

Equationally **nontrivial** algebras

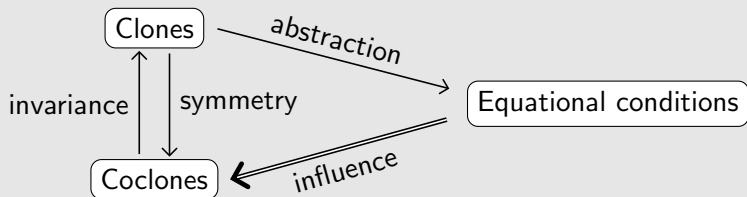
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Charles University in Prague

BLAST

Nashville, 15 Aug 2017

Our playground



Goal: Understand \Leftarrow in “the most general” case
equationally nontrivial clones, finite, idempotent

Q: Why understand \Leftarrow ?

A: to understand clones and coclones

Q: Why understand clones or coclones?

A: to understand algebras

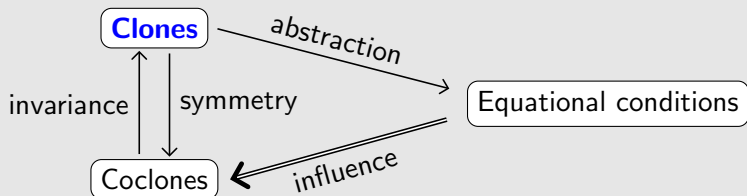
A: to understand symmetries

Q: Why the most general setting?

A: sometimes exactly what's needed

A: nontrivial results possible

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Definition

Operation on A = function $A^n \rightarrow A$, $n \geq 1$

Clone on A = set of operations on A closed under forming term operations

For each clone \mathbf{A} on A :

- ▶ for each $i \leq n$

$$(x_1, \dots, x_n) \mapsto x_i$$

is in \mathbf{A}

- ▶ if f, g are binary operations from \mathbf{A} , then

$$(x, y, z) \mapsto f(g(f(z, x), y), g(x, x))$$

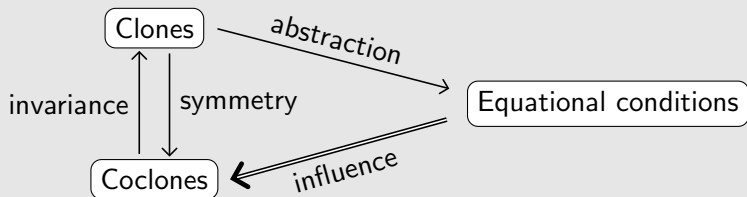
is in \mathbf{A}

Notation: For algebra \mathbf{A} , $\text{Clo}(\mathbf{A}) =$ all term operations of \mathbf{A}

Compute Clo(**A**)

- ▶ $(\{0, 1\}; \vee)$
- ▶ $(\{0, 1\}; \vee, \wedge)$
- ▶ $(\{0, 1\}; \text{majority})$
- ▶ $(\{0, 1\}; \vee, \wedge, \neg)$
- ▶ $(\mathbb{Z}_p; x + y)$
- ▶ $(\mathbb{Z}_p; x - y + z)$

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Q: Why clones? A: Abstraction of algebra

- ▶ carries important information – subuniverses, congruences
- ▶ for some purposes carries all necessary information (sometimes term equivalent algebras are essentially the same)

Q: Why understand algebras?

Too popular viewpoint

Group theory, Semigroup theory

- ▶ **group**: algebraic structure $\mathbf{G} = (G; \cdot, ^{-1}, 1)$ satisfying ...
- ▶ **permutation group**: when G happens to be a set of bijections, \cdot is composition, ...
- ▶ **monoid**: algebraic structure $\mathbf{M} = (M; \cdot, 1)$ satisfying ...
- ▶ **transformation monoid**: ...

Universal algebra

- ▶ **algebra**: any algebraic structure $\mathbf{Z} = (Z; \text{some operations})$

Rants

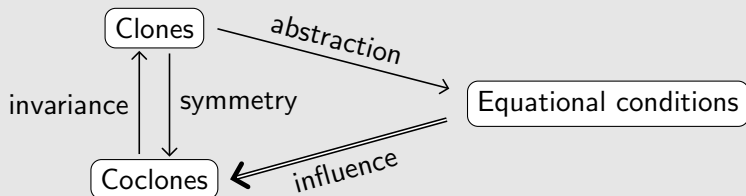
- ▶ Model theorist: models of purely algebraic signature, why do you avoid relations?
- ▶ Algebraist: groups are complicated enough, nothing interesting can be said about general algebras
- ▶ All: have you ever seen a 37-ary operation? You shouldn't study such a nonsense

Alternative viewpoint

	concrete	abstract
unary invert. symmetries	permutation group	group
unary symmetries	transformation monoid	monoid
higher arity symmetries	clone	abstract clone

- ▶ **permutation group**: Subset of $\{f : A \rightarrow A\}$ closed under composition and id_A and inverses...
can be given by a generating unary algebra
- ▶ **group**: Forget concrete mappings, remember composition
- ▶ **clone**: Subset of $\{f : A^n \rightarrow A : n \in \mathbb{N}\}$ closed under composition and projections
can be given by a generating algebra
- ▶ **abstract clone**: Forget concrete mappings, remember composition
aka variety, finitary monad over SET, Lawvere theory

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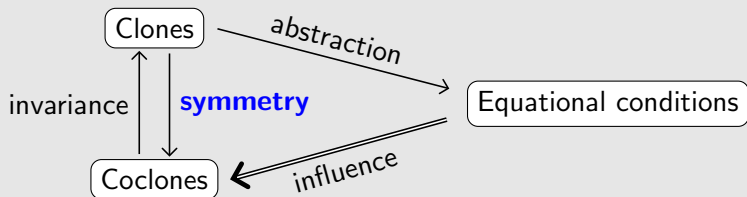
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Clones:

- ▶ classical algebraic structures \rightarrow general algebras \rightarrow clones
- ▶ permutation group \rightarrow clone

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Symmetries of relational structures

Definition

$f : A^n \rightarrow A$ is **compatible** with $R \subseteq A^k$

(f is a **symmetry** of R , f is a **polymorphism** of R ,
 R is **invariant under f**)

if $f(\mathbf{a}_1, \dots, \mathbf{a}_n) \in R$ whenever $\mathbf{a}_1, \dots, \mathbf{a}_n \in R$

Notation: For a set of relations \mathbb{A} ,

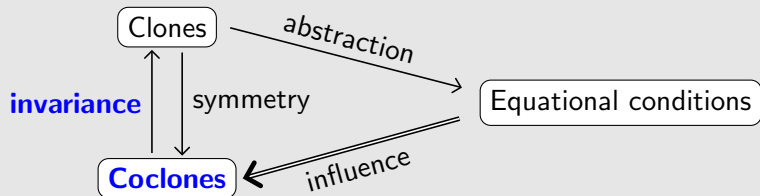
$\text{Pol}(\mathbb{A}) =$ all operations compatible with all relations in \mathbb{A}

Fact: $\text{Pol}(\mathbb{A})$ is a clone.

Compute $\text{Pol}(\mathbb{A})$

- ▶ $(\{0, 1\}; x \wedge y \rightarrow z, x \wedge y \rightarrow \neg z)$
- ▶ $(\{0, 1\}; \leq)$
- ▶ $(\{0, 1\}; \text{all binary relations})$
- ▶ $(\{0, 1, 2\}; \neq)$
- ▶ $(\mathbb{Z}_p; \text{vector subspaces of } \mathbb{Z}_p^3)$
- ▶ $(\mathbb{Z}_p; \text{affine subspaces of } \mathbb{Z}_p^3)$

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Invariant relations of a clone

Notation: For a set of relations \mathbb{A} ,
 $\text{Inv}(\mathbf{A})$ = all relations invariant under all operations in \mathbf{A}

Fact: It is closed under **pp-definitions** =
1st order definitions using $\exists, =$, and

Example: If binary R, S in $\text{Inv}(\mathbf{A})$, then

$$\{(x, y, z) : (\exists u)(\exists v) R(x, u) \text{ and } S(u, v) \text{ and } R(y, y)\}$$

is in $\text{Inv}(\mathbf{A})$

Definition

Coclone on A = set of (nonempty) relations on A closed under pp-definitions

Compute Cocolo(\mathbb{A})

- ▶ $(\{0, 1\}; x \wedge y \rightarrow z, x \wedge y \rightarrow \neg z)$
- ▶ $(\{0, 1\}; \leq)$
- ▶ $(\{0, 1\}; \text{all binary relations})$
- ▶ $(\{0, 1, 2\}; \neq)$
- ▶ $(\mathbb{Z}_p; \text{vector subspaces of } \mathbb{Z}_p^3)$
- ▶ $(\mathbb{Z}_p; \text{affine subspaces of } \mathbb{Z}_p^3)$

Compute $\text{Inv}(\mathbf{A})$

- ▶ $(\{0, 1\}; \vee)$
- ▶ $(\{0, 1\}; \vee, \wedge)$
- ▶ $(\{0, 1\}; \text{majority})$
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Theorem ([Geiger]; [Bodnarchuk et al.]

For finite A , Pol, Inv are (mutually inverse) bijections

$$\text{Clones on } A \quad \leftrightarrow \quad \text{Coclones on } A$$

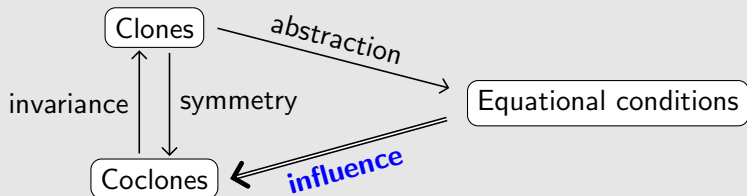
Remarks:

- ▶ $\text{Clo}(\mathbf{A}) = \text{Pol}(\text{Inv}(\mathbf{A}))$, $\text{Coclo}(\mathbb{A}) = \text{Inv}(\text{Pol}(\mathbb{A}))$
- ▶ Clones determined by invariant relations
- ▶ Coclones determined by symmetries
- ▶ Understanding clones = understanding coclones

Proof: Regard the set of n -ary operations in \mathbf{A} as $|A|^n$ -ary relation

From now on: \mathbf{A} clone, \mathbb{A} corresponding coclone

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Theorem ([Maltsev])

A contains a Maltsev operation ($m(x, y, y) = m(y, y, x) = x$)
iff

Each R in \mathbb{A} is rectangular ($\mathbf{ab}, \mathbf{ab}', \mathbf{a'b} \in R \Rightarrow \mathbf{ab}' \in R$)

Theorem ([Baker, Pixley])

A contains a majority operation

($m(x, x, y) = m(x, y, x) = m(y, x, x) = x$)

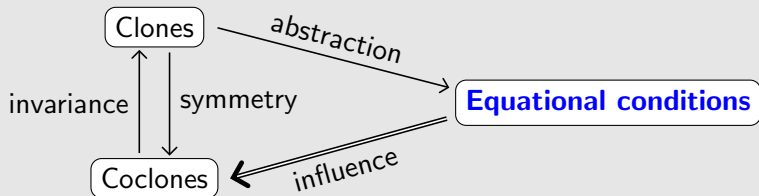
iff

Each R is determined by its projections to pairs of coordinates.

Does Clo(**A**) have Maltsev or majority operation?

- ▶ $(\{0, 1\}; \vee)$
- ▶ $(\{0, 1\}; \vee, \wedge)$
- ▶ $(\{0, 1\}; \text{majority})$
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Definition

Equational condition = condition of the form
“there exists operations ... satisfying equations”
(infinitely many operations or equations allowed)

Examples: the existence of a Maltsev term, the existence of a majority term

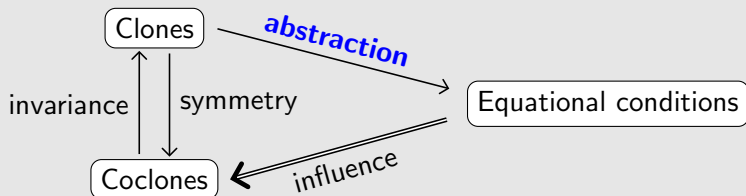
Remarks:

- ▶ Equational conditions are ordered by strength
- ▶ Equational condition is **nontrivial** if it is not satisfied in some clone
- ▶ Clone is **equationally nontrivial** if it satisfies some nontrivial equational condition

Is Clo(**A**) equationally trivial?

- ▶ $(\{0, 1\}; \vee)$
- ▶ $(\{0, 1\}; \vee, \wedge)$
- ▶ $(\{0, 1\}; \text{majority})$
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Abstraction: clones \rightarrow equational conditions

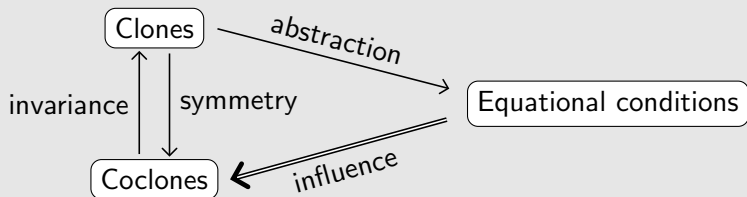
Consider two clones equivalent if they satisfy the same equational conditions

Abstraction: clone \rightarrow its equivalence class

Remarks:

- ▶ The set of equivalence classes is lattice ordered
- ▶ Simple formalization: It is the order induced by clone homomorphisms

Our playground



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Definition

Clone \mathbf{A} is **idempotent**

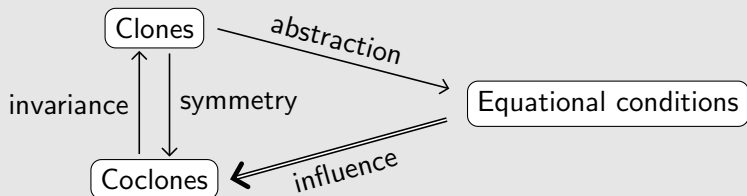
if $f(x, x, \dots, x) = x$ for each f in \mathbf{A}

\Leftrightarrow unary part of \mathbf{A} is trivial

Why this assumption?

- ▶ Complementary to group/semigroup theory
- ▶ Many useful equational conditions are idempotent
- ▶ Gives some information about general clones

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Equationally nontrivial clones

Theorem ([Birkhoff]; [Bodirsky]; [Bulatov])

For finite idempotent clone \mathbf{A} TFAE

(i) *Some part of the domain is purely combinatorial:*

*Formally $\mathbf{TRIV} \in \text{HSP}(\mathbf{A})$ (\mathbf{TRIV} is the clone of
projections on a 2-element set)*

Equivalently $\mathbf{TRIV} \in \text{HS}(\mathbf{A})$

(ii) *\mathbb{A} has the highest expressive power*

Formally, \mathbb{A} pp-interprets all finite relational structures

(iii) *\mathbf{A} is equationally trivial*

Definition (Just for this talk)

\mathbf{A} is a **Taylor clone** if it is finite, idempotent, and equationally **nontrivial**

Theorem ([Taylor])

For an idempotent clone \mathbf{A} TFAE

- ▶ \mathbf{A} is equationally nontrivial
- ▶ \mathbf{A} satisfies nontrivial height 1 equational condition involving a single operation symbol:

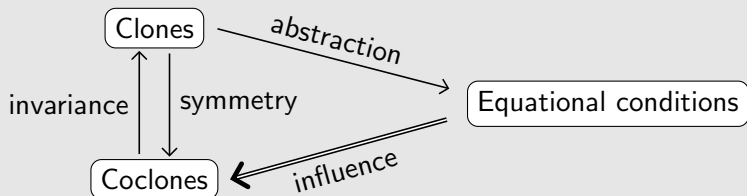
$$t(x, -, -, \dots) = t(y, -, -, \dots)$$

$$t(-, x, -, \dots) = t(-, y, -, \dots)$$

$$\vdots$$

$$t(\dots, -, -, x) = t(\dots, -, -, y)$$

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Sometimes it's exactly what's needed

Complexity of constraint satisfaction problems (CSPs)
... **large** part of computational complexity

Fixed finite template CSP

- ▶ a class of computational problems, one for each finite relational structure \mathbb{A}
CSP(\mathbb{A}) = membership in $\{\mathbb{X} : \mathbb{X} \rightarrow \mathbb{A}\}$
- ▶ tiny fraction of computational complexity (from global perspective)
- ▶ very broad class, base case for more optimistic goals

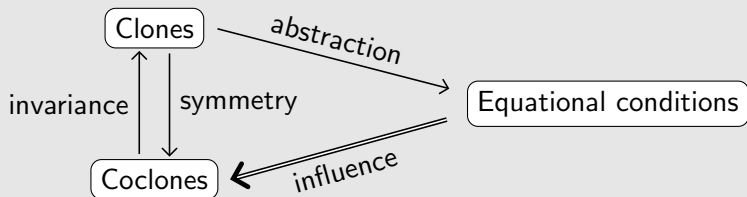
Theorem: [Bulatov, Jeavons, Krokhin] The complexity depends only on the position of \mathbf{A} in the order.

Consequence: If \mathbf{A} is not Taylor, the CSP(\mathbb{A}) is NP-complete.

THEOREM ([Bulatov]; [Zhuk])

If \mathbf{A} is Taylor, then CSP(\mathbb{A}) is solvable in polynomial time.

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Classic:

- ▶ Commutator theory [Smith], ...
- ▶ Tame congruence theory [Hobby, McKenzie], ...

More recent:

- ▶ Absorption theory [Barto, Kozik], ...
- ▶ Bulatov's theory
- ▶ Zhuk's theory

Theorem ([Maróti, McKenzie]; [Barto, Kozik, Niven]; [BK])

If \mathbf{A} is Taylor, then:

- ▶ *There exists $n > 1$ such that every symmetric n -ary relation in \mathbb{A} contains a constant tuple*
- ▶ *Every linked binary relation in \mathbb{A} contains a loop*
- ▶ *For every prime $p > |A|$, every cyclically-symmetric p -ary relation in \mathbf{A} contains a constant tuple*

Consequences for equations

Theorem ([MM], [Kearnes,Marković,McKenzie], [BK])

For an idempotent clone \mathbf{A} on finite set TFAE

(i) \mathbf{A} is Taylor

(ii) \mathbf{A} has a weak near unanimity operation of some arity $n > 1$

$$w(x, \dots, x, y) = w(x, \dots, x, y, x) = \dots = w(y, x, \dots, x)$$

(iii) \mathbf{A} has a 4-ary Sigger's operation

$$t(r, a, r, e) = t(a, r, e, a)$$

(iv) \mathbf{A} has a cyclic operation of each prime arity $p > |A|$

$$t(x_1, x_2, \dots, x_p) = t(x_2, \dots, x_p, x_1)$$

Note: (iii): A weakest nontrivial equation!

Digression to infinite

Weakest nontrivial equation for finite idempotent clones

Recall: From any set of idempotent operations f_1, \dots on a finite set A satisfying nontrivial equations

one can build a term operation s such that

$$s(r, a, r, e) = s(a, r, e, a)$$

Intuition: A is finite \Rightarrow composition is not sufficiently free

“Obviously”: It is impossible to find a weakest equations without the restriction to finite sets

Wrong!

Theorem ([Olšák])

For an idempotent clone \mathbf{A} TFAE

- ▶ *\mathbf{A} is equationally nontrivial*
- ▶ *\mathbf{A} contains a 6-ary t such that*

$$t(x, x, y, y, y, x) = t(x, y, x, y, x, y) = t(y, x, x, x, y, y)$$

End of digression

Absorption theory

Definition

Let $B \subseteq A$ be in \mathbb{A} .

B **absorbs** \mathbf{A} if \exists n -ary t in \mathbf{A} such that
 $t(A, B, B, \dots, B) \subseteq B$, $t(B, A, B, \dots, B) \subseteq B$, \dots

Theorem

If binary R in \mathbb{A} is subdirect and linked, then $B = A^2$ or \mathbf{A} has a proper absorbing set.

Theorem

*If B, C are minimal absorbing sets of \mathbf{A} ,
binary R in \mathbb{A} is subdirect and linked, and
 $R \cap (B \times C) \neq \emptyset$,
then $B \times C \subseteq R$.*

Bulatov's theory

$\mathbf{A} \rightarrow$ digraph on A , 3 types of edges: semilattice, majority, affine

Definition

(a, b) is a **semilattice edge** if \exists binary s in \mathbf{A} such that

$$s(a, b) = s(b, a) = b$$

(a, b) is a **majority edge** if ... [a more complex condition] ...

(a, b) is an **affine edge** if ... [even more complex condition] ...

Theorem

\mathbf{A} is Taylor iff the digraph of \mathbf{B} is connected for each subalgebra \mathbf{B} .

Theorem

If B, C are minimal affine & semilattice upward-closed subsets of A , binary R in \mathbb{A} is subdirect and linked, and

$$R \cap (B \times C) \neq \emptyset,$$

then $B \times C \subseteq R$.

Zhuk's theory

Strong structure theorems on “indecomposable” relations.

Crucial concepts: binary absorption, center, ...

Definition

A subset B of A is a center of \mathbf{A} if [such and such relation] is compatible with [something weird] and [some other condition].

Theorem

*If B, C are minimal centers of \mathbf{A} ,
binary R in \mathbb{A} is subdirect and linked, and
 $R \cap (B \times C) \neq \emptyset$,
then $R \cap (B \times C)$ is subdirect and linked.*

Methods

- ▶ absorption: heavily relational, lightly algebraic
- ▶ Bulatov: extremely algebraic, heavily relational
- ▶ Zhuk: heavily relational, lightly algebraic

Common: Some results look **very** similar
(different concepts, same assumptions and conclusions)

But: There is no clear connection, e.g. adding operations to a clone does not destroy absorption, can change colors, or centers

Also: Bulatov and Zhuk sometimes need to remove some operations (while remaining Taylor)

Work of:

- ▶ Zarathustra Brady (great write-up on his website)
- ▶ B + Bulatov + Kozik + Zhuk (last 2 weeks)

Definition

An idempotent clone is **Taylor minimal** if

- ▶ it is Taylor
- ▶ no proper subclone is Taylor

Fact: Each Taylor clone contains a Taylor minimal clone

Fun facts I (no guarantee!!!)

Theorem?

For a Taylor minimal clone \mathbf{A} and a unary B in \mathbb{A} TFAE

- ▶ *B is binary absorbing*
- ▶ *for each operation f in \mathbf{A} and each essential variable i , $f(A, \dots, A, B, A, \dots, A) \subseteq B$*
- ▶ *B is a cube term blocker*
- ▶ *B is semilattice & majority & affine upward-closed*

Fun facts II (no guarantee!!!)

Theorem?

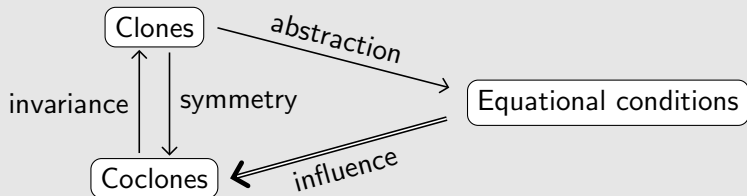
For a Taylor minimal clone \mathbf{A} and $B \subseteq A$, $(i) \Rightarrow (ii) \Rightarrow (iii) \dots$

- (i) B is ternary absorbing
- (ii) B is a center
- (iii) B is absorbing
- (iv) B is semilattice & affine upward-closed
- (v) B is “singleton ternary absorbing”

Fun facts III (no guarantee!!!)

- ▶ [Brady] Complete classification of 3-element affine-free Taylor minimal clones
- ▶ Complete classification on small domains possible \rightarrow source of examples
- ▶ Taylor minimality closed under H,S,P
- ▶ 5 omitting colors theorem, e.g.
 - ▶ semilattice & affine free \Leftrightarrow majority \Leftrightarrow near-unanimity
 - ▶ majority & affine free \Leftrightarrow 2-semilattice \Leftrightarrow binary cyclic
- ▶ 1 missing
- ▶ many questions...

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What next?

- ▶ Organize, unify, simplify, ...
- ▶ → (slightly) infinite
- ▶ → weighted
- ▶ → clonoids