that any relation pp-definable from \mathbb{A} is invariant under every polymorphism of \mathbb{A} .

Problem 8. Find all polymorphisms of the structure \mathbb{B} in Problem Set 2, Problem 4. (3–SAT). Hint: only projections; possible approach: (1) pp-define the four-ary relations of the form $R_{a,b,c,d} = \{0,1\}^4 \setminus \{(a,b,c,d)\}, (2)$ pp-define all four-ary relations (3) similarly, pp-define every relation, (4) use the results of other problems in this problem set.

Problem 9. Prove that a relation invariant under every polymorphism of A is pp-definable from A. Proof strategy:

(i) Denote $R = \{(c_{11}, \dots, c_{1k}), \dots, (c_{m1}, \dots, c_{mk})\}$

- (ii) Let $\mathbf{a}_1, \ldots, \mathbf{a}_n$ be a complete list of *m*-tuples of elements of *A* (i.e. $n = |A|^m$)
- (iii) Prove that the relation

 $S = \{(f(\mathbf{a}_1), \dots, f(\mathbf{a}_n)) : f \text{ is an } m \text{-ary polymorphism}\}\$

is pp-definable from \mathbb{A} (no need to use existential quantification)

- (iv) Existentially quantify over all coordinates but those corresponding to $(c_{11}, \ldots, c_{m1}), \ldots, (c_{1k}, \ldots, c_{mk})$
- (v) Prove that the obtained relation contains R (because of projections) and is contained in R (because of compatibility)

Problem 10. Observe that

- A pp-defines \mathbb{B} iff $\operatorname{Pol}(\mathbb{A}) \subseteq \operatorname{Pol}(\mathbb{B})$ and in such a case $\operatorname{CSP}(\mathbb{B}) \leq_P \operatorname{CSP}(\mathbb{A})$;
- any CSP over a two-element structure is polynomially reducible to 3-SAT
- if $\operatorname{Pol}(\mathbb{A}) \subseteq \operatorname{Pol}(\mathbb{B})$, then the proof of Problem 9 gives an explicit pp-formulas defining relations in \mathbb{B} from relations in \mathbb{A} .
- In particular, for structures from Problem Set 4, Problem 2, we get CSP(ℂ) ≤ CSP(𝔅). How large are the explicit formulas defining relations in ℂ from relations in 𝔅?