Instances of the Constraint Satisfaction Problem in Universal Algebra

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Fruitful cooperation UA \leftrightarrow Computer Science:

- Applications of UA to the complexity of the CSP (not in this lecture)
- ► Study of the CSP has a great impact on (parts of) UA
 - Surprisingly strong properties of quite general classes of algebras
 - New important classes of algebras discovered (FS - few subpowers, CS - congruence singular)
 - Fundamental new results about classic classes (CP - congruence permutable, CD - congruence distributive)
 - $\blacktriangleright \Rightarrow \mathsf{CSP}$ is not just a fashion

Apologies

- Picture of a naked mathematician
- Instance of the CSP
- Examples CSP in UA

$$CS
ightarrow CP
ightarrow FS
ightarrow CM
ightarrow Taylor \ \uparrow \qquad \uparrow \qquad \uparrow \ NU
ightarrow CD
ightarrow CSD(\wedge)$$

Useful technique - absorbing subalgebras (+ some news)

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Definition

A . . . finite idempotent algebra (always)

Instance of $CSP(\mathbf{A}) = finite set V + set C of constraints$ $Constraint = subalgebra of <math>\mathbf{A}^{I}$, where $I \subseteq V$ is the scope

Solution of the instance = mapping $f : V \to A$ such that $f_{|I} \in R_{|I}$ for every constraint $R \leq \mathbf{A}^{I}$

Example

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every scope is equal to V
\downarrow\downarrow
set of solutions = intersection of constraints
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to study CSP = to study intersection properties of subpowers

A is Taylor, if

HSP(**A**) (equivalently HS(**A**) ?Szendrei, ?Bulatov) doesn't contain a two-element algebra whose every operation is a projection \Leftrightarrow HSP(**A**) satisfies a nontrivial Maltsev condition \Leftrightarrow **A** has a Taylor term Taylor 77, i.e. a term *t* satisfying a set of identities in two variables *x*, *y* of the form

$$t(x, \cdot, \cdot, \dots) \approx t(y, \cdot, \cdot, \dots)$$

 $t(\cdot, x, \cdot, \dots) \approx t(\cdot, y, \cdot, \dots)$
 \dots
 $t(\cdot, \cdot, \dots, x) \approx t(\cdot, \cdot, \dots, y)$

 $\Leftrightarrow \mathsf{HSP}(\mathbf{A}) \text{ omits } \mathbf{1} \text{ Hobby, McKenzie } 88 \\ \Leftrightarrow \dots$

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Smooth theorem (Barto, Kozik, Niven 08)

Let **A** be a Taylor algebra, $R \leq \mathbf{A}^2$ subdirect and assume that $\exists k, l \in \mathbb{N}$ such that $(R^k \circ R^{-k})^l = A^2$. Then $\exists a \in A \quad (a, a) \in R$.

Corollary (Siggers, Kearnes, Marković, McKenzie 10)

A is Taylor iff **A** has a term t satisfying t(x, y, y, z) = t(y, z, x, x).

Proof.

F ... free algebra on $\{x, y, z\}$ *R* ... subalgebra of **F**² generated by (x, y), (y, z), (y, x), (z, x). Apply the theorem

Collapses of Maltsev conditions for finite algebras

Other intersection property (a generalization of Maróti, McKenzie 06):

Theorem (Barto, Kozik 09)

Let **A** be a Taylor algebra, p a prime, p > |A|, and $\emptyset \neq R \leq \mathbf{A}^p$. If R is invariant under cyclic shift of coordinates, then $\exists a \in A \ (a, a, ..., a) \in R$.

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Corollary

Let **A** be an algebra, p a prime, p > |A|. Then **A** is Taylor iff **A** has a cyclic term of arity p(i.e. a term satisfying $t(x_1, \ldots, x_p) = t(x_2, \ldots, x_p, x_1)$) Other intersection property (a generalization of Maróti, McKenzie 06):

Theorem (Barto, Kozik 09)

Let **A** be a Taylor algebra, p a prime, p > |A|, and $\emptyset \neq R \leq \mathbf{A}^p$. If R is invariant under cyclic shift of coordinates, then $\exists a \in A \ (a, a, ..., a) \in R$.

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Problem

Find a common generalization of this theorem and the smooth theorem.

Reading: L. Barto, M. Kozik: Absorbing subalgebras, cyclic terms and the constraint satisfaction problem

A is $CSD(\land)$, if HSP(A) (or HS(A)) doesn't contain a reduct of a module \Leftrightarrow HSP(A) is meet semi-distributive $\alpha \land \beta = \alpha \land \gamma \Rightarrow \alpha \land (\beta \lor \gamma) = \alpha \land \beta$ \Leftrightarrow HSP(A) omits 1, 2

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 \Leftrightarrow **A** has Willard terms

Definition

An instance of CSP(A) is (2,3)-minimal, if

- ▶ \forall three-element $I \subseteq V \exists R \subseteq \mathbf{A}^{K}$ in C such that $I \subseteq K$
- ▶ \forall at most two-element $I \subseteq V$ $\forall R \subseteq \mathbf{A}^{K}$, $R' \subseteq \mathbf{A}^{K'}$ such that $I \subseteq K, K'$ we have $R_{|I} = R'_{|I}$

Theorem (Barto, Kozik 09 Bulatov 09)

A is $\mathrm{CSD}(\wedge)$ iff every (2,3)-minimal instance of $\mathrm{CSP}(A)$ has a solution

Corollary

If **A** is $CSD(\wedge)$ and $R_1, \ldots, R_n \leq \mathbf{A}^n$ have the same binary projections, then $\cap R_i \neq \emptyset$

Example 2: $CSD(\land)$ algebras

WNU = operation f satisfying $f(x,...,x,y) = f(x,...,x,y,x) = \cdots = f(y,x,...,x)$

Corollary (Kozik, Valeriote)

A is $CSD(\wedge)$ iff **A** has WNUs of all arities ≥ 3 .

Proof.

Take V big enough. Let $\mathbf{F} = \text{free algebra in HSP}(\mathbf{A})$ over $\{x, y\}$. for every three element I we include one constraint $R_I \leq \mathbf{F}^I$, where $R_I = \langle (x, x, y), (x, y, x), (y, x, x) \rangle$ It is (2, 3)-minimal instance of $\text{CSP}(\mathbf{F}) \Rightarrow \exists$ solution $f : V \to F$. V is big $\Rightarrow \exists i, j, k \ f(i) = f(j) = f(k) = b$ $b \in R_{\{i,j,k\}} \Rightarrow \mathbf{A}$ has WNU of arity 3

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Reading: Doesn't exist yet :(Collapses of Maltsev conditions for finite algebras \mathbb{A} ... relational structure (on a finite set A)

 $\mathsf{Pol}(\mathbb{A})$... clone of all operations compatible with all relations in \mathbb{A}

Theorem (Geiger, Bodnarchuk, Kaluznin, Kotov, Romov 68)

 \forall finite algebra $A \exists A$ such that Pol(A) = Clo(A)

Definition

Finite A is finitely related, if \exists $\mathbb A$ with finitely many relations such that $\mathsf{Pol}(\mathbb A)=\mathsf{Clo}(A)$

Example

Algebras with near-unanimity term (by Baker-Pixley)

(Recall: near-unanimity = operation f satisfying $x = f(x, ..., x, y) = f(x, ..., x, y, x) = \cdots = f(y, x, ..., x)$)

Theorem (Barto 09)

If A is finitely related and HSP(A) is congruence distributive, then A has a near-unanimity term.

Proof.

Say A has at most k-ary relations, $Clo(\mathbf{A}) = Pol(A)$. $n \dots$ big enough natural number $V = A^n$ $F \leq \mathbf{A}^V \dots$ free algebra on *n*-generators (=*n*-ary operations) For every at most k-element $I \subseteq V$ we include the constraint $F_{|I|}$ Solutions of this instance = *n*-ary operations of **A**

Reading: L. Barto: Finitely related algebras in congruence distributive varieties have near unanimity terms More collapses of Maltsev conditions for finitely related algebras

Example 4: few subpowers \Rightarrow finitely related

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A has few subpowers, if $|\{R \le A^n\}| \le 2^{polynomial(n)}$ \Leftrightarrow subpowers of A have small generating sets \Leftrightarrow A has a cube term \Leftrightarrow ...

Example

Maltsev algebras, algebras with near-unanimity operation Few subpowers \Rightarrow HSP(**A**) is congruence modular

Theorem (Aichinger, Mayr, McKenzie 09)

Every finite algebra with few subpowers is finitely related.

Proof.

Use compact representations of subpowers developed for CSP by Dalmau, Bulatov; Berman, Idziak, Markovic, McKenzie, Valeriote, Willard

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Corollary

On a finite set, there is countably many clones with few subpowers (in particular, there is countably many Maltsev clones on a finite set).

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(2 years ago open for expansions of $\mathbb{Z}_8!$)

Conjecture (Valeriote's conjecture, Edinburgh conjecture)

If A is finitely related and HSP(A) is congruence modular, then A has few subpowers.

Reading: A. Bulatov, V. Dalmau: A simple algorithm for Mal'tsev constraints

J. Berman, P. Idziak, P. Markovic, R. McKenzie, M. Valeriote and R. Willard: Varieties with few subalgebras of powers

E. Aichinger, P. Mayr, R. McKenzie: On the number of finite algebraic structures

Congruence uniform $|x/\alpha| = |y/\alpha|$ Congruence singular $|x/\alpha||x/\beta| = |x/\alpha \lor \beta||x/\alpha \land \beta|$ strong malcev conditions.... Bulatov, Dalmau....

Reading: M. Dyer, D. Richerby: An effective dichotomy for the counting constraint satisfaction problem

Definition

B is an absorbing subuniverse of an algebra A, if

- ► *B* ≤ **A**
- A has a term t such that t(a₁,..., a_n) ∈ B whenever all but (at most) 1 of the a_i's are in B.

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Example

Singletons are absorbing subuniverses iff $\boldsymbol{\mathsf{A}}$ has a near-unanimity operation

Useful because:

- Algebras often have proper absorbing subuniverses
- Some connectivity properties of subpowers can be pushed inside absorbing subuniverses

Theorem

Let **A** be a Taylor algebra, $R \leq \mathbf{A}^2$ subdirect and assume that $\exists l \in \mathbb{N}$ such that $(R \circ R^{-1})^l = A^2$ and $R^l = A^2$. Then $\exists a \in A$ $(a, a) \in R$.

Proof.

- (0) Assume |A| > 1
- (1) Find a proper absorbing subuniverse B of **A**
- (2) Walk with B to find a proper absorbing subuniverse C of A such that $S \leq \mathbf{C}^2$ is subdirect, where $S = R \cap C^2$

(3) Prove that
$$(S \circ S^{-1})^{l'} = A^2$$
 and $S^{l'} = A^2$

Proof.

- (0) Assume |A| > 1
- (1) Find a proper absorbing subuniverse B of **A**
- (2) Walk with B to find a proper absorbing subuniverse C of A such that $S \leq \mathbf{C}^2$ is subdirect, where $S = R \cap C^2$
- (3) Prove that $(S \circ S^{-1})^{l'} = A^2$ and $S^{l'} = A^2$

Ad (1):

(Special case of) Absorption Theorem

Let **A** be a Taylor algebra, $R \leq \mathbf{A}^2$ subdirect, $R \neq A^2$, and $(R \circ R^{-1})^I = A^2$ for some *I*. Then **A** has a proper absorbing set.

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Studying CSPs in NL

Theorem (Barto, Kozik, Willard 11)

Let **A** be an algebra (no assumptions on **A**!). Then there exists $a \in A$ such that for any $n \in \mathbb{N}$ and any $B \leq C \leq \mathbf{A}^n$ such that B is subdirect in \mathbf{A}^n C contains all constant tuples B absorbs Cwe have $(a, a, \dots, a) \in B$

Ross has a problem

Let $Ab(\mathbf{A})$ be the set of all such a's. Easy to see that $Ab(\mathbf{A}) \leq \mathbf{A}$. What is this??????? (what is it good for? what characterizes this subalgebra?)

Problem

Is the following problem decidable? Input is a finite algebra **A** and a subset. Question is whether the subset is an absorbing subuniverse of **A**.

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Affirmative answer would generalize Maróti's result that NU is decidable

Problem

Is the following problem decidable? Input is a finite relational structure \mathbb{A} and a subset. Question is weather the subset is an absorbing subuniverse of $Pol(\mathbb{A})$.